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#### UNIVERSITY OF CALIFORNIA RIVERSIDE

Adiabatic Motion of Fault Tolerant Qubits

A Dissertation submitted in partial satisfaction of the requirements for the degree of

Doctor of Philosophy

 $\mathrm{in}$ 

Physics

by

David Edward Drummond

August 2014

Dissertation Committee:

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This work is dedicated to my wife and parents; their love and support made this

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#### ABSTRACT OF THE DISSERTATION

Adiabatic Motion of Fault Tolerant Qubits

by

David Edward Drummond

Doctor of Philosophy, Graduate Program in Physics University of California, Riverside, August 2014 Dr. Leonid P. Pryadko, Chairperson

This work proposes and analyzes the adiabatic motion of fault tolerant qubits in two systems as candidates for the building blocks of a quantum computer. The first proposal examines a pair of electron spins in double quantum dots, finding that the leading source of decoherence, hyperfine dephasing, can be suppressed by adiabatic rotation of the dots in real space. The additional spin-orbit effects introduced by this motion are analyzed, simulated, and found to result in an infidelity below the error-correction threshold. The second proposal examines topological qubits formed by Majorana zero modes theorized to exist at the ends of semiconductor nanowires coupled to conventional superconductors. A model is developed to design adiabatic movements of the Majorana bound states to produce entangled qubits. Analysis and simulations indicate that these adiabatic operations can also be used to demonstrate entanglement experimentally by testing Bell's theorem.

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### Chapter 1

# Introduction to Quantum Information

#### 1.1 Historical Overview of Quantum Computing

The fundamentally different physics at the atomic scale has interested and confounded scientists ever since the discovery of quantum mechanics in the first quarter of the twentieth century. While physicists did not fully reconcile the philosophical ramifications of the theory of quantum mechanics, they came to terms with it and made incredibly successful predictions in the decades that followed. However, in 1964 Bell made the philosophical differences between the classical and quantum worlds more concrete by conceiving his famous inequality that could be experimentally tested [7]. Subsequent experimental results were consistent with quantum mechanical predictions, suggesting that nanoscopic systems are inherently different from the macroscopic systems that our every-day intuition is based on.

This difference led physicists like Poplavskii and Feynman to believe that it was physically impossible to simulate a quantum mechanical system on a classical computer [35, 88], suggesting that a new type of computer constructed from a quantum system may be more powerful than a classical Turing machine. Thoughout the 1970's and 1980's many physicists established quantum information theory and elaborated on the specific requirements and potential of a quantum computer [119, 53, 51, 75, 9, 120, 29, 8]. Specifically, Deutsch showed that a quantum Turing machine could not only simulate a classical Turing machine, but could fundamentally out-perform conventional computers at certain tasks [25].

In the 1990's, physicists developed several algorithms that could be implemented on a quantum computer exponentially faster than classical computers [34, 27, 109, 46]. In particular, Shor discovered a quantum algorithm in 1994 capable of factoring large numbers in reasonable times, which was previously considered impossible and thus the crux of many of the cryptographic systems still in use today [106]. These discoveries spurred further interest into the theory of quantum information and the pursuit of a physical realization of a quantum computer.

Over the past 20 years, there have been many proposals for systems that could perform quantum computations and numerous experimental attempts at constructing them. While important developments in the field of quantum error correction made building a quantum computer more feasible, it has proven to be one of the most challenging goals in contemporary physics. At the heart of the difficulty is the delicate requirement of controlling individual particles, yet shielding them from their local environment. In recent years, experimental progress has been made through both incredible effort and advances in nano-technology and material science [80, 112, 28, 3, 30, 4]. Nonetheless, it is far from clear when, or whether, these advances will produce a quantum computer that surpasses the power of classical computers. At the same time, theorists have proposed innovative new approaches that use the concept of fault tolerance, building a quantum computer from systems that are inherently immune to certain environmental factors [89]. Rather than using a single two-state quantum system (i.e., qubits), these proposals encode quantum information in multiple qubits; some non-local degrees of freedom of a collection of qubits are tolerant to errors from any single local disturbances. Though these proposals often require more complicated setups or more advanced physics, fault tolerant quantum information offers an additional avenue towards a quantum computer.

This dissertation introduces and analyzes two proposals for fault tolerant quantum systems: the spin state of two electrons in a double quantum dot, and the particle parity of a topological state in a Majorana wire. In particular, this dissertation analyzes the gradual, adiabatic motion of their constituents that is crucial to both these proposals. The remainder of this introduction will briefly review the fundamentals of quantum information necessary to understand the two proposals discussed in the later chapters. For a more thorough review of quantum information, the interested reader should see the texts by Nielsen and Chuang[84] or Mermin[79].

#### **1.2** Fundamentals of Quantum Information

#### 1.2.1 Qubits

Rather than the classical bit that only takes the values of 0 or 1, the fundamental unit of quantum information, known as a quantum bit or qubit, can be a superposition state,  $|\psi\rangle$ , of two orthonormal basis states  $|0\rangle$  and  $|1\rangle$  of a quantum system

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle \tag{1.1}$$

where  $\alpha$  and  $\beta$  are complex coefficients that satisfy

$$|\alpha|^2 + |\beta|^2 = 1. \tag{1.2}$$

When the physical system is measured, the state is projected to either  $|0\rangle$  or  $|1\rangle$  with probability  $|\alpha|^2$  or  $|\beta|^2$ , respectively. The overall phase of  $|\psi\rangle$  is not physically relevant since the probabilities depend only on the absolute value of inner products. For example, the probability of measuring a qubit in the state  $|0\rangle$  is

$$\left|\langle 0|\psi\rangle\right|^2 = \left|\langle 0|(\alpha|0\rangle + \beta|1\rangle)\right|^2 \tag{1.3}$$

$$= \left| \alpha \langle 0|0 \rangle + \beta \langle 0|1 \rangle \right|^2 \tag{1.4}$$

$$= |\alpha|^2. \tag{1.5}$$

The choice of basis states used to describe a qubit is not unique; any qubit can be written in terms of any other orthonormal bases such as  $|\tilde{0}\rangle = (|0\rangle + |1\rangle)/\sqrt{2}$  and  $|\tilde{1}\rangle = (|0\rangle - |1\rangle)/\sqrt{2}$ . In practice, the preferred basis choice usually depends on the physical method of measurement. For example, when encoding information in the spin of electrons, the measurement device typically separates the spin in two directions, which are conventionally defined as  $|0\rangle$  and  $|1\rangle$ .

Mathematically, the superposition in Eq. (1.1) can be represented by a column vector

$$|\phi\rangle = \alpha|0\rangle + \beta|1\rangle \implies \phi = \left(\begin{array}{c} \alpha\\ \beta \end{array}\right). \tag{1.6}$$

This representation is generalized to multiple qubits through the Kronecker product. For example, two general qubits can be represented as

$$|\psi_{1}\rangle \otimes |\psi_{0}\rangle \sim \begin{pmatrix} \alpha_{1} \\ \beta_{1} \end{pmatrix} \otimes \begin{pmatrix} \alpha_{0} \\ \beta_{0} \end{pmatrix} = \begin{pmatrix} \alpha_{1} \begin{pmatrix} \alpha_{0} \\ \beta_{0} \end{pmatrix} \\ \beta_{1} \begin{pmatrix} \alpha_{0} \\ \beta_{0} \end{pmatrix} \end{pmatrix} = \begin{pmatrix} \alpha_{1}\alpha_{0} \\ \alpha_{1}\beta_{0} \\ \beta_{1}\alpha_{0} \\ \beta_{1}\beta_{0} \end{pmatrix}, \quad (1.7)$$

which maintains the binary convention from classical bits. For example, the state  $|10\rangle \equiv |1\rangle \otimes |0\rangle$ , which is the binary version of 2, is represented as

$$\begin{pmatrix} 0\\1 \end{pmatrix} \otimes \begin{pmatrix} 1\\0 \end{pmatrix} = \begin{pmatrix} 0\begin{pmatrix}1\\0\\1 \end{pmatrix}\\1\begin{pmatrix}1\\0 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} 0\\0\\1\\0 \end{pmatrix}, \qquad (1.8)$$

with only element 2 non-zero (with the count starting from 0). The Kronecker product and this binary rule both generalize as one would expect, for any number of qubits.

#### 1.2.1.1 Superposition

While the probabilistic interpretation of the coefficients in Eq. (1.1) may seem similar to a statistical ensemble of states, superposition is fundamentally different. In addition to the ratio of the coefficient magnitudes, the relative phase difference between  $\alpha$  and  $\beta$  also contains information that can lead to physically measurable results. For example, it is tempting to consider the state  $(|0\rangle + |1\rangle)/\sqrt{2}$  as simply an equal mixture of  $|0\rangle$  and  $|1\rangle$ . However, this description doesn't discern between the previous state and  $(|0\rangle - |1\rangle)/\sqrt{2}$ , despite the fact that these two states are distinct in another orthogonal basis. To make this distinction concrete, a superposition such as Eq. (1.1) is known as a pure state, while a classical mixture of states from an ensemble is known as a mixed state.

Part of the power of quantum algorithms comes from performing computations with superpositions of the basis states. For example, instead of separately calculating the values f(0) and f(1) of some computationally expensive function f(x), a quantum computer can find both values simultaneously by computing the function with a superposition of  $|0\rangle$  and  $|1\rangle$ . More so, a superposition of the two-qubit states  $|00\rangle$ ,  $|01\rangle$ ,  $|10\rangle$ , and  $|11\rangle$  can be used to simultaneously calculate four values and generally a quantum computer with n qubits could simultaneously calculate  $2^n$  values in principle. This exponential scaling, known as quantum parallelism, is at the heart of the incredible speed-up of quantum algorithms.

However, this description is slightly specious since only one of the results is accessible; when the result is measured the final state is projected to just one of the calculated values. This projection also makes it impossible to check a calculation in the middle of a computation since any intermediate measurement would ruin the superposition required in later steps. Fortunately, this restriction can be overcome with clever algorithms that solve a given problem with just a small subset of the function values. In practice, non-trivial algorithms like those designed by Grover or Shor don't solve problems deterministically, but rather find the correct solution in a relatively small number of attempts, with a probability that tends towards one [46, 106].

It should be noted that this "collapse" of a superposition to a single state upon measurement can also be useful. For example, if a secure cryptographic key is encoded in a superposition, any unintended eavesdropping will collapse the state and can be detected before transimitting sensitive messages [9]. Alternatively, a specific quantum state can be initialized by projectively measuring the relevant system until the desired state is found.

#### 1.2.1.2 Entanglement

Another feature of quantum mechanics that is responsible for the potential of quantum computers is the property that some multi-particle systems can only be described as a single "entangled" state. For example, the maximally entangled twoqubit system

$$|\Phi^{-}\rangle = \frac{|00\rangle - |11\rangle}{\sqrt{2}} \tag{1.9}$$

cannot be written as separable states (i.e., there are no  $|\psi_1\rangle$ ,  $|\psi_0\rangle$  such that  $|\Phi^-\rangle = |\psi_1\rangle \otimes |\psi_0\rangle$ ). Thus measuring one of the qubits automatically determines the outcome of the other, even if the qubits have been spatially separated after they were entangled [33]. This "spooky action at a distance" can be used to implement useful applications like quantum teleportation and cryptography that would be impossible on classical computers [34]. More so, entanglement plays a crucial role in many algorithms and other applications such as superdense coding [10, 55]. For this reason, preparing and demonstrating entanglement is an essential ingredient for any quantum computing proposal. The fundamental distinction between entangled and classical systems can be shown using Bell's theorem, as discussed further in Sec. 1.3.

#### 1.2.1.3 Bloch Sphere Visualization

A single pure state can be visualized by using the real parameters

$$\theta = 2 \tan^{-1}(|\beta|/|\alpha|) \tag{1.10}$$

$$\phi = \arg(\beta) - \arg(\alpha) \tag{1.11}$$

to write a qubit as

$$|\phi\rangle = \cos(\theta/2)|0\rangle + e^{i\phi}\sin(\theta/2)|1\rangle.$$
(1.12)

Thus  $\theta$  and  $\phi$  can serve as the polar and azimuthal angles, respectively, on the surface of what is known as the Bloch sphere, with  $|0\rangle$  and  $|1\rangle$  at the North and South pole, respectively (see Fig. 1.1). The Bloch sphere representation is quite helpful for understanding single qubits and how they evolve. Unfortunately, this simple visualization doesn't easily generalize to multiple qubits since any visualization would have to contain exponentially more information and account for the possibility of entanglement.



Figure 1.1: A single qubit can be represented as a point on the surface of the Bloch sphere. The latitude of the point describes the probability of measuring  $|0\rangle$  or  $|1\rangle$ ; Northern points have higher probabilities of measuring  $|0\rangle$ . The longitude describes the phase difference between the states  $|0\rangle$  and  $|1\rangle$  in the superposition.

#### 1.2.2 Quantum Operations

Despite the complexity of qubits compared to classical bits, any quantum computation can be performed with arbitrary accuracy with a small set of operations known as universal quantum operations [26, 22].

#### 1.2.2.1 Single Qubit Operations

The simplest operations involve a single qubit at a time, performing operations such as a bit flip that switches  $|0\rangle \leftrightarrow |1\rangle$  or a phase shift that switches  $|1\rangle \leftrightarrow -|1\rangle$ . More generally, any single qubit operation can be thought of as an a rotation of the states on the Bloch sphere. With this in mind, Euler's rotation theorem ensures that any rotation can be generated if one can perform arbitrary rotations about two orthogonal axes.

Mathematically, single qubit rotations can be generated by the orthogonal Pauli spin matrices

$$\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \qquad \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \qquad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, \qquad (1.13)$$

which are sometimes referred to as the X, Y, and Z operations. It is easy to see that the X operation performs the bit flip when acting on a general qubit

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \beta \\ \alpha \end{pmatrix}, \qquad (1.14)$$

effectively rotating the state  $\pi$  radians about the *x*-axis of the Bloch sphere. Similarly, the *Y* and *Z* operations correspond to  $\pi$ -rotations about their respective axes.

In general, any single-qubit operation corresponds to a unitary matrix and can be performed by allowing an appropriate Hamiltonian to evolve the qubit state. For example, a two-state quantum system driven by a harmonic electric field with a frequency near the splitting energy of the two states,  $\hbar\omega \sim \Delta E$ , causes the occupation probabilities of the two states to oscillate. This periodic behavior, known as Rabi oscillation, can be described as the operation of the unitary evolution

$$U(t) = e^{-i\omega t\sigma_x} \tag{1.15}$$

on the two-level system. The matrix representation of this exponential can be found by noting that even powers of the Pauli matrices are just the identity,  $\sigma_i^2 = 1$ , so the Taylor series can be written

$$e^{-i\omega t\sigma_x} = 1 + (-i\omega t)\sigma_x + \frac{(-i\omega t)^2}{2!} 1 + \frac{(-i\omega t)^3}{3!}\sigma_x + \dots$$
(1.16)

$$= \mathbb{1}\left(1 - \frac{(\omega t)^2}{2!} + ...\right) - i\sigma_x\left((\omega t) - \frac{(\omega t)^3}{3!} + ...\right)$$
(1.17)

$$= \lim_{\alpha \to \infty} \cos(\omega t) - i\sigma_x \sin(\omega t) \tag{1.18}$$

$$= \begin{pmatrix} \cos(\omega t) & -i\sin(\omega t) \\ -i\sin(\omega t) & \cos(\omega t) \end{pmatrix}.$$
 (1.19)

Thus an X rotation of any given angle can be performed in this system by oscillating an electric field for the appropriate duration. Similarly, Y and Z operations can be performed by finding or designing a system with the appropriate terms in the Hamiltonian.

#### 1.2.2.2 Multi-Qubit Operations

Single-qubit operations alone are not sufficient to perform most quantum algorithms; multi-qubit operations are necessary in order to construct a universal quantum computer [27]. Specifically most algorithms use "controlled" gates that perform an operation on a target qubit if and only if another control qubit is  $|1\rangle$ . For example, the controlled-not operation, also known as CNOT, given by

performs an X rotation (traditionally known as the NOT gate in the classical computing context) on the right, target qubit if and only if the left, control qubit is 1. Explicitly,

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} \alpha \\ \beta \\ 0 \\ 0 \\ 0 \end{pmatrix}, \quad (1.21)$$

while

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ \beta \\ \alpha \end{pmatrix}.$$
(1.22)

More generally, any single-qubit operation U can be applied to the target qubit, forming the controlled-U operation, represented in block matrix form as

$$\left(\begin{array}{cc}
1 & 0\\
0 & U
\end{array}\right).$$
(1.23)

Similarly, it is easy to check that the matrix

also represents a CNOT operation, but with the target and control qubits interchanged. These controlled operations add multi-particle interactions, enabling logical operations where the outcome of one qubit depends on the input of another qubit. Together with the single-qubit operations, any non-trivial controlled operation forms a universal set of quantum operations capable of performing any quantum computation [27].

The importance of controlled operations can be seen by observing that they can be used, together with single-qubit rotations, to entangle two qubits. For example, a qubit initialized in the state  $|0\rangle$ , then acted on by a  $\pi/2$ -angle Y rotation, given by

$$e^{-i\frac{\pi}{2}\sigma_y} = \frac{\sqrt{2}}{2} \begin{pmatrix} 1 & 1 \\ & \\ 1 & -1 \end{pmatrix},$$
 (1.25)

[cf. Eq. (1.16)], produces the superposition  $(|0\rangle - |1\rangle)/\sqrt{2}$ . If this state is used as the control qubit of a CNOT operation, with a target qubit initialized in the state  $|0\rangle$ , the resulting state is

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} \sqrt{2}/2 \\ 0 \\ -\sqrt{2}/2 \\ 0 \end{pmatrix} = \begin{pmatrix} \sqrt{2}/2 \\ 0 \\ 0 \\ -\sqrt{2}/2 \end{pmatrix}.$$
(1.26)

which represents  $(|00\rangle - |11\rangle)/\sqrt{2}$ , the maximally entangled state  $|\Phi^{-}\rangle$ .

The fact that controlled operations like CNOT can entangle qubits leads to a very useful consequence: together with entangled qubits, single-qubit rotations are sufficient to form a universal set of quantum operations [12]. For example, in the context of qubits formed by Ising anyons discussed in Ch. 3, universal quantum computation can only be achieved by complementing the topologically protected operations with entangled qubits in a process known as magic state distillation [11].

#### 1.2.3 Decoherence

Perhaps the most difficult aspect of constructing a quantum computer is to overcome decoherence, undesired interactions of the qubits with their environment. Almost all scalable physical systems are incredibly sensitive to environmental influences such as stray electromagnetic fields, disorder, and finite-temperature effects. One can consider the environment as a large set of unknown qubits that entangle themselves with the computational qubits and thus introduce unpredictable errors into the calculations.

In order to reliably execute algorithms, a quantum computer needs to perform several hundreds or thousands of operations without errors. Unfortunately, there is a difficult trade-off; systems with weaker environmental interactions tend to have longer operation times. Thus, longer decoherence times usually come at the cost of longer operation times.

#### 1.2.3.1 Fidelity

In order to track progress in preventing decoherence, a number of mathematical distances and measures have been introduced in quantum information [41]. While the details of these measures often depend on the relevant physical system, the most ubiquitous measure is known as "fidelity", which generally describes the similarity of two quantum states. This dissertation only considers the fidelity that compares a pure state before ( $\rho_0$ ) and after unitary evolution ( $U\rho_0U^{\dagger}$ ), which leads to the simpler definition [84]

$$F = \operatorname{Tr}(\rho_0 U \rho_0 U^{\dagger}), \qquad (1.27)$$

where  $\rho_0$  is the density matrix of the pure state, U is the time-evolution operator, and "Tr" is the trace operator. Since the initial state is not always known a priori, it is often useful to average the above fidelity over the subspace of possible pure states to calculate the average fidelity,  $\langle F \rangle$ , as discussed in more detail for the relevant systems in Sec. 2.4 and Sec. 4.4. Under these assumptions, the fidelity describes the ability of a system to preserve the value of its computational qubits. In practice, the fidelity can be quite close to 1, so the infidelity, 1 - F, is usually reported instead.

#### 1.2.3.2 Quantum Error Correction

Even with significant improvements in fidelity, quantum algorithms would be severely restricted without the ability to correct errors during computations. However, quantum error correction is much more complicated than its classical counterpart. A quantum error, which can involve both a bit-flip and a phase-shift, cannot be simply reversed between steps of an algorithm since measuring a superposition projects the state to either  $|0\rangle$  or  $|1\rangle$ , losing the information contained in the superposition and destroying any entanglement required for further steps. Furthermore, it is impossible to copy a general superposition prior to measuring for errors, due to a quantum mechanical rule known as the no-cloning theorem [120, 29].

Nevertheless, there are more advanced ways to correct errors by encoding the information of a single qubit into several entangled qubits using quantum error correction codes [107, 111]. While these codes require many more qubits, they make it possible to find the error syndrome, which can be used to detect and correct errors without directly measuring the state. Since these codes are self-correcting in a sense, quantum error correction codes are one form of fault tolerant qubits. For a more thorough review of the field of quantum error correction, the interested reader should read the overview by Gottesman [44].

Crucially, quantum error correction codes can be concatenated by encoding one logical qubit with several physical qubits, then using several of those coded qubits to create another layer of coded qubits, continuing in this manner indefinitely until the fidelity is sufficiently high. This repeated use of codes only helps if the increased overhead of operations needed for the coding scheme does not cause more errors than it corrects. While the details of this requirement depend on the specific system and type of errors possible, this important threshold for the infidelity per operation typically ranges from  $10^{-4}$  to  $10^{-2}$ . In practice, the infidelity needs to be well below the threshold to avoid needing several levels of error correction codes, and thus, a prohibitve number of qubits. For this reason, the goal of most quantum computing proposals is to perform quantum operations while remaining below this infidelity threshold, which is the focus of the proposal discussed in Ch. 2.

#### 1.3 Bell's Theorem

This section briefly reviews Bell's theorem and the relevant entanglement inequalities. It only covers the basic aspects needed for the proposal in Ch. 4; the interested reader should refer to the numerous works on the topic for a more comprehensive review [90, 78, 104, 79, 84].

#### 1.3.1 Bell's Inequality

To derive a simplified version of Bell's inequality, consider an experiment that separately measures each qubit of the state  $|\Phi^-\rangle = (|00\rangle - |11\rangle)/\sqrt{2}$  with one of three different methods, denoted *a*, *b*, and *c*. When both qubits are measured using the same method, the results are always the same. However, when the two qubits are measured using different methods, the results are completely uncorrelated. Thus the possible results for one qubit measurement depend on what method is used for the other qubit, even if the measurement events are well separated spatially (i.e., space-like).

Einstein, Podolsky, and Rosen famously objected to this type of non-local behavior [33], citing it as motivation for a more complete theory that removes the probabilistic nature of quantum mechanics by introducing "hidden variables". Hidden variable theories predict, with full certainty, the outcomes of different measurement methods on a single qubit, even though only one measurement at a time is possible. Bell's theorem states that any local hidden variable theory makes predictions that are inconsistent with quantum mechanics [7]. Thus any experiment that agrees with quantum mechanics rather than hidden variable theories, implies that the qubits in the system are entangled.

To see where the two theories are inconsistent, consider the interpretation of  $|\Phi^-\rangle$  in hidden variable theories. Instead of a pure state, it is viewed as one instance from a classical ensemble of states, prepared with different hidden variables. If an experimentalist could measure a single preparation with all three methods at once, the two qubits' results would match for each method. In this view the two qubits only seem uncorrelated when using different methods, but are actually correlated regardless of the method chosen. Thus the possible results of one qubit measurement don't depend on the method chosen for the other.

While this interpretation avoids non-local behavior, it replaces a superposition of states with a classical ensemble. Thus any single preparation in the ensemble must be either 0 or 1, not a more general superposition. Since there are three measurement methods, but only two possible outcomes, the pigeonhole principle states that at least two of the methods must give matching results. By defining  $P_{=}(a, b)$  as the probability that the results match when one qubit is measured with a and the other is measured with b, this statement can be written

$$P_{=}(a,b) + P_{=}(b,c) + P_{=}(a,c) \ge 1, \tag{1.28}$$

which is one version of Bell's inequality. Meanwhile, quantum mechanics predicts that this inequality is invalid for certain measurement methods, which demonstrates Bell's theorem. Specifically, each of the probabilities can be 1/4, leading to a violation of the inequality by 25%.

#### 1.3.2 Clauser-Horne-Shimony-Holt Inequality

While this inequality can be tested experimentally in principle, it requires method b to be tested for both qubits, which would be difficult to accomplish exactly in many proposals. Instead, consider the case where the left qubit of  $|\Phi^-\rangle$  is measured with either method  $L_1$  or  $L_2$ , while the right qubit is measured using either method  $R_1$ or  $R_2$ . Without superposition, each hidden variable preparation of the left qubit must have either  $L_1 = 0$  or  $L_1 = 1$ , meaning that measuring the left qubit with method  $L_1$ would yield 0 or 1, respectively. It is simpler to derive the inequality by considering the parity of these quantities so we use 1 and -1 for even and odd parity, respectively, for the remainder of this appendix. Thus, each preparation must have  $L_1$ ,  $L_2$ ,  $R_1$ , and  $R_2$ as either 1 or -1 according to the hidden variable interpretation.

Consider the quantities  $L_1 + L_2$  and  $L_1 - L_2$ ; either  $L_1 + L_2 = \pm 2$  and  $L_1 - L_2 = 0$ , or  $L_1 + L_2 = 0$  and  $L_1 - L_2 = \pm 2$ . This implies that the quantity

$$|(L_1 + L_2)R_1 + (L_1 - L_2)R_2| = 2 (1.29)$$

for each preparation since one of the terms vanishes in either case.

If an experimentalist could measure the qubits with more than one method at a time, this prediction could be tested directly. Instead, one must extract a statistical prediction that only requires a single measurement of each qubit for any given preparation. With that in mind, note that the expectation value of any constant is simply that constant, and any variable X satisfies  $|\langle X \rangle| \leq \langle |X| \rangle$  for any probability distribution. Applying these arguments to Eq. (1.29) yields the eponymous inequality first derived by Clauser, Horne, Shimony, and Holt[19]

$$|\langle L_1, R_1 \rangle + \langle L_2, R_1 \rangle + \langle L_1, R_2 \rangle - \langle L_2, R_2 \rangle| \le 2, \tag{1.30}$$

where  $\langle L, R \rangle = P_{=}(L, R) - P_{\neq}(L, R)$  is the expectation value for the combined parity of the left and right qubits when measured with methods L and R, respectively. Since each term only involves one measurement per qubit, it is possible to predict the left side of the inequality with quantum mechanics. For several measurement method combinations the left side can be as large as  $2\sqrt{2} \simeq 2.8$ , a violation of over 40%. Thus the quantum mechanical predictions are inconsistent with the local hidden variably theories and any experiment that violates the CHSH inequality negates the local hidden variable theories, demonstrating entanglement in the system.

## Chapter 2

# Double Quantum Dots Spatial Exchange Proposal

#### 2.1 Introduction

In recent years there has been great interest in the prospect of using scalable solid-state devices to implement qubits for potential applications such as quantum computation. One promising candidate for a qubit is a pair of electron spins in quantum dots, which forms a fault-tolerant subspace that is immune to collective decoherence [69]. This section briefly reviews double quantum dots, though the interested reader should refer to texts such as Marder [76] or Refs. [63, 97] for more details.

#### 2.1.1 Quantum Dot Background

The quantum dots in quantum computing proposals are built from semiconductor heterostructures that combine different semiconductor alloys with very similar properties to realize novel band structures not typically possible in nature. For example, since GaAs and AlAs have similar lattice sizes and structures, yet different band gaps, constructing a layer of  $Ga_{0.7}Al_{0.3}As$  on GaAs using molecular beam epitaxy leads to an abrupt change in the band structure at the junction (see Fig. 2.1B).



Figure 2.1: Band structure of a heterostructure showing the band bending and abrupt change at the junction in (A) and (B). The chemical potential can be tuned to occupy only a small inversion layer, shown in (C) and (D). Figure from Ref. [76].

With appropriate doping of the semiconductor, the chemical potential can be tuned to produce an inversion layer (see Fig. 2.1C), a small set of states that are bound for sufficiently small temperatures (typically on the order of a Kelvin). Thus the junction of the two alloys of the heterostructure forms a plane of confined electrons known as a two-dimensional electron gas (2DEG).

The electrons can be further confined by placing electrostatic gates on the top or bottom of the sample, parallel to the 2DEG. By placing a long, thin gate across the sample with a small opening, known as a quantum point contact, the path of electrons from a source to a drain can be narrowed so much that the resulting conductance through the opening is quantized in units of  $2e^2/h$ , allowing for very accurate detection of conductance [76].



Figure 2.2: Quantum dot confining a small area of a two-dimensional electron gas between two quantum point contacts and a plunger gate. Figure from Ref. [76].

A quantum dot, which is essentially a zero-dimensional quantum well, can be built with two quantum point contacts that confine an area of a 2DEG with a radius on the order of 100nm. Together with the quantum point contacts, another "plunger" gate is used to control the number of electrons in the interior of the dot using Coulomb repulsion, known as the Couloumb blockade in this context. With careful control of the voltages on the plunger and across the two sides of the sample (see Fig. 2.2), it is possible to control and detect the presence of a single electron in the dot [63, 97, 87].

#### 2.1.2 Double Quantum Dot Qubit

With the ability to confine a single electron in a quantum dot, it is theoretically possible to encode a qubit with the spin of that electron. Unfortunately, the spin of the electron would couple to any stray magnetic field in the system and quickly decohere before any meaningful operations could be performed.

However, Levy proposed a fault-tolerant modification of this qubit that is immune to a uniform magnetic field [69]. By using a pair of electrons shared between a


Figure 2.3: Double quantum dots confining a pair of electrons to encode a fault-tolerant qubit immune to a uniform magnetic field formed from the  $m_s = 0$  triplet and singlet spin states. From Ref. [87].

double quantum dot (see Fig. 2.3), with the  $m_s = 0$  triplet and singlet states as the logical basis states

$$|0\rangle = \frac{|\uparrow\downarrow\rangle + |\downarrow\uparrow\rangle}{\sqrt{2}}, \qquad |1\rangle = \frac{|\uparrow\downarrow\rangle - |\downarrow\uparrow\rangle}{\sqrt{2}}. \tag{2.1}$$

Since these states have an equal superposition of up and down spin states, a uniform magnetic field has the same effect on the spins of each dot, producing an innocuous overall phase. Nonetheless, each electron is still subject to the local hyperfine interaction from the nuclear spins of the semiconductor lattice of each individual dot, which are generally different for each quantum dot and thus lead to dephasing of the individual electron spins [14].

There have been several proposals to suppress this dephasing such as nuclear polarization [14, 59, 47], state narrowing [62], and spin-echo pulse correction [60, 123]. While improved coherence has been experimentally demonstrated using these techniques, the coherence times desired for applications have proven very difficult to achieve. For example, pumping methods have been used to partially polarize nuclei, but the nearly full polarization needed has yet to be achieved [62, 52, 113, 97].

Promising results have been shown using spin-echo sequences through the exchange interaction between two spins [87, 5, 117]. The exchange is controlled by lowering the tunneling barrier between the two quantum dots using quickly-controlled electric gates. This leads to Rabi oscillations; a single  $\pi$ -pulse corresponds to exchanging the two spins. Sequences of such pulses effectively couple both spins to the same average hyperfine interaction resulting in improved coherence times. While several echo sequences can be performed using this exchange, it is likely that the electrical gate noise and spatial variations in the Overhauser field remain the dominant sources of dephasing [5].

Rather than relying on interdot tunneling, the current proposal uses a spatial exchange of the two quantum dots, allowing the two electrons to traverse the same path, spending the same time coupled to the local nuclei, as shown in Figure 2.4. Compared to exchange via tunneling, ideally, this should eliminate the effect of the electrical gate noise. On the other hand, in the presence of spin-orbit coupling, the motion of electrons may introduce additional errors. In the remainder of this chapter, we analyze how these errors depend on the parameters of the motion and discuss the constraints and possible parameters for potential implementation of this proposal.

It should be noted that a similar setup with spatial exchange of electrostaticallydefined quantum dots has been discussed in relation to holonomic quantum computation [42, 105]. However, the corresponding coherence estimates have been done in the absence of a magnetic field. This proposal focuses on the parameter range characteristic of double-quantum dot qubits and accounts for typical magnetic fields of order  $\geq 0.1$  T.



Figure 2.4: Suggested electrode geometry for a rotating double-quantum-dot qubit with top and bottom gates in different shades of gray. Exchange gates via real space rotation, as opposed to tunneling, are expected to strongly reduce the qubit sensitivity to charge noise.

The outline of the remainder of this chapter is as follows. In Sec. 2.2, the Hamiltonian of the double-quantum dot qubit with spatial exchange is defined. The effective spin-only Hamiltonian for a single electron in a moving quantum dot is derived in Sec. 2.3, and the single-qubit fidelity associated with a sequence of double-dot rotations is found in Sec. 2.4. Simulations of a possible protocol are shown in Sec. 2.5, and the constraints and corresponding characteristic time and distance scales are discussed in Sec. 2.6, followed by a summary in Sec. 2.7.

# 2.2 Double Quantum Dot Setup

Consider a qubit formed by a pair of quantum dots electrostatically defined in a III-V semiconductor (e.g., GaAs/AlGaAs) heterostructure using a system of top and bottom gates similar to that discussed in the previous section and illustrated in Fig. 2.4, with the parameters of the dots similar to the experiments in Refs. [87, 5, 117]. Specifically, each dot contains a single electron, with a typical dot-size quantization energy  $\hbar \omega_d \sim 1 \text{meV}$ . The qubit is defined as the subspace of the singlet and  $m_s = 0$ triplet states of the two electrons. The triplet degeneracy is removed by a uniform, constant magnetic field  $B_0$  of at least  $\sim 0.1$ T, applied perpendicular to the sample, which creates a Zeeman gap of  $\Delta \sim 2.5 \ \mu \text{eV}$ . The electrons interact with  $N \sim 10^6$  spins of the lattice nuclei, leading to a different local hyperfine interaction for each electron. This can be approximated as a Zeeman interaction with a random, fluctuating, nonuniform magnetic "Overhauser" field  $B_N \sim 1$ mT.

To prevent dephasing, both quantum dots will be moved along the same trajectory in a time much shorter than the relaxation time of the nuclear spins,  $t_{\rm nuc} \sim 100 \ \mu s$ , so it is assumed that the Overhauser field is quasi-static. In order to reduce the sensitivity to charge noise the distance between the dots must be significantly greater than the size of each quantum dot,  $a \sim 100$  nm. Due to this spatial separation between the dots, the Hamiltonians of each electron can be treated separately,

$$H = H_{\rm d}(\mathbf{r}_0) + H_{\rm Z} + V_{\rm Z}(\mathbf{r}, t) + H_{\rm SO} + V(\mathbf{r}, t).$$
(2.2)

Here the dot Hamiltonian is given by

$$H_{\rm d}(\mathbf{r}_0) = \frac{p^2}{2m} + U(\mathbf{r} - \mathbf{r}_0(t)), \qquad (2.3)$$

with the canonical momentum  $\mathbf{p} = \mathbf{P} + e\mathbf{A}/c$  and the confining potential U centered at  $\mathbf{r}_0 \equiv \mathbf{r}_0(t)$ ; the Zeeman Hamiltonians for the externally-applied and Overhauser fields are respectively

$$H_{\rm Z} = \frac{g\mu_{\rm B}}{2} \mathbf{B}_0 \cdot \boldsymbol{\sigma},\tag{2.4}$$

$$V_{\rm Z}(\mathbf{r},t) = \frac{g\mu_{\rm B}}{2} \mathbf{B}_{\rm N}(\mathbf{r},t) \cdot \boldsymbol{\sigma}, \qquad (2.5)$$

and the spin-orbit Hamiltonian is given by

$$H_{\rm SO} = \sigma^i C^{ij} p^j. \tag{2.6}$$

The last term,  $V(\mathbf{r}, t)$ , accounts for additional effects originating from disorder, variation of the dot potential as it moves due to imperfections of the confining potential, as well as phonons. The spin orbit coupling coefficients  $C_{ij}$  in Eq. (2.6) incorporate both Dresselhaus (originating from the lack of the inversion symmetry of the lattice) with  $C^{yy} = -C^{xx} = \beta$ , and Rashba terms (structural inversion asymmetry due to the quantum well) with  $C^{xy} = -C^{yx} = \alpha$ . This specific form assumes that the growth of the semiconductor heterostructure and the quantum well asymmetry are in the positive z direction, [108, 121] so all the matrix elements that involve z are zero.

For numerical estimates the effective electron mass used is  $m \sim 0.067 m_{\rm e}$ , and the spin-orbit parameter is assumed to satisfy  $\alpha \simeq \beta$  with values ranging from  $10^3$  to  $10^4$ m/s, as appropriate for typical GaAs heterostructures.[108] Based on the above values,  $m\beta^2 \ll g\mu_{\rm B}B \ll \hbar\omega_{\rm d}$ , so terms quadratic in the spin-orbit coupling can be ignored for the following analysis.

# 2.3 Effective Single-Dot Hamiltonian

Consider the moving reference frame of the dot using the translation operator

$$\Psi(t) \to e^{-\frac{i}{\hbar} \mathbf{P} \cdot \mathbf{r}_o(t)} \Psi(t), \qquad (2.7)$$

$$H \to e^{\frac{i}{\hbar} \mathbf{P} \cdot \mathbf{r}_o(t)} H e^{-\frac{i}{\hbar} \mathbf{P} \cdot \mathbf{r}_o(t)} - \mathbf{v}_0(t) \cdot \mathbf{P}, \qquad (2.8)$$

which introduces the additional term proportional to the dot's velocity,  $\mathbf{v}_0 \equiv \dot{\mathbf{r}}_0$ , and removes the  $\mathbf{r}_0$  from the confining potential,  $U(\mathbf{r} - \mathbf{r}_0) \rightarrow U(\mathbf{r})$ . This also affects the vector potential,  $\mathbf{A}(\mathbf{r}) \rightarrow \mathbf{A}(\mathbf{r} + \mathbf{r}_0)$ , in the dot Hamiltonian and spin-orbit term, which can be reversed with an appropriate gauge transformation. Using the symmetric gauge,  $\mathbf{A}(\mathbf{r} + \mathbf{r}_0) = \frac{1}{2}\mathbf{B} \times (\mathbf{r} + \mathbf{r}_0)$ , and transforming  $\mathbf{A} \to \mathbf{A} + \nabla f$  with  $f = -\frac{1}{2}\mathbf{r} \cdot (\mathbf{B} \times \mathbf{r}_0)$ results in

$$\Psi \to \Psi \exp(-\frac{ie}{\hbar c}f),$$
 (2.9)

$$\mathbf{P} \to \mathbf{P} - \frac{e}{c} \nabla f. \tag{2.10}$$

These transformations introduce two additional terms in the time-dependent Schrödinger equation  $i\hbar\partial\Psi/\partial t = H\Psi$ . The first term,  $-(e/2c)\mathbf{v}_0 \cdot (\mathbf{B} \times \mathbf{r}_0)$ , arises when Eq. (2.10) is substituted into the  $-\mathbf{v}_0(t) \cdot \mathbf{P}$  term in Eq. (2.8). The second term,  $-(e/2c)\mathbf{r} \cdot (\mathbf{B} \times \mathbf{v}_0)$ , appears on the left hand side as a result of taking the time derivative of the exponent in Eq. (2.9). These two terms can be moved to one side of the equation and combined using the cyclic property of the mixed product:

$$-\frac{e}{2c}\mathbf{v}_0\cdot(\mathbf{B}\times\mathbf{r}_0) + \frac{e}{2c}\mathbf{r}\cdot(\mathbf{B}\times\mathbf{v}_0) = -\frac{e}{2c}\mathbf{v}_0\cdot\mathbf{B}\times(\mathbf{r}+\mathbf{r}_0)$$
(2.11)

which is precisely the vector potential term in  $-\mathbf{v}_0 \cdot \mathbf{p}$ . This leads to the moving-frame Hamiltonian

$$H = H_{\rm d} + H_{\rm Z} + H_{\rm SO} + V(\mathbf{r} + \mathbf{r}_0(t), t) + V_{\rm Z}(\mathbf{r} + \mathbf{r}_0(t), t) - \mathbf{v}_0 \cdot \mathbf{p}.$$
 (2.12)

Following Golovach et al.[43], a canonical transformation is performed,  $H \to e^S H e^{-S} \simeq (1+S)H(1-S) = H + [S,H]$ , where S is anti-Hermitian and chosen to eliminate the original spin-orbit term. Splitting  $S = S_0 + S_1$  such that

$$[S_0, H_d] + H_{\rm SO} = 0, (2.13)$$

$$[S_1, H_d] + [S_0, H_Z] = 0, (2.14)$$

and choosing

$$S_0 = \frac{im}{\hbar} \sigma^i C^{ij} r^j, \qquad (2.15)$$

satisfies Eq. (2.13). This can be verified, noting that  $S_0$  has no momentum dependence so it clearly commutes with the confining potential, and

$$\begin{bmatrix} \frac{im}{\hbar} \sigma^i C^{ij} r^j, \frac{p^2}{2m} \end{bmatrix} = \frac{i}{\hbar} \sigma^i C^{ij} [r^j, p^k] p^k$$
$$= -\sigma^i C^{ij} p^j.$$
(2.16)

With  $S_0$  known, Eq. (2.14) is used to define  $S_1$ ,

$$[H_d, S_1] = [S_0, H_Z] = \frac{img\mu_B}{2\hbar} C^{ij} r^j [\sigma^i, \sigma^k] B_0^k$$
$$= -\frac{mg\mu_B}{\hbar} C^{ij} r^j \epsilon^{ikl} B_0^k \sigma^l$$
$$= -\frac{g\mu_B}{2} \mathbf{Q} \cdot (\mathbf{B}_0 \times \boldsymbol{\sigma})$$
$$= \frac{g\mu_B}{2} (\mathbf{B}_0 \times \mathbf{Q}) \cdot \boldsymbol{\sigma}, \qquad (2.17)$$

where  $Q^i \equiv (2m/\hbar)C^{ij}r^j$ . This equation can be written in terms of the electron's orbital states  $|n\rangle$  in the dot potential,

$$\langle n | [H_d, S_1] | m \rangle = (S_1)_{nm} (E_n - E_m)$$
$$= \frac{g\mu_{\rm B}}{2} \boldsymbol{\sigma} \cdot (\mathbf{B}_0 \times \langle \mathbf{Q} \rangle_{nm}).$$
(2.18)

As long as the relevant dot quantization energies are non-degenerate,  $E_n \neq E_m$ , this can be written as

$$(S_1)_{nm} = \frac{g\mu_{\rm B}}{2} \boldsymbol{\sigma} \cdot \frac{(\mathbf{B}_0 \times \langle \mathbf{Q} \rangle_{nm})}{E_n - E_m}$$
$$= \frac{g\mu_{\rm B}}{2} \boldsymbol{\sigma} \cdot \mathbf{W}_{nm}, \qquad (2.19)$$

where  $W_{nm}$  is defined as

$$\mathbf{W}_{nm} \equiv \frac{\mathbf{B}_0 \times \langle \mathbf{Q} \rangle_{nm}}{E_n - E_m}.$$
(2.20)

Expanding the canonical transformation to first order in the spin-orbit parameter contained in  $S_0$  and  $S_1$ , the transformed Hamiltonian is

$$\bar{H} \simeq H + [S_0, V_Z] - [S_0, \mathbf{v}_0 \cdot \mathbf{p}] + [S_1, V] + [S_1, H_Z] - [S_1, \mathbf{v}_0 \cdot \mathbf{p}] + [S_1, V_Z] \quad (2.21)$$

Using the definition of  $S_0$ , the first two commutators are

$$\frac{img\mu_{\rm B}}{2\hbar}C^{ij}r^{j}[\sigma^{i},\sigma^{k}]B_{\rm N}^{k} = -\frac{mg\mu_{\rm B}}{\hbar}C^{ij}r^{j}(\epsilon^{ikl}B_{\rm N}^{k}\sigma^{l})$$
$$= \frac{g\mu_{\rm B}}{2}(\mathbf{B}_{\rm N}\times\mathbf{Q})\cdot\boldsymbol{\sigma}$$
(2.22)

and

$$-\frac{im}{\hbar}\sigma^i C^{ij}[r^j, p^k]v_0^k = m\sigma^i C^{ij}v_0^j = \frac{1}{2}\mathbf{Q}_{\mathbf{v}} \cdot \boldsymbol{\sigma}, \qquad (2.23)$$

respectively, with  $Q_v^i \equiv 2mC^{ij}v_0^j$  defined analogously to  $Q^i$ . One can now define the effective spin Hamiltonian by projecting onto the orbital ground state,  $H_S \equiv \langle 0|\bar{H}|0\rangle$ . This makes it possible to express the remaining commutators in the transformed Hamiltonian using the definition of  $S_1$ . The first commutator involving  $S_1$  in Eq. (2.21) is simplified by explicitly writing out the commutator and inserting a complete set of states,

$$\langle 0|S_1 V|0\rangle - \langle 0|VS_1|0\rangle = \sum_{n>0} \langle 0|S_1|n\rangle \langle n|V|0\rangle - \langle 0|V|n\rangle \langle n|S_1|0\rangle = \sum_{n>0} (S_1)_{0n} V_{n0} - V_{0n}(S_1)_{n0}.$$
 (2.24)

Assuming V has no momentum dependence, it commutes with  $S_1$  and the order of the second term above can be reserved. Since  $S_1$  is anti-Hermitian, while V is Hermitian,

$$\langle 0|S_1 V|0\rangle - \langle 0|VS_1|0\rangle = 2\sum_{n>0} (S_1)_{0n} V_{n0}$$
  
$$= g\mu_{\mathrm{B}}\boldsymbol{\sigma} \cdot \sum_{n>0} \mathbf{W}_{0n} V_{n0}.$$
 (2.25)

This technique can be used on the remaining terms in Eq. (2.21). The second term involving  $S_1$  in Eq. (2.21) vanishes because

$$(H_{\rm Z})_{n0} = \frac{g\mu_{\rm B}}{2}\boldsymbol{\sigma} \cdot \langle n | \mathbf{B}_0 | 0 \rangle = 0.$$
 (2.26)

However, the final commutator in Eq. (2.21) contains the terms

$$(S_1)_{0n}(V_Z)_{n0} - (V_Z)_{0n}(S_1)_{n0}, (2.27)$$

which do not commute because they contain two spin terms, but can be treated using the spin identity

$$(\boldsymbol{\sigma} \cdot \mathbf{a})(\boldsymbol{\sigma} \cdot \mathbf{b}) = \mathbf{a} \cdot \mathbf{b} + i\boldsymbol{\sigma} \cdot (\mathbf{a} \times \mathbf{b}).$$
(2.28)

The first term in Eq. (2.27) becomes

$$\left(\frac{g\mu_{\rm B}}{2}\right)^2 \left\{ \mathbf{W}_{0n} \cdot (\mathbf{B}_{\rm N})_{n0} + i\boldsymbol{\sigma} \cdot [\mathbf{W}_{0n} \times (\mathbf{B}_{\rm N})_{n0}] \right\},\tag{2.29}$$

and the second term looks quite similar, except the anti-commutator of the cross product causes the spin dependent term to cancel with the one above, while the spin-indepedent terms is doubled. The Hamiltonian contains several of these spin-independent terms that can be taken as constants. The effective spin Hamiltonian, up to a constant, can now be written simply as

$$H_S = \frac{1}{2}\hbar \left[\boldsymbol{\omega}_0 + \boldsymbol{\omega}_1(t)\right] \cdot \boldsymbol{\sigma}, \qquad (2.30)$$

where

$$\boldsymbol{\omega}_0 = \frac{g\mu_{\rm B}}{\hbar} \mathbf{B}_0, \tag{2.31}$$

is the Larmor frequency, and the time-dependent term  $\omega_1 \equiv \omega_1(t)$  is

$$\boldsymbol{\omega}_{1} = \frac{g\mu_{\mathrm{B}}}{\hbar} \left[ \mathbf{B}_{\mathrm{N}} + (\mathbf{B}_{\mathrm{N}})_{0n} \times (\mathbf{Q})_{n0} \right] + \frac{4g\mu_{\mathrm{B}}}{\hbar} \mathbf{W}_{0n} V_{n0} + \frac{1}{2\hbar} \mathbf{Q}_{\mathrm{v}} - \frac{2g\mu_{\mathrm{B}}}{\hbar} \mathbf{W}_{0n} (\mathbf{v}_{0} \cdot \mathbf{p}_{n0}), \qquad (2.32)$$

where the index n is implicitly summed over all the excited states of the dot. If the phonons are ignored and the Overhauser fields are approximated as static, the time dependence of the terms in the first line of Eq. (2.32) comes only from the position, parameterized by the known trajectory of the dot,  $\mathbf{r}_0(t)$ . Similarly, the time-dependence of the terms in the second line comes from both the position  $\mathbf{r}_0(t)$  and the dot velocity  $\mathbf{v}_0(t)$ , [in fact, these terms are all linear in components of  $\mathbf{v}_0(t)$ ]. This simple spin Hamiltonian is the key result of this derivation; it is correct to linear order in the spinorbit couplings. It should be noted that including cubic spin-orbit terms in the original Hamiltonian introduces additional terms proportional to  $\mathbf{v}_0^2(t)$  and  $\mathbf{v}_0^3(t)$  in Eq. (2.32), but these terms are smaller by at least an order of magnitude [64] and the general spin Hamiltonian form in Eq. (2.30) is preserved.

### 2.4 Average Fidelity

### 2.4.1 General Expression

In order to analyze the implications of the additional terms in the effective spin Hamiltonian (Eq. 2.30), one needs to take into account that the qubit is actually formed by two electron spins. It will be convenient to assume that the dots' velocities are small compared to the speed of sound,  $v_0 \ll s \sim 5 \times 10^3$  m/s, meaning the phonon effects should decouple from the dots' motion and can be approximated as contributing to the same "intrinsic" decoherence times as one would have without the motion. The effect of such decoherence terms on dynamical decoupling has been considered in detail in Ref. [91]; in the following we assume that these decoherence times are large compared to the characteristic period T of the dots' motion and therefore can be ignored.

In the absence of phonons, and approximating the Overhauser field as classical, the time evolution of the two-spin wavefunction with N = 4 components can be characterized by a unitary matrix U(t). The qubit subspace Q is formed by the  $m_z = 0$ component of the two-spin wavefunction; it has M = 2 dimensions. The standard assumption is that, at the beginning of the experiment, the two spins are initiated in a pure state  $|\psi\rangle$  which belongs to the qubit subspace,  $|\psi\rangle \in Q$ . Therefore, when computing the average fidelity, one only needs to average over the original wavefunctions in  $\mathcal{Q}$ .

More generally, consider an M-dimensional subspace  $\mathcal{Q}$  of an N-dimensional Hilbert space  $\mathcal{H}$ . Introducing an  $N \times M$  matrix T whose columns are formed by the components of orthonormal vectors that form a basis of  $\mathcal{Q}$ , then the components of an arbitrary wavefunction  $|\psi\rangle \in \mathcal{Q}$  are a linear combination of the columns of T; namely  $\psi = T\varphi$ , where  $\varphi$  is an M-dimensional column vector,  $\|\varphi\| = 1$ . The corresponding density matrix can be written in this basis as  $\rho_0 \equiv T\varphi \varphi^{\dagger}T^{\dagger}$ . The fidelity corresponding to the evolution matrix U is

$$F = \operatorname{Tr}(\rho_0 U \rho_0 U^{\dagger}) = (\varphi^{\dagger} W \varphi)(\varphi^{\dagger} W^{\dagger} \varphi), \qquad (2.33)$$

where  $W = T^{\dagger}UT$  can be thought of as the projection of U onto the subspace Q. The average fidelity in the subspace can now be calculated using the averaging identities for components  $\varphi_i$  of the normalized wavefunction  $|\varphi\rangle$ 

$$\langle \varphi_i \varphi_j^* \rangle = \delta_{ij} / M,$$
 (2.34)

$$\langle \varphi_i \varphi_j^* \varphi_k \varphi_l^* \rangle = \frac{\delta_{ij} \delta_{kl} + \delta_{il} \delta_{jk}}{M^2 + M}.$$
(2.35)

This leads to the average fidelity

$$\langle F \rangle = \frac{|\operatorname{Tr} W|^2 + \operatorname{Tr}(WW^{\dagger})}{M^2 + M}.$$
(2.36)

For the special case of the qubit formed by the singlet and m = 0 triplet states of two spins, assuming no interdot tunneling, the net evolution matrix is just the Kronecker product of evolution matrices corresponding to the two qubits,  $U = U_1 \otimes U_2$ . Further, it will be convenient to decompose each single-spin matrix in the interaction representation with respect to the precession in the net effective magnetic field along the z-axis,

$$U_i = U_{0i}S_i, \quad U_{0i} \equiv e^{-i\sigma^z \,\varphi_i(t)/2},$$
 (2.37)

where

$$\varphi_i(t) \equiv \omega_0 t + \int_0^t dt' \,\omega_{1i}^z(t'), \qquad (2.38)$$

 $\omega_0$  is the Larmor frequency,  $\omega_{1i}^z(t)$ , with i = 1, 2 [cf. Eq. (2.30)], are the two dot's effective perturbing fields in the z-direction, and the matrices

$$S_i \equiv e^{-i\gamma_i - i\phi_i \cdot \sigma/2}, \quad i = 1, 2, \tag{2.39}$$

are parametrized as rotations by angle  $\phi_i \equiv |\phi_i|$  around the unit vectors  $\hat{\phi}_i$ , with extra phases  $\gamma_i$ . These rotations come entirely from transverse,  $\mu = x, y$ , components of  $\omega_{1i}^{\mu}$ in the rotating frame, largely due to the Larmor frequency. Since the Larmor frequency is large, the additional rotation angles are expected to be small; we expand the average fidelity (2.36) to quadratic order in components of  $\phi_i$ ,

$$\langle F \rangle = 1 - f_0 - f_1 - f_2^z - f_2^\perp + \dots,$$
 (2.40)

with the infidelity terms

$$f_0 = \frac{2}{3}\sin^2(\Delta\varphi/2),$$
 (2.41)

$$f_1 = \frac{1}{3}(\phi_2^z - \phi_1^z)\sin(\Delta\varphi), \qquad (2.42)$$

$$f_2^z = \frac{1}{6} (\phi_2^z - \phi_1^z)^2 \cos(\Delta \varphi), \qquad (2.43)$$

$$f_2^{\perp} = \frac{2 + \cos(\Delta\varphi)}{12} \left[ (\phi_1^{\perp})^2 + (\phi_2^{\perp})^2 \right].$$
 (2.44)

Note that, as expected, the fidelity only depends on the differences  $\Delta \varphi \equiv \varphi_2 - \varphi_1$  and  $\phi_2^z - \phi_1^z$  of the two precession angles around the z-axis. To the same quadratic accuracy in the small angles, one can also write

$$f_0 + f_1 + f_2^z = \frac{1}{3} [1 - \cos(\Delta \varphi + \phi_2^z - \phi_1^z)], \qquad (2.45)$$

which only depends on the total phase difference, and is exact in the limiting case when  $\phi_i^{\perp} = 0.$ 

### 2.4.2 Rotating-Frame Approximation

Returning to the analysis of the single-spin Hamiltonian (Eq. 2.30), the Larmor frequency,  $\omega_0 \gtrsim 4 \times 10^9 \text{ rad/s}$ , is the dominant term,  $\omega_1 \ll \omega_0$ . Given that the phonons have been excluded, one can also assume that the dot trajectory is such that the time dependence in  $\omega_1(t)$  is slow on the scale of  $\omega_0$ . This implies that the interaction representation Eq. (2.37) is valid, where S(t) is the slow part of the unitary evolution operator; it obeys the equation

$$i\hbar \dot{S} = H_{\rm int}(t)S, \quad S(0) = 1,$$
(2.46)

where the dot index is temporarily omitted, and

$$\frac{1}{\hbar}H_{\rm int}(t) \equiv \frac{1}{2}\omega_1(t) \cdot U_0^{\dagger}(t)\boldsymbol{\sigma}U_0(t) \qquad (2.47)$$

$$= \frac{1}{2} \left\{ \sigma^x [\omega_1^x(t)\cos\varphi(t) + \omega_1^y(t)\sin\varphi(t)] + \sigma^y [\omega_1^y(t)\cos\varphi(t) - \omega_1^x(t)\sin\varphi(t)] \right\} \qquad (2.48)$$

is the perturbing Hamiltonian in the interaction representation.

The remaining analysis is performed perturbatively with the Magnus (cumulant) expansion,

$$S(t) = \exp(\mathcal{C}^{(1)} + \mathcal{C}^{(2)} + \ldots), \qquad (2.49)$$

where the first two cumulants are

$$\mathcal{C}^{(1)}(t) = -\frac{i}{\hbar} \int_0^t dt_1 H_{\text{int}}(t_1), \qquad (2.50)$$

$$\mathcal{C}^{(2)}(t) = -\frac{1}{\hbar^2} \int_{0 < t_1 < t_2 < t} dt_1 dt_2 [H_{\text{int}}(t_2), H_{\text{int}}(t_1)].$$
(2.51)

The integration is performed explicitly to leading order in  $\omega_1/\omega_0$ , also assuming that the time-dependence of  $\omega_1$  is slow on the scale of  $\omega_0$ :

$$i\mathcal{C}^{(1)}(t) = \frac{1}{2\omega_0} \Big\{ \sigma^x [a\sin\varphi(t) + b_0 - b\cos\varphi(t)] \\ + \sigma^y [b\sin\varphi(t) - a_0 + a\cos\varphi(t)] \Big\},$$
(2.52)

$$i\mathcal{C}^{(2)}(t) = \frac{1}{4\omega_0}\sigma^z \int_0^t dt_1 \left[a^2(t_1) + b^2(t_1)\right], \qquad (2.53)$$

denoting  $a \equiv \omega_1^x(t)$ ,  $b \equiv \omega_1^y(t)$ , and  $a_0$ ,  $b_0$  as the corresponding values at t = 0 [ $\varphi(0) \equiv 0$  by definition]. Eq. 2.40 gives the expression for the average qubit infidelity

$$1 - \langle F \rangle = f_z + f_1^{\perp} + f_2^{\perp}, \qquad (2.54)$$

$$f^{z} = \frac{1}{3} [1 - \cos(\Delta \varphi + \delta_{2} - \delta_{1})],$$
 (2.55)

$$f_i^{\perp} = \frac{2 + \cos(\Delta\varphi)}{12\omega_0^2} \left[ (B_i - b_{0i})^2 + (A_i - a_{0i})^2 \right], \qquad (2.56)$$

where the index i = 1, 2 refers to the two spins,  $A_i \equiv a_i \cos \varphi(t) + b_i \sin \varphi(t)$  and  $B_i \equiv b_i \cos \varphi(t) - a_i \sin \varphi(t)$  are the rotated components of the transverse angularvelocity vectors,  $(\omega_1^x, \omega_1^y)$ , for the corresponding spins, and the additional phases  $\delta_i$  are given by the integrals.

$$\delta_i \equiv \int_0^t \frac{dt'}{2\omega_0} \, [\omega_{1i}^{\perp}(t')]^2. \tag{2.57}$$

One immediately recognizes the additional phases in Eqs. (2.55) and (2.57) as the effect of level repulsion, or equivalently as the gap for the spins, driven adiabatically by the total instantaneous magnetic field  $\propto [(\omega_0 + \omega_{1i}^z)^2 + (\omega_{1i}^{\perp})^2]^{1/2}$ . Then, Eq. (2.56) can be interpreted as the effect of the basis change between the original quantization z-axis and the direction of the instantaneous magnetic field; it is similar in nature to the initial decoherence associated with any periodic decoupling sequence, see, e.g. Ref. [92].

The terms of the next order in  $1/\omega_0$  expansion, omitted in Eqs. (2.52) and (2.53), include the trivial correction to Eq. (2.53)  $\propto \int dt \,\omega_1^z [\omega_1^{\perp}]^2/\omega_0^2$ , as well as the ge-

ometrical phase  $\propto \int dt' W[\omega_1^x, \omega_1^y]$ , where  $W[x, y] \equiv W[x(t), y(t)] \equiv x(t)y'(t) - x'(t)y(t)$ is the Wronskian.

Since  $\Delta \varphi = \int_0^t dt' [\omega_{12}^z(t') - \omega_{11}^z(t')]$ , the term  $f^z$  does not contain any rapid oscillations at the Larmor precession frequency  $\omega_0$ , while the terms  $f_i^{\perp}$  can be averaged over the period of Larmor precession by replacing  $A_i^2 + B_i^2$  with  $(\omega_{1i}^{\perp})^2$  and dropping all of the terms linear in  $A_i$ ,  $B_i$ ,

$$\overline{f}_{i}^{\perp} = \frac{2 + \cos(\phi_{2}^{z} - \phi_{1}^{z})}{3\omega_{0}^{2}} \left\{ [\omega_{1i}^{\perp}(t)]^{2} + [\omega_{1i}^{\perp}(0)]^{2} \right\}.$$
(2.58)

### 2.4.3 Sequences

Overall, the dynamical decoupling should be designed to null the difference between the accumulated phases  $\varphi_i + \delta_i$  of the two spins which suppresses the main contribution to the infidelity, see Eq. (2.55). For a static Overhauser field, this can be achieved just by ensuring that each spin spends the same amount of time at each position, e.g., via the solid adiabatic trajectory in Fig. [2.5]. This is also sufficient to suppress the effect of the velocity-dependent terms in the second line of Eq. (2.32). A more complicated set of dot rotation involving motion in both directions, e.g., see the dashed trajectory of Fig. [2.5], are required to suppress a time-dependent Overhauser field.

To analyze the effect of a sequence of  $\pi$ -rotations in the presence of a time and position-dependent Overhauser field, consider its Fourier expansion at a position on the trajectory parametrized by the rotation angle  $\theta$ ,

$$B^{z}(\theta, t) = A_{0}(t) + \sum_{m} A_{m}(t) \cos m\theta + B_{m}(t) \sin m\theta.$$
(2.59)



Figure 2.5: Time dependent adiabatic trajectories used in the simulations. Plotted is the position of the first dot parametrized in terms of the angle  $\theta$  as a function of time t. The angle-dependent positions were defined as a sum of properly scaled and shifted hyperbolic tangents. Single direction rotations suppress the effect of a static Overhauser field but not of a time-varying one. Longer rotation sequences like alternating forwardand-back suppress the effect of a linear in time Overhauser field. Arbitary rotations on the Bloch sphere may be accomplished using additional operations when the dot is stationary, corresponding to the flat segments in the sequence.

Only the difference between the fields corresponding to the two dots (located at  $\theta$  and  $\theta + \pi$ ) is relevant for the infidelity Eq. (2.55). This leaves only the terms with m odd in the Fourier expansion (2.59).

For a term with  $\cos m\theta$  (an even function of  $\theta$ ), a sequence of  $\pi$  rotations acts the same way, independent of the direction. It is easy to check that with an equidistant sequence of rotations centered at  $T_0/2$ ,  $3T_0/2$ , ...,  $(2s + 1)T_0/2$ , where the number of rotations s is even; any time-independent and linear in t contributions to  $A_m(t)$  are suppressed, but a quadratic term would generally remain. Unlike the usual dynamical decoupling problem[60, 116], it is not generally possible to suppress the quadratic term of  $A_m(t)$ .

The rotation direction starts to matter for a term with  $\sin m\theta$  which is an odd function of  $\theta$ . Here a sequence of  $\pi$  rotations in the same direction picks up a sum of contributions from consecutive time intervals with alternating signs, suppressing the time-independent contribution to  $B_m(t)$  but not the linear contribution. As an alternative prescription, a symmetrized sequence of two forward rotations by angle  $\pi$ , followed by two rotations in the opposite direction can be used to suppress the effect of the linear in t term in  $B_m(t)$ . Generally, it is possible to design more complicated sequences analogous to concatenated or Uhrig's sequences to suppress the effect of any fixed-degree polynomial in time  $B_m(t)$ [60, 116]. However, this is not expected to be useful since the quadratic time contribution of  $A_m(t)$  would still remain.

### 2.5 Simulations

The above analytical results were corroborated by simulating the two-spin unitary evolution with the effective Hamiltonian (Eq. 2.30). Specifically, the dot trajectory was parametrized by the rotation angle  $\theta = \theta_1(t)$ ; the other dot is assumed to have the symmetric position,  $\theta_2(t) = \theta_1(t) + \pi$  [see Fig. 2.5 for samples of actual trajectories.] The position-dependent terms in the first line of Eq. 2.32 were simulated in terms of a three-component correlated magnetic field  $\mathbf{B}(\theta)$  drawn from the Gaussian distribution with zero average and the correlation function  $\langle B_{\mu}(\theta)B_{\nu}(\theta')\rangle = \sigma_{\mu}^2 \delta_{\mu\nu}\vartheta(\theta - \theta')$  (no implicit summation in  $\mu, \nu = x, y, z$ ), where

$$\vartheta(\theta) \equiv \sum_{m=-\infty}^{\infty} e^{-(\theta - 2\pi m)^2/\ell^2}$$

is an infinite sum of Gaussian functions (which can also be represented in terms of the Jacobi theta function). These were obtained by applying a Gaussian filter to a discrete set of uncorrelated random numbers drawn from the Gaussian distribution, and using the standard cubic spline interpolation with the result. To simulate the components of the time-dependent magnetic field  $B(\mathbf{r}, t)$ , explicit order-r spline interpolation are used between several such angle-dependent functions, where r = 1, 2.

For all simulations, the time units were chosen to correspond to the Larmor precession period,  $\tau_0 \equiv 2\pi/\omega_0$ , and the correlation time of the Overhauser field,  $4 \cdot 10^4 \tau_0$ , with each component of its r.m.s. value corresponding to rotation frequency,  $\langle |\omega_1|^2 \rangle = 0.025/\tau_0^2$ . The adiabatic trajectories of the dots were simulated using a sum of appropriately shifted hyperbolic tangents, scaled so that the dot is in motion for approximately half of the protocol. The leading velocity-dependent term in Eq. 2.32 was simulated using the corresponding derivatives and the parameter  $Q_v/\hbar = 0.075/\tau_0$ , assuming equal contributions from the Rashba and Dresselhaus parameters.

For the case of the static Overhauser field (see Fig. 2.6) the average infidelity is dominated by the contribution from the z-component of the field, and the



Figure 2.6: Simulated qubit infidelity  $1 - \langle F \rangle$  (Eq. 2.36) in the vicinity of the first full rotation period of the double-dot qubit at t = T. Position-dependent magnetic field  $B_{\mu}(\theta)$  is assumed static, and the rotation period T is chosen commensurate with the Larmor frequency,  $\omega_0 T/2\pi = 400$ . Dashed line: only  $B_z(\theta)$  is included; the infidelity minimum is exactly at t = T, in agreement with Eq. 2.56 which is exact in this situation. Dotted line: only the transverse components  $B_{\mu}(\theta)$ ,  $\mu = x, y$  are included. The infidelity minimum is slightly off the commensurate time t = T due to the terms not included in Eq. 2.54. All three components of the field  $B_{\mu}(\theta)$  are included for the red solid line.



Figure 2.7: Simulation results for fidelity measured at the end of each cycle, for a linear time-dependence of  $B_{1i}^{\mu}$  with (a) forward, (b) forward-back dot rotations (cf. Fig. 2.5), as well as (c) forward-back for quadratic time-dependence.

infidelity nearly vanishes at the end of the spatial exchange protocol. For a linearly time-interpolated Overhauser field, the infidelity increases over several cycles of the single-direction  $\pi$  pulses(see Fig. 2.7a), but maintains a low value  $\sim 10^{-5}$  after alternating between a sequence of two forward rotations, followed by two rotations in the opposite direction (see Fig. 2.7b). However, for the quadratic interpolation (see Fig. 2.7c) the infidelity gradually increases even for the alternating protocol, though it stays below  $10^{-4}$  for several cycles. The code for the above simulation is provided in Appendix A.

## 2.6 Possible Experimental Setup

Aspects of this proposal, such as the precise construction of few-electron quantum dots in III-V semiconductor heterostructures, have already been demonstrated in experiments [87, 97, 117, 65]. However, the precise adiabatic rotation required in our proposal may be quite difficult to accomplish experimentally. This section discusses possible design implementations for an experimental realization, as well as physical constraints.

As discussed above, our proposal is suitable in materials with relatively weak spin-orbit coupling such as GaAs/AlGaAs heterostructures. It should be possible for the electrons in a 2DEG to be confined in the radial direction by creating a circular depletion layer by placing electrostatic gates with a constant voltage in the center and outer edge of the circle as sketched in Figure 2.4. Since the tunneling should be suppressed, the normal interdot spacing should be much larger than the dot size, a, so the circular trajectory can have a radius  $r_0 \sim 15a \sim 1 \ \mu\text{m}$ . Confinement and rotation in the angular direction could be accomplished by placing appropriately chosen time-dependent voltages on the "wedge-gates" on both sides of the electron (Fig. 2.4). Several of these wedges will be needed to accomplish the smooth and adiabatic trajectory needed. The combined use of wedge and circular gates may require gates on both the top and bottom of the sample. The typical confining potential is approximated by  $U \sim m\omega_d^2 a^2$ , which only requires reasonable gate voltages on the order of 100 mV.

The averaging of the hyperfine interaction is only valid in the quasi-static approximation of the Overhauser field. In general, the hyperfine interaction between the electron and the nuclei leads to a Knight shift. However, in the presence of the magnetic field, fluctuating corrections to the quasi-static approximation are inversely proportional to  $B_0$  and can be neglected [114]. Thus, the most significant effect in this case is due to the dipole-dipole interaction of nearest neighboring nuclei which requires that  $T \ll t_{\rm nuc} \sim 10^{-4}$ s [77]. This places a lower bound on the velocity, while there is also an upper bound necessary to ensure that the Lorentz force from the dot rotation only deforms the actual path of the electron by a distance much smaller than the correlation length of the Overhauser field. This results in the restriction,  $10^{-1} \ll v \ll 10^5$  in m/s. This rather lenient constraint is due to the assumption that the confinement potential be large compared to the other potentials in our system. This allows one to neglect trajectory deviations from perturbations such as charge noise. In these estimates, v is assumed to be small compared to the speed of sound,  $s \sim 5 \times 10^3$  m/s. With  $a/r_0 \sim 15$ and  $v \sim 10$  m/s, a rotation period of  $T \sim 1 \ \mu$ s, which easily satisfies the above conditions and, according to the simulations, should result in the infidelities lower than  $10^{-4}$  for significant time-scales. The rotation period could potentially be decreased to allow more operations to be performed before the states decohere.

# 2.7 Quantum Dot Proposal Summary

The real-space exchange of quantum dots was analyzed as a possible substitute for the tunneling exchange. Ideally, exchange eliminates the hyperfine dephasing from the Overhauser field parallel to the applied field, leaving only the smaller effects from the in-plane field. The real-space exchange accomplishes the same suppression of the hyperfine interaction, but avoids the problematic sensitivity to charge noise present in exchange via tunneling. While spatial exchange does introduce additional effects such as spin-orbit coupling, simple tricks like using pairs of  $\pi$  rotations in alternating directions can be used to suppress these so that the decoherence is still dominated by the hyperfine interaction. In particular, this field only enters as the small ratio of the average in-plane Overhauser field to the externally applied field. Perhaps the simplest way to suppress the hyperfine interaction in this approach is to reduce this ratio by increasing the externally applied field. In addition, this spatial exchange is compatible with some of the methods already being attempted such as nuclear polarization via pumping. If the hyperfine interaction can be further suppressed, the next largest contribution from our spatial exchange approach would come from the disorder of the sample or the electron-phonon coupling. The analysis of this spatial exchange also remains valid in systems that use additional quantum dots [65, 117], with universal quantum operations in mind, as long as each operation is applied to only two dots at a time.

While the movement of quantum dots requires the precise control of the confining potential, which may be difficult to realize experimentally, the analysis shows that the construction of such a system is viable. With realistic parameter values from current experiments, the analysis produces infidelities smaller than  $10^{-4}$  after ten decoupling cycles. This setup could also be a productive step towards the experimental realization of more complicated exchange systems, with many more interesting applications.

# Chapter 3

# Majorana Bound States Background

# 3.1 Historical Review: Topological Quantum Information

### 3.1.1 Majorana Fermions

The experimental observation of a self-conjugate fermionic particle with equal creation and annihilation operators,  $\hat{\gamma}^{\dagger} = \hat{\gamma}$ , known as a Majorana fermion, has been a goal in physics since it was first theorized by Ettore Majorana over 75 years ago as a real solution to the relativistic Dirac equation [74]. It should be noted that the creation  $(\hat{c}^{\dagger})$ and annihilation  $(\hat{c})$  operators of any conventional fermion can be trivialy decomposed into two self-conjugate forms

$$\hat{\gamma}^A = -i(\hat{c} - \hat{c}^{\dagger}) \tag{3.1}$$

$$\hat{\gamma}^B = \hat{c} + \hat{c}^{\dagger}. \tag{3.2}$$

While these operators are mathematically valid, they do not neccesarily correspond to physically meaningful particles that occur in nature. While the existence of a self-conjugate, elementary particle is still an unsolved problem [98, 58, 115], there has been recent evidence, and subsequent interest, that a selfconjugate quasiparticle may also exist in condensed matter systems [37]. Besides being self-conjugate, these particles are novel because they arise from topological principles rather than broken symmetries alone. The rest of this section briefly reviews some of these systems; refer to the articles by Nayak et. al. and Hasan et. al. [83, 48] for a more thorough review.

### 3.1.2 Two-Dimensional Spinless p-wave Superconductor

One of the first systems theorized to host a self-conjugate fermion is a twodimensional p-wave superconductor with spinless, or spin-polarized, particles as noted by Read and Green [95], given by the effective Hamiltonian

$$H = \sum_{k} \left[ t_k \hat{c}_k^{\dagger} \hat{c}_k + \frac{1}{2} \left( \Delta_k^* \hat{c}_{-k} \hat{c}_k + \Delta_k \hat{c}_k^{\dagger} \hat{c}_{-k}^{\dagger} \right) \right],$$
(3.3)

where  $t_k = \epsilon_k - \mu$ , with single-particle kinetic energy  $\epsilon_k$ , and  $\Delta_k$  has p-wave symmetry (e.g.  $\Delta_k \propto k_x - ik_y$ ). This system can be diagonalized by making a Bogoliubov-de Gennes transformation

$$\hat{a}_k = u_k \hat{c}_k - v_k \hat{c}_{-k}^{\dagger} \tag{3.4}$$

$$\hat{a}_{k}^{\dagger} = u_{k}^{*} \hat{c}_{k}^{\dagger} - v_{k}^{*} \hat{c}_{-k}$$
(3.5)

with the requirement that  $|u_k|^2 + |v_k|^2 = 1$ . The Bogoliubov coefficients  $u_k$  and  $v_k$  can be found by requiring that  $[\hat{a}_k, H] = E_k \hat{a}_k$ , which is consistent with

$$H = E_0 + \sum_k E_k \hat{a}_k^{\dagger} \hat{a}_k, \qquad (3.6)$$

for some constant  $E_0$  and positive quasiparticle energies,  $E_k$ . Together with the transformation, the commutation relation leads to the Bogoliubov-de Gennes(BdG) equations

$$t_k u_k - \Delta_k^* v_k = E_k u_k, aga{3.7}$$

$$-\Delta_k u_k - t_k^* v_k = E_k v_k, (3.8)$$

or equivalently in matrix form,

$$\begin{pmatrix} t_k & -\Delta_k^* \\ -\Delta_k & -t_k \end{pmatrix} \begin{pmatrix} u_k \\ v_k \end{pmatrix} = E_k \begin{pmatrix} u_k \\ v_k \end{pmatrix},$$
(3.9)

which lead to the equations

$$E_k^2 = t_k^2 + |\Delta_k|^2, (3.10)$$

$$\frac{v_k}{u_k} = \frac{t_k - E_k}{\Delta_k^*}, \tag{3.11}$$

$$|u_k|^2 = \frac{1}{2} \left( 1 + \frac{t_k}{E_k} \right), \qquad (3.12)$$

$$|v_k|^2 = \frac{1}{2} \left( 1 - \frac{t_k}{E_k} \right).$$
 (3.13)

The topological nature of this system can be seen by noting the asymptotic behavior of k. For large values of |k|,  $\Delta_k \simeq 0$  and  $\epsilon_k \gg \mu$ , which implies  $t_k \simeq E_k$ . Thus as  $|k| \to \infty$ ,  $|u_k| \to 1$  and  $|v_k| \to 0$  [see Eq.(3.13)]. For small k,  $\epsilon_k \simeq k^2/(2m^*)$  and  $\Delta_k \simeq \Delta(k_x - ik_y)$ , so as  $k \to 0$  then  $t_k \to -\mu$  and  $\Delta_k \to 0$ , which together lead to  $E_k^2 \to t_k^2$ , yielding two distinct cases.

### 3.1.2.1 Topological Phases

When  $\mu < 0$ ,  $t_k > 0$  and  $t_k \to E_k$  implies  $|u_k| \to 1$  and  $|v_k| \to 0$ , while when  $\mu > 0$ ,  $t_k < 0$  and  $t_k \to -E_k$  implies  $|u_k| \to 0$  and  $|v_k| \to 1$ . These two cases have physically distinct phases: the first case, known as strong-pairing, has shortrange wavefunctions, while the second case, known as weak-pairing, has long-range wavefunctions [95]. More so, these two phases are topologically distinct, which can be seen by looking at the relationship between k and the Bogoliubov coefficients  $u_k$  and  $v_k$ .

Specifically, values of the Bogoliubov coefficients  $u_k$  and  $v_k$ , which satisfy  $|u_k|^2 + |v_k|^2 = 1$ , correspond to points on the surface of a sphere,  $S^2$ , with the point  $|u_k| = 1$ ,  $|v_k| = 0$  defined as the North pole, while the opposite case,  $|u_k| = 0$ ,  $|v_k| = 1$ , is defined as the South pole. By adding the point at  $|k| \to \infty$  as the North pole, the two-dimensional plane spanned by  $k_x$  and  $k_y$  corresponds topologically to  $S^2$  as well. This defines a map from the k-space  $S^2$  to u,v-space  $S^2$  that forms the homotopy group  $\pi_2(S^2)$ , which is equivalent to the group of integers Z. This mapping has an integer topological invariant, known as the degree, that counts the minimum number of times a mapping must pass through any point other than the North pole [118, 95].

In the  $\mu < 0$  case, both  $|k| \to \infty$  and  $k \to 0$  can be mapped to the point with  $|u_k| = 1$ ,  $|v_k| = 0$  at the North pole, meaning that this mapping can be smoothly deformed to the trivial mapping of all points to the North pole. Thus, the mapping does not need to pass through any point other than the North pole, and the degree of this phase is 0. For the  $\mu > 0$  case, on the other hand, since  $k \to 0$  corresponds to  $|u_k| = 0$ ,  $|v_k| = 1$ , this mapping must include the South pole and cannot be smoothly deformed to the trivial mapping. Thus, this phase has degree 1 and these two phases are topologically distinct; the strong and weak pairing phases cannot be smoothly deformed into each other without changing the sign of  $\mu$ .

#### 3.1.2.2 Majorana Bound States and Non-Abelian Statistics

Since these two phases are topologically distinct, any system that includes both phases must contain a boundary where the topological parameter changes sign, i.e.,  $\mu = 0$ . At this boundary value, the k = 0 solution has  $t_0 = -\mu = 0$ , meaning  $E_0 = 0$  since  $\Delta_0 = 0$ . Solving the BdG equations when  $\mu = 0$  also yields  $u_k = v_k^*$ , meaning the k = 0 quasiparticle operator satisfies  $\hat{a}_0^{\dagger} = \hat{a}_0$  [see Eq. (3.5)] and is thus a self-conjugate Majorana fermion. As a consequence, a system in the topologically non-trivial weak-pairing phase that contains a defect in the trivial phase, such as the core of a superconducting vortex, must have a zero-energy Majorana state bound to that defect.

While the guarantee of Majorana bound states (MBS) pinned to defects is interesting in itself, the real promise of these states come from their non-trivial exchange statistics. Unlike the wavefunctions of conventional bosons or fermions that only gain a phase of 1 or -1 when they are exchanged, these topological bound states can take on "any" phase upon exchange, and are known as "anyons" for this reason [95, 83].

Just as any conventional fermion can be decomposed into two self-conjugate fermions [see Eq. (3.2)], a pair of Majorana bound states  $\hat{\gamma}_A$  and  $\hat{\gamma}_B$ , bound to two different defects, can be combined to form a conventional, albeit non-local, fermion by writing

$$\hat{d} = \frac{1}{2}(\hat{\gamma}_A + i\hat{\gamma}_B).$$
 (3.14)

Since the MBS has zero-energy, the presence or absence of this fermion does not change the energy, meaning there are two degenerate ground states. More generally, 2n Majorana bound states form n conventional fermions, which can each be present or absent, forming a  $2^n$ -fold ground state degeneracy. Physically exchanging any two of the bound states, known as braiding, applies a non-trivial unitary transformation to the degenerate manifold of ground states and can be used to perform quantum operations. Furthermore, when exchanging multiple pairs of Majorana bound states, the order of the braiding matters; the braiding operations do not commute, also known as "non-Abelian" in the group theory context [83]. Since the MBS bound to topological defects are inherently non-local, they are immune to local pertubations, and thus, excellent candidates for fault-tolerant qubits.

### 3.1.3 Additional Topological Systems

Many other physical systems have been proposed to advance towards a more feasible experimental setup, though the underlying topological arguments often remain the same. For example, the non-Abelian statistics of Majorana bound states, also known as Ising anyons, were also discussed in the Moore-Read Pfaffian state of the  $\nu = 5/2$ fractional quantum Hall state [81]. In addition, it is worth noting that the  $\nu = 12/5$ Read-Reyazi state [96] may host more powerful topological states, known as Fibonacci anyons, which are capable of universal quantum computation [38, 83], unlike Ising anyons which require entangled states via magic state distillation [11].

Ivanov showed that equivalent MBS exist in the cores of half-quantum vortices (vortices where the superconducting phase shifts by  $\pi$ , rather than  $2\pi$ , as it circles the core) of *p*-wave superconductors with spin [54], rather than the spinless or spinpolarized case discussed previously. Similarly, Kane and Mele found that similar states should exist in quantum spin Hall states of graphene [56].

Another important step came when Fu and Kane found that the required superconducting Cooper pairs with the rare p-wave symmetry could be replaced by Cooper pairs that tunnel from a conventional s-wave superconductor into a topological insulator via the proximity effect [39]. This proposal was advanced even further by Sau et. al. when they replaced the topological insulator with a ferromagnetic insulator and semiconductor with strong Rashba spin-orbit coupling [102]. Alicea simplified that proposal even more by showing that the ferromagnetic insulator could be omitted if the semiconductor also had Dresselhaus spin-orbit coupling and a magnetic field was applied [2].

All of these gradual improvements culminated in the proposals by Lutchyn et. al. [73] and Oreg et al. [86] of a one-dimensional semiconductor nano-wire with strong spin-orbit coupling, on top of an *s*-wave superconductor in the presence of a magnetic field. Since it only involves relatively conventional materials and a onedimensional system, their proposal was experimentally attempted by several groups. These groups reported zero-bias conductance peaks when connecting the nano-wire to a normal lead [82, 24, 20, 36, 18], indicating a single level at zero-energy that is consistent with a MBS at the wire-ends. This exciting development has spurred a great amount of interest in this system, and is the subject of the proposal in Ch. 4.

The remainder of this chapter will introduce and discuss some details of this "Majorana wire" system. However, it is helpful to first understand a simplified model analyzed by Kitaev [61], as well as Lieb et. al. in the context of the transverse Ising model [70].

# 3.2 Review of Kitaev's Toy Model

Kitaev's toy model is a 1D chain of N spinless electrons, with tight-binding hopping parameter  $t_0$ , chemical potential  $\mu$ , and superconducting pairing  $\Delta e^{i\theta}$  with phase  $\theta$  on bonds (i.e. p-wave).

$$H = \mu \sum_{1 \le j \le N} \left( c_j^{\dagger} c_j - \frac{1}{2} \right) + \sum_{1 \le j \le N-1} \left[ -t_0 (c_j^{\dagger} c_{j+1} + c_{j+1}^{\dagger} c_j) + \Delta e^{i\theta} c_j c_{j+1} + \Delta e^{-i\theta} c_{j+1}^{\dagger} c_j^{\dagger} \right]$$
(3.15)

This Hamiltonian is not meant to represent a physical system, but is useful in easily demonstrating the existence of localized unpaired Majorana modes at the ends of the wire. With that said, one can think of a spinless model as a real system where one spin species is energetically favored, effectively eliminating spin as a true degree of freedom in the low-energy Hamiltonian. Indeed, the semiconductor wire system discussed in the next section can be directly mapped to the Kitaev model in certain parameter regimes.

Consider the self-conjugate Majorana operators (i.e. those that satisfy  $\gamma^{\dagger} = \gamma$ )

$$\gamma_{j}^{A} = -i(e^{i\frac{\theta}{2}}c_{j} - e^{-i\frac{\theta}{2}}c_{j}^{\dagger}) \qquad \gamma_{j}^{B} = e^{i\frac{\theta}{2}}c_{j} + e^{-i\frac{\theta}{2}}c_{j}^{\dagger}.$$
(3.16)

. One can easily check the anti-commutation relations

$$\{\gamma_j^A, \gamma_{j'}^A\} = -(\{c_j, -c_{j'}^\dagger\} + \{-c_j^\dagger, c_{j'}\}) = 2\delta_{jj'}$$
(3.17)

$$\{\gamma_j^B, \gamma_{j'}^B\} = \{c_j, c_{j'}^{\dagger}\} + \{c_j^{\dagger}, c_{j'}\} = 2\delta_{jj'}$$
(3.18)

$$\{\gamma_j^A, \gamma_{j'}^B\} = -i(\{c_j, c_{j'}^\dagger\} + \{-c_j^\dagger, c_{j'}\}) = 0$$
(3.19)

to give the general  $\{\gamma_j^{\alpha}, \gamma_{j'}^{\alpha'}\} = 2\delta_{jj'}\delta_{\alpha\alpha'}$ . Note that there are two "species" of Majorana fermions that commute with each other, but which one is A or B is an arbitrary distinction. Equivalently, one can solve for the electron operators

$$c_{j} = \frac{1}{2}e^{-i\frac{\theta}{2}}(\gamma_{j}^{B} + i\gamma_{j}^{A}) \qquad c_{j}^{\dagger} = \frac{1}{2}e^{i\frac{\theta}{2}}(\gamma_{j}^{B} - i\gamma_{j}^{A})$$
(3.20)

and substitute these into the Hamiltonian. Before doing this, note the useful property

$$\gamma_j^{\alpha}\gamma_j^{\alpha} = \frac{1}{2}(\gamma_j^{\alpha}\gamma_j^{\alpha} + \gamma_j^{\alpha}\gamma_j^{\alpha}) = \frac{1}{2}\{\gamma_j^{\alpha}, \gamma_j^{\alpha}\} = 1.$$
(3.21)

Thus the Hamiltonian in terms of the Majorana operators is given by

$$\begin{split} H &= \frac{\mu}{4} \sum_{1 \leq j \leq N} \left[ (\gamma_j^B - i\gamma_j^A)(\gamma_j^B + i\gamma_j^A) - 2 \right] \\ &+ \frac{1}{4} \sum_{1 \leq j \leq N-1} \left\{ -t_0 \Big[ (\gamma_j^B - i\gamma_j^A)(\gamma_{j+1}^B + i\gamma_{j+1}^A) + (\gamma_{j+1}^B - i\gamma_{j+1}^A)(\gamma_j^B + i\gamma_j^A) \Big] \right. \\ &+ \Delta \Big[ (\gamma_j^B + i\gamma_j^A)(\gamma_{j+1}^B + i\gamma_{j+1}^A) + (\gamma_{j+1}^B - i\gamma_{j+1}^A)(\gamma_j^B - i\gamma_j^A) \Big] \Big\} \\ &= \frac{\mu}{4} \sum_{1 \leq j \leq N} \left[ 1 + 1 + i(\gamma_j^B \gamma_j^A - \gamma_j^A \gamma_j^B) - 2 \right] \\ &+ \frac{1}{4} \sum_{1 \leq j \leq N-1} \left\{ -t_0 \Big[ \gamma_j^B \gamma_{j+1}^B + \gamma_j^A \gamma_{j+1}^A + i(\gamma_j^B \gamma_{j+1}^A - \gamma_j^A \gamma_{j+1}^B) \Big] \right. \\ &- t_0 \Big[ \gamma_{j+1}^B \gamma_j^B + \gamma_{j+1}^A \gamma_j^A + i(\gamma_j^B \gamma_{j+1}^A - \gamma_j^A \gamma_{j+1}^B) \Big] \\ &+ \Delta \Big[ \gamma_j^B \gamma_{j+1}^B - \gamma_j^A \gamma_{j+1}^A + i(\gamma_j^B \gamma_{j+1}^A + \gamma_j^A \gamma_{j+1}^B) \Big] \\ &+ \Delta \Big[ \gamma_j^B \gamma_j^B - \gamma_{j+1}^A \gamma_j^A - i(\gamma_{j+1}^B \gamma_j^A + \gamma_{j+1}^A \gamma_j^B) \Big] \Big\} \\ &= \frac{i}{2} \Big\{ -\mu \sum_{1 \leq j \leq N} \gamma_j^A \gamma_j^B + \sum_{1 \leq j \leq N-1} \Big[ (t_0 + \Delta) \gamma_j^A \gamma_{j+1}^B + (t_0 - \Delta) \gamma_{j+1}^A \gamma_j^B \Big] \Big\}. \end{split}$$

It is useful to analyze this Hamiltonian in two extreme parameters regimes. First, if  $t_0 = \Delta = 0$  while  $\mu \neq 0$  we have

$$H = -\frac{i}{2}\mu \sum_{1 \le j \le N} \gamma_j^A \gamma_j^B = -\frac{1}{2}\mu \sum_{1 \le j \le N} (e^{i\frac{\theta}{2}}c_j - e^{-i\frac{\theta}{2}}c_j^{\dagger})(e^{i\frac{\theta}{2}}c_j + e^{-i\frac{\theta}{2}}c_j^{\dagger}) \quad (3.22)$$

$$= \mu \sum_{1 \le j \le N} c_j^{\dagger} c_j, \qquad (3.23)$$

which is, not surprisingly, just a trivial 1D chain of fermions on the usual sites, with a ground state with no fermions present. This is a reminder that the substitution of Majorana operators is always mathematically possible, but in most cases doesn't lead to any physical significance. However, if  $\mu = 0$  and  $\Delta = t_0$  then

$$H = it_0 \sum_{1 \le j \le N-1} \gamma_j^A \gamma_{j+1}^B, \qquad (3.24)$$

which importantly lacks the operators  $\gamma_N^A$  and  $\gamma_1^B$ . One can introduce new fermion operators that are localized between two adjacent sites

$$d_j = \frac{1}{2} (\gamma_j^A + i \gamma_{j+1}^B), \qquad (3.25)$$

which are conventional in the sense that they are not self-conjugate and they satisfy the usual anti-commutation relations

$$\{d_j, d_{j'}\} = \frac{1}{4}\{\gamma_j^A + i\gamma_{j+1}^B, \gamma_{j'}^A + i\gamma_{j'+1}^B\} = \frac{1}{2}(\delta_{jj'} - \delta_{jj'}) = 0, \qquad (3.26)$$

$$\{d_{j}^{\dagger}, d_{j'}^{\dagger}\} = \frac{1}{4}\{\gamma_{j}^{A} - i\gamma_{j+1}^{B}, \gamma_{j'}^{A} - i\gamma_{j'+1}^{B}\} = \frac{1}{2}(\delta_{jj'} - \delta_{jj'}) = 0, \qquad (3.27)$$

$$\{d_{j}^{\dagger}, d_{j'}\} = \frac{1}{4}\{\gamma_{j}^{A} - i\gamma_{j+1}^{B}, \gamma_{j'}^{A} + i\gamma_{j'+1}^{B}\} = \frac{1}{2}(\delta_{jj'} + \delta_{jj'}) = \delta_{jj'}.$$
 (3.28)

Using

$$d_{j}^{\dagger}d_{j} = \frac{1}{4}(\gamma_{j}^{A} - i\gamma_{j+1}^{B})(\gamma_{j}^{A} + i\gamma_{j+1}^{B}) = \frac{1}{4}\left[2 + i(\gamma_{j}^{A}\gamma_{j+1}^{B} - \gamma_{j+1}^{B}\gamma_{j}^{A})\right] = \frac{1}{2} + \frac{i}{2}\gamma_{j}^{A}\gamma_{j+1}^{B} \quad (3.29)$$

the Hamiltonian can be written in terms of the "conventional" fermions operators

$$H = 2t_0 \sum_{1 \le j \le N-1} \left( d_j^{\dagger} d_j - \frac{1}{2} \right).$$
(3.30)

The significant point is that the missing end Majorana operators can be combined to form a conventional, but non-local, fermion operator

$$d_{\rm end} = \frac{1}{2} (\gamma_N^A + i \gamma_1^B), \qquad (3.31)$$

with

$$d_{\text{end}}^{\dagger} d_{\text{end}} = \frac{1}{4} (\gamma_N^A - i\gamma_1^B) (\gamma_N^A + i\gamma_1^B) = \frac{1}{4} (2 + 2i\gamma_N^A \gamma_1^B)$$
(3.32)

$$= \frac{1}{2} + \frac{i}{2}\gamma_N^A \gamma_1^B, \qquad (3.33)$$

or equivalently,

$$i\gamma_N^A \gamma_1^B = 2d_{\text{end}}^\dagger d_{\text{end}} - 1.$$
(3.34)

In order to include every operator in the Hamiltonian, one could write

$$H = E_{\rm g} + \epsilon_0 d_{\rm end}^{\dagger} d_{\rm end} + 2t_0 \sum_{1 \le j \le N-1} d_j^{\dagger} d_j \tag{3.35}$$

with  $E_{\rm g} = -2t_0(N-1)$  and  $\epsilon_0 = 0$ . Thus, Majorana operators at the ends of the chain form a zero-energy fermion, while the conventional localized fermions in the bulk of the chain require energy  $2t_0$ . Thus the system has a ground-state degeneracy consisting of the states with, and without, the non-local fermion. The operator that measures the presence of the non-local fermion is

$$-i\gamma_N^A \gamma_1^B = 1 - 2d_{\text{end}}^\dagger d_{\text{end}}, \qquad (3.36)$$

with eigenvalues -1 and +1 corresponding to the presence and absence of the non-local fermion, respectively. Thus, the Kitaev toy model also has two distinct phases, with and without the presence of a zero-energy Majorana mode, sometimes known as a Majorana zero mode.

While the case of general parameters requires more complicated calculations, the above result generalizes for all parameter values. However, instead of the MBS existing only at the end sites, the distribution of these two states decays exponentially into the wire, and thus have a small overlap that slightly breaks the degeneracy of the two ground states.

## 3.3 Majorana Wire System Band Structure

While the more realistic Majorana wire proposal by Lutchyn et. al. [73] and Oreg et al. [86] is more complicated then the Kitaev model, the band structure in the uniform case shows that similar states emerge. The full Hamiltonian of the system is given by

$$H_{\rm S} = H_{\rm TB} + H_{\rm SO} + H_{\rm Z} + H_{\rm SC},$$
 (3.37)

where

$$H_{\rm TB} = \sum_{j\sigma} \left[ (2t_0 - \mu) \hat{c}^{\dagger}_{j\sigma} \hat{c}_{j\sigma} - t_0 \hat{c}^{\dagger}_{j\pm 1,\sigma} \hat{c}_{j\sigma} \right], \qquad (3.38)$$

$$H_{\rm SO} = \sum_{j\sigma} \left[ \frac{\alpha}{2} s(\sigma) \left( \hat{c}^{\dagger}_{j\bar{\sigma}} \hat{c}_{j+1,\sigma} - \hat{c}^{\dagger}_{j+1,\bar{\sigma}} \hat{c}_{j\sigma} \right) \right], \qquad (3.39)$$

$$H_{\rm Z} = \sum_{j\sigma} \left[ s(\sigma) V^z \hat{c}^{\dagger}_{j\sigma} \hat{c}_{j\sigma} + V^{s(\sigma)} \hat{c}^{\dagger}_{j\bar{\sigma}} \hat{c}_{j\sigma} \right], \qquad (3.40)$$

$$H_{\rm SC} = \sum_{j} \left( \Delta \hat{c}_{j\uparrow}^{\dagger} \hat{c}_{j\downarrow}^{\dagger} + \Delta^* \hat{c}_{j\downarrow} \hat{c}_{j\uparrow} \right), \qquad (3.41)$$

are the tight-binding, spin-orbit, Zeeman, and proximity-effect superconducting terms, respectively. Here  $\hat{c}_{j\sigma}$  is the electron annihilation operator for spin  $\sigma$  at site j,  $t_0 = \hbar^2/(2m^*a^2)$  is the tight-binding coefficient with effective mass  $m^*$  and lattice size a,  $\mu$  is the chemical potential,  $\alpha/2$  is the Rashba coupling,  $\mathbf{V} = g\mu_{\rm B}\mathbf{B}/2$  is the Zeeman coupling with  $V^{\pm} = V^x \pm iV^y$  used for the terms perpendicular to the wire axis, and  $\Delta$  is the *s*-wave pairing potential. The coefficient  $s(\sigma)$  stands for + and - when  $\sigma$  is  $\uparrow$ and  $\downarrow$ , respectively, and  $\bar{\sigma}$  denotes the opposite spin.

Consider the tight-binding terms in the lattice Hamiltonian, transferring to momentum space by using

$$c_{j\sigma} = \sum_{k} \langle j\sigma | k \rangle c_{k\sigma} = \sum_{k} e^{ikja} c_{k\sigma}$$
(3.42)

where  $|j\sigma\rangle$  is the state at site j, at x = ja, and  $c_{k\sigma}$  annihilates an electron of momentum k and spin  $\sigma$ . Thus, the tight-binding term in k-space can be written

$$H = \sum_{j\sigma kk'} \left[ -t_0 (e^{-ik(j+1)a} e^{ik'ja} + e^{-ik(j-1)a} e^{ik'ja}) + (2t_0 - \mu) e^{-ikja} e^{ik'ja} \right] c^{\dagger}_{k\sigma} c_{k'\sigma}$$

$$= \sum_{j\sigma kk'} \left[ -t_0 (e^{-ika} e^{-i(k-k')ja} + e^{ika} e^{-i(k-k')ja}) + (2t_0 - \mu) e^{-i(k-k')ja} \right] c^{\dagger}_{k\sigma} c_{k'\sigma}$$

$$= \sum_{\sigma k} \left[ -t_0 (e^{-ika} + e^{ika}) + (2t_0 - \mu) \right] c^{\dagger}_{k\sigma} c_{k\sigma}$$

$$= \sum_{\sigma k} \left[ -2t_0 \cos(ka) + (2t_0 - \mu) \right] c^{\dagger}_{k\sigma} c_{k\sigma}$$

$$= \sum_{\sigma k} \left\{ 2t_0 [1 - \cos(ka)] - \mu \right\} c^{\dagger}_{k\sigma} c_{k\sigma}, \qquad (3.43)$$

while the spin-orbit term is written

$$\sum_{j\sigma kk'} \frac{\alpha}{2} s(\sigma) \left[ e^{-ikja} e^{ik'(j+1)a} c^{\dagger}_{k\bar{\sigma}} c_{k'\sigma} - e^{-ik(j+1)a} e^{ik'ja} c^{\dagger}_{k\bar{\sigma}} c_{k'\sigma} \right]$$
(3.44)

$$= \sum_{j\sigma kk'} \frac{\alpha}{2} s(\sigma) [e^{ik'a} - e^{-ika}] e^{-i(k-k')ja} c^{\dagger}_{k\bar{\sigma}} c_{k'\sigma}$$
(3.45)

$$= \sum_{\sigma k} \frac{\alpha}{2} s(\sigma) [e^{ika} - e^{-ika}] c_{k\bar{\sigma}}^{\dagger} c_{k\sigma} = \sum_{\sigma k} i\alpha s(\sigma) \sin(ka) c_{k\bar{\sigma}}^{\dagger} c_{k\sigma}, \qquad (3.46)$$

and the superconducting terms are

$$\sum_{jkk'} (\Delta e^{-ikja} e^{-ik'ja} c^{\dagger}_{k\uparrow} c^{\dagger}_{k\downarrow} + \Delta^* e^{ikja} e^{ik'ja} c_{k\downarrow} c_{k'\uparrow})$$
(3.47)

$$= \sum_{ikk'} (\Delta e^{-i(k+k')ja} c^{\dagger}_{k\uparrow} c^{\dagger}_{k'\downarrow} + \Delta^* e^{i(k+k')ja} c_{k\downarrow} c_{k'\uparrow})$$
(3.48)

$$= \sum_{k} (\Delta c_{k\uparrow}^{\dagger} c_{-k\downarrow}^{\dagger} + \Delta^{*} c_{-k\downarrow} c_{k\uparrow}). \qquad (3.49)$$

This can be repeated for the remaining terms in the lattice Hamiltonian to give

$$H = \sum_{\sigma k} \left\{ 2t_0 [1 - \cos(ka)] - \mu + s(\sigma) V^z \right\} c^{\dagger}_{k\sigma} c_{k\sigma}$$
(3.50)

$$+\sum_{\sigma k} \left[ i\alpha s(\sigma) \sin(ka) + V^{s(\sigma)} \right] c^{\dagger}_{k\bar{\sigma}} c_{k\sigma}$$
(3.51)

$$+\sum_{k} (\Delta c_{k\uparrow}^{\dagger} c_{-k\downarrow}^{\dagger} + \Delta^{*} c_{-k\downarrow} c_{k\uparrow}).$$
(3.52)

Separating the terms with negative-momentum and reordering gives

$$H = \sum_{k \ge 0,\sigma} \{ 2t_0 [1 - \cos(ka)] - \mu + s(\sigma) V^z \} (c_{k\sigma}^{\dagger} c_{k\sigma} + c_{-k\sigma}^{\dagger} c_{-k\sigma})$$
(3.53)

$$+\sum_{k\geq 0,\sigma} \left[i\alpha s(\sigma)\sin(ka) + V^{s(\sigma)}\right] c^{\dagger}_{k\bar{\sigma}}c_{k\sigma}$$
(3.54)

$$+\sum_{k\geq 0,\sigma} \left[-i\alpha s(\sigma)\sin(ka) + V^{s(\sigma)}\right] c^{\dagger}_{-k\bar{\sigma}}c_{-k\sigma}$$
(3.55)

$$+\sum_{k\geq 0} \Delta (c_{k\uparrow}^{\dagger} c_{-k\downarrow}^{\dagger} + c_{-k\uparrow}^{\dagger} c_{k\downarrow}^{\dagger}) + \Delta^{*} (c_{-k\downarrow} c_{k\uparrow} + c_{k\downarrow} c_{-k\uparrow})$$
(3.56)

$$= \sum_{k \ge 0,\sigma} \left\{ 2t_0 [1 - \cos(ka)] - \mu + s(\sigma) V^z \right\} \left[ c_{k\sigma}^{\dagger} c_{k\sigma} + (1 - c_{-k\sigma} c_{-k\sigma}^{\dagger}) \right] \quad (3.57)$$

$$+\sum_{k\geq 0,\sigma} \left[i\alpha s(\sigma)\sin(ka) + V^{s(\sigma)}\right] c^{\dagger}_{k\bar{\sigma}}c_{k\sigma}$$
(3.58)

$$+\sum_{k\geq 0,\sigma} \left[i\alpha s(\sigma)\sin(ka) - V^{s(\sigma)}\right] c_{-k\sigma} c^{\dagger}_{-k\bar{\sigma}}$$
(3.59)

$$+\sum_{k\geq 0} \Delta (c_{k\uparrow}^{\dagger}c_{-k\downarrow}^{\dagger} - c_{k\downarrow}^{\dagger}c_{-k\uparrow}^{\dagger}) + \Delta^{*}(c_{-k\downarrow}c_{k\uparrow} - c_{-k\uparrow}c_{k\downarrow})$$
(3.60)
By defining  $h(k) = 2t_0[1 - \cos(ka)] - \mu$ , one can write the momentum-space Hamiltonian concisely in matrix form as  $H = E_g^k + \sum_{k\geq 0} K^{\dagger} H_k K$  with  $E_g^k = 2 \sum_{k\geq 0} h(k)$ ,

$$K^{\dagger} = \begin{pmatrix} c^{\dagger}_{k\uparrow} & c^{\dagger}_{k\downarrow} & c_{-k\uparrow} & c_{-k\downarrow} \end{pmatrix}$$
(3.61)

and

$$H_{k} = \begin{pmatrix} h(k) + V^{z} & -i\alpha\sin(ka) + V^{-} & 0 & \Delta \\ i\alpha\sin(ka) + V^{+} & h(k) - V^{z} & -\Delta & 0 \\ 0 & -\Delta^{*} & -h(k) - V^{z} & i\alpha\sin(ka) - V^{+} \\ \Delta^{*} & 0 & -i\alpha\sin(ka) - V^{-} & -h(k) + V^{z} \end{pmatrix},$$
(3.62)

which has the expected block form

$$\begin{pmatrix} H(k) & \Delta \\ -\Delta^* & -H^*(-k) \end{pmatrix}.$$
 (3.63)

The eigenvalues can be found, but are generally quite complicated. Instead, consider the simpler case of just the tight-binding, Zeeman, and spin-orbit terms. Without the superconducting terms, the eigenvalues of the top-left  $2 \times 2$  block are

$$\epsilon_{\pm}(k) = h(k) \pm \sqrt{V_{\perp}^2 + [V_y + \alpha \sin(ka)]^2},$$
(3.64)

where  $V_{\perp}^2 = V_x^2 + V_z^2$  is the Zeeman field perpendicular to the spin-quantization axis caused by the spin-orbit coupling. The relevant physics occurs near k = 0, so consider the  $k \ll a$  limit. The tight-binding and spin-orbit terms give

$$\epsilon_{\pm}(k) = t_0(ka)^2 \pm \alpha ka - \mu \tag{3.65}$$

$$\iff \epsilon/t_0 = x^2 \pm \frac{\alpha}{t_0} x - \frac{\mu}{t_0} = x \left( x \pm \frac{\alpha}{t_0} \right) - \frac{\mu}{t_0}$$
(3.66)



Figure 3.1: Band structure with  $\alpha = 6.666$ ,  $t_0 = 11.3$ , and  $V_{\perp} = 0.5$ : Dashed is tight-binding and spin-orbit only, while solid also includes perpendicular Zeeman field, plotting with  $\mu = 0$  and implicitly understanding that  $\mu$  is the y = 0 axis. The units for the y and x axis are  $t_0$  and 1/a respectively.

with  $x \equiv ka$ , so the spin-orbit shifts the spin bands to be centered at  $\pm \frac{\alpha}{2t_0}$  with  $k_{\rm F} = \pm (\alpha + \sqrt{\alpha^2 + 4\mu t_0})/2t_0 a$ . The introduction of the perpendicular Zeeman terms gives

$$\epsilon_{\pm}/t_0 = x^2 \pm \frac{1}{t_0}\sqrt{V_{\perp}^2 + \alpha^2 x^2} - \frac{\mu}{t_0},$$
(3.67)

which opens up a gap of  $2V_{\perp}/t_0$  at k = 0.

For the introduction of the superconducting pairing terms, consider the whole  $4 \times 4$  matrix, which has the particle-hole symmetry that results in symmetric eigenvalues about E = 0. For this reason, consider the square of the excitation spectrum

$$\epsilon_{\pm}^{2}(k) = h^{2} + V_{\perp}^{2} + \Delta^{2} + \alpha^{2} \sin^{2}(ka) \pm 2\sqrt{V_{\perp}^{2}(h^{2} + \Delta^{2}) + h^{2}\alpha^{2} \sin^{2}(ka)}, \quad (3.68)$$

which at k = 0 is

$$\epsilon_{\pm}^{2}(0) = \mu^{2} + V_{\perp}^{2} + \Delta^{2} \pm 2V_{\perp}\sqrt{\mu^{2} + \Delta^{2}}$$
(3.69)

$$= (V_{\perp} \pm \sqrt{\mu^2 + \Delta^2})^2. \tag{3.70}$$

So  $\epsilon_{-}(0) = |V_{\perp} - \sqrt{\mu^2 + \Delta^2}|$ . Thus the gap closes when  $V_{\perp} = \sqrt{\mu^2 + \Delta^2}$  and, just like the Kitaev toy model, this marks the transition between the two distinct phases.

To see which side of this transition yields localized modes, consider the band structure in the figure above. With just the tight-binding and spin-orbit terms present, the Kramer's degeneracy is split, forming two spin-bands parallel and anti-parallel to the spin-orbit quantization axis (shown in the figure as up and down arrows). The introduction of the Zeeman term perpendicular to that axis results in a gap which creates two bands with different "spin-momentum" species. In the upper band, the electrons with positive momentum have one spin (down in the figure), while the electrons with the opposite momentum have the opposite spin (up in the figure). Similarly in the lower band, the momentum and spin are tied together. This relationship between spin and momentum, which effectively reduces the degrees of freedom of the system, is known as a helical liquid. If the chemical potential lies in the gap (and  $k_{\rm B}T \ll (V_{\perp} - \mu)/t_0$ ) the upper band becomes energetically unavailable and the system can be projected to the lower band. While the lower band still contains opposite spins for opposite momentum (as required for the s-wave pairing which will be introduced) the spin is not a true degree of freedom; in some sense the system is "spinless" like Kitaev's toy model.

As mentioned above, the energy spectrum for more general parameters in this system can be found analytically, but are quite complicated and not particularly helpful. Instead, the lowest energy level and the gap to next highest energy level was simulated using the code in App. B, and are plotted below with  $V_z = 0.9$ ,  $V_y = V_x = 0$ ,  $\Delta = 0.5$ ,  $\mu = 0$  and  $\alpha = 3$ .



Figure 3.2: Modes near zero energy are robust near  $V_z \sim 0.5$ , and small  $V_y < 0.3$ . It should be noted that  $V_x$  could be used equivalently since it is also perpendicular to the spin-orbit quantization axis.



Figure 3.3: Modes near zero energy are robust near  $V_z \sim 0.9$ , and large  $\alpha > 1.5$ . There seems to be little benefit to increasing  $\alpha$  beyond 3.



Figure 3.4: The results are in general agreement that the topological threshold is given by  $\mu^2 < V_z^2 - \Delta^2$ , which for  $\mu = 0$  is below the line of slope 1. On the other hand, a  $\Delta$ value that is too small ( $\Delta < 0.3$ ) doesn't create a sufficiently large gap, so this equation doesn't apply for small  $\Delta$ . Thus, robust modes only appear for  $V_z \ge 0.5$ , and decent window of feasible  $\Delta$  doesn't begin until  $V_z \ge 0.9$ .



Figure 3.5: The results are in general agreement that the topological threshold is given by  $\mu^2 < V_z^2 - \Delta^2$ , which is a circle of radius  $V_z$ . On the other hand, a  $\Delta$  value that is too small ( $\Delta < 0.2$ ) doesn't create a sufficiently large gap, so this equation doesn't apply for small  $\Delta$ . Thus, robust modes are best achieved with higher  $V_z \sim 0.9$ , where  $\Delta \sim 0.5$  is ideal for small  $\mu$  values.



Figure 3.6: The results are in general agreement that the topological threshold is given by  $\mu^2 < V_z^2 - \Delta^2$ , which is plotted as a red line. Robust modes are best achieved with higher  $V_z > 0.75$  for small  $\mu$  values.

### Chapter 4

# Majorana Wire

#### 4.1 Introduction

While Majorana zero modes bound to the ends of semiconductor nanowires as discussed in the previous chapter are theoretically supported by models, further evidence is needed to rule out alternative explanations [71, 101, 68]. Perhaps the most definitive signature of Majorana bound states in these "Majorana wires" would be the demonstration of their non-trivial braiding statistics. While braiding is ultimately needed for topological quantum computation, it remains an ambitious experimental task. With this in mind, simpler experiments are desired to provide insight and direct further research before braiding is attempted. Though there have been feasible tests proposed and performed on several aspects of the system [50, 49, 15, 21, 99, 72, 103], such as qubit measurement, there is still no clear consensus on the presence of Majorana bound states [37].

Observing entanglement of these states in Majorana wires would not only be a significant step towards their verification, but would also demonstrate their potential utility for topological quantum computation. While tests of quantum entanglement with Ising anyons have been discussed formally [13, 16], the goal of this chapter is to devise and analyze a more concrete protocol motivated by the recent experimental developments discussed in the previous chapter.

Thus we propose a procedure for demonstrating Bell's theorem with three pairs of Majorana bound states in semiconductor nanowire systems [see Fig. 4.1]. Specifically, our procedure can be used to test the Bell [7] and Clauser-Horne-Shimony-Holt (CHSH) [19] inequalities using only two operations on maximally entangled states, which can be prepared using the same operations and projective measurement [see Fig. 4.2]. These operations are accomplished by moving the domain walls along the axis of the wire using "keyboard" gates already needed for braiding [1]. Hence, our proposal may also serve as a step towards experiments that perform topological operations.

The remainder of this chapter will proceed as follows: Sec. 4.2 introduces a simplified model for the Majorana wire and defines a qubit basis. Sec. 4.3 summarizes the entanglement inequalities and lay out the procedure for testing them. Sec. 4.4 introduces a more realistic description of the semiconductor nanowire system, discuss corresponding simulation results, and introduces a simpler version of the CHSH experiment. Sec. 4.5 discusses experimental considerations and Sec. 4.7 summarizes the proposal. Modifications to this procedure for different measurement outcomes are discussed in Section 4.6.



Figure 4.1: The wires are segmented into three regions where the wires are in the topological phase (solid lines). Majorana bound states, represented by red x's, are localized at the ends of these regions. Majorana bound states at the ends of the same topological region are coupled by  $\eta$ , while neighboring topological regions are coupled by  $\Gamma$ .



Figure 4.2: Preparation of the maximally entangled states of even total parity. A) The occupation of all three topological regions is measured, represented by rectangles around each region. B) The topological regions (solid lines) are expanded to perform  $\pi/2$  rotations about the *x*-axis for the left and right logical qubits. C) The middle qubit is measured, projecting to one of the four maximally entangled states of even total parity. Different measurement outcomes are shown; when both middle measurements are 0, the  $|\Phi_{\rm E}^-\rangle$  state is prepared as discussed in Sec. 4.3 (solid arrows), while other outcomes (dashed arrows) are discussed in Section 4.6.

#### 4.2 Majorana Model Hamiltonian

To discuss the salient features of the Majorana wire system, begin by considering a description similar to the toy model analyzed by Kitaev [61]. With appropriate parameters, the wire is driven into a topological phase with an unpaired Majorana fermion at each end [73, 86, 1]. If the parameters vary spatially (e.g., non-uniform chemical potential) there may be multiple topological regions separated by non-topological regions, with a Majorana fermion localized at each domain wall separating the two regions. For this proposal, consider the case with three topological regions separated by two non-topological regions [see Fig. 4.1] with the Majorana Hamiltonian

$$H = i\eta_1 \hat{\gamma}_{1,A} \hat{\gamma}_{1,B} + i\eta_2 \hat{\gamma}_{2,A} \hat{\gamma}_{2,B} + i\eta_3 \hat{\gamma}_{3,A} \hat{\gamma}_{3,B} + i\Gamma_{12} \hat{\gamma}_{1,B} \hat{\gamma}_{2,A} + i\Gamma_{23} \hat{\gamma}_{2,B} \hat{\gamma}_{3,A}, \quad (4.1)$$

where  $\eta$  describes the coupling between Majorana bound states at the ends of a single topological region, while  $\Gamma$  describes the coupling of neighboring topological regions, assuming that all couplings decay exponentially as the Majorana bound states separate from their nearest neighbors. Each topological region has two types of Majorana operators, denoted by index A or B, that form a conventional fermion operator  $\hat{d}_n = \frac{1}{2}(\hat{\gamma}_{n,A} + i\hat{\gamma}_{n,B})$ , and satisfy  $\{\hat{\gamma}_i, \hat{\gamma}_j\} = 2\delta_{ij}$ , where i, j specifies both the region and type.

The parity of the occupation number for the conventional fermions, (i.e., the eigenstate of  $\hat{N}_n \equiv \hat{d}_n^{\dagger} \hat{d}_n$ ), will serve as the degree of freedom for our qubits. A computational basis is specified with the conventional fermions by defining the state  $|000\rangle$  such that  $\hat{d}_n |000\rangle = 0$  for all n and using the ordering conventions given by

$$|000\rangle \qquad |010\rangle = \hat{d}_{2}^{\dagger}|000\rangle$$

$$|011\rangle = \hat{d}_{2}^{\dagger}\hat{d}_{3}^{\dagger}|000\rangle \qquad |001\rangle = \hat{d}_{3}^{\dagger}|000\rangle$$

$$|110\rangle = \hat{d}_{1}^{\dagger}\hat{d}_{2}^{\dagger}|000\rangle \qquad |100\rangle = \hat{d}_{1}^{\dagger}|000\rangle \qquad (4.2)$$

$$|101\rangle = \hat{d}_{1}^{\dagger}\hat{d}_{3}^{\dagger}|000\rangle \qquad |111\rangle = \hat{d}_{1}^{\dagger}\hat{d}_{2}^{\dagger}\hat{d}_{3}^{\dagger}|000\rangle.$$

Since this model describes a system with superconductivity, the total number of particles is conserved modulo 2. This restriction splits the basis into two sub-bases,  $S_{\rm E}$  and  $S_{\rm O}$ , with an even and odd number of total particles (i.e., total parity), which are the left and right columns of Eqs. (4.2), respectively. A state from one basis cannot evolve into a state from the other basis since they differ by a single particle. Strictly speaking, the two bases can interact if we account for quasi-particle poisoning in our model [94], but this occurs on a much longer time-scale than our proposed operations as discussed in Sec. 4.5. The middle occupation number is used to preserve the total parity rather than storing unique quantum information. Thus, two logical qubits are encoded in the left and right topological regions while the occupation of the middle region is forfetied as the "parity qubit".

By writing the Majorana operators in terms of the conventional fermions with  $\hat{\gamma}_{n,A} = \hat{d}_n + \hat{d}_n^{\dagger}$  and  $i\hat{\gamma}_{n,B} = \hat{d}_n - \hat{d}_n^{\dagger}$ , the Hamiltonian in this basis is

$$H = -\eta_1(\sigma^z \otimes \sigma^0 \otimes \sigma^0) - \eta_2(\sigma^0 \otimes \sigma^z \otimes \sigma^0) - \eta_3(\sigma^0 \otimes \sigma^0 \otimes \sigma^z) - \Gamma_{12}(\sigma^x \otimes \sigma^x \otimes \sigma^0) - \Gamma_{23}(\sigma^0 \otimes \sigma^x \otimes \sigma^x).$$
(4.3)

The  $\eta$  terms for each topological region perform the  $\sigma^z$  operation for their corresponding qubits, while the  $\sigma^x$  operation is performed on the neighboring qubits involved in the  $\Gamma$ terms. Thus rotations on the Bloch spheres of the qubits can be made by adjusting the parameters of the wire to suppress the couplings of all but one term in the Hamiltonian. For example, if all the couplings other than  $\Gamma_{12}$  are negligible, the evolution operator after time T is

$$r_{12}^{x}(\theta) \equiv \exp\left[i\frac{\theta}{2}(\sigma^{x}\otimes\sigma^{x}\otimes\sigma^{0})\right]$$
(4.4)

$$= \cos\frac{\theta}{2}(\sigma^0 \otimes \sigma^0 \otimes \sigma^0) + i\sin\frac{\theta}{2}(\sigma^x \otimes \sigma^x \otimes \sigma^0), \qquad (4.5)$$

where  $\theta = 2\Gamma_{12}T/\hbar$  is the angle that qubits 1 and 2 rotate about the *x*-axis of their respective Bloch spheres. By adjusting the parameters appropriately, one can perform all the Bloch sphere rotations necessary for this proposal.

#### 4.3 Entanglement Inequalities

Before testing the Bell and CHSH inequalities, this section discusses the preparation of one of the four maximally entangled states of even parity,

$$|\Phi_{\rm E}^{\pm}\rangle = \frac{|000\rangle \pm |101\rangle}{\sqrt{2}}, \qquad |\Psi_{\rm E}^{\pm}\rangle = \frac{|011\rangle \pm |110\rangle}{\sqrt{2}} \tag{4.6}$$

using the operations already discussed and projective measurement. To begin the preparation, the parity of each topological region is measured, fixing the total parity and projecting to one of the basis states. The inequalities can be tested equivalently with any of the maximally entangled states from either parity, but for conciseness, assume the total parity is even for the rest of the body of this paper, and consider only the inequalities with  $|\Phi_{\rm E}^-\rangle$ , assuming the initially measured state is  $|000\rangle$ . This proposal can be accomplished for general initial conditions by altering the procedure slightly as described in Section 4.6. If a  $\pi/2$  rotation about the *x*-axis is performed for both logical qubits the resulting state is  $r_{12}^x(\pi/2)r_{23}^x(\pi/2)|000\rangle$ , or

$$\frac{|000\rangle - |101\rangle + i|011\rangle + i|110\rangle}{2} = \frac{|\Phi_{\rm E}^-\rangle + i|\Psi_{\rm E}^+\rangle}{\sqrt{2}},\tag{4.7}$$

which will project to  $|\Phi_{\rm E}^-\rangle$  if the middle parity qubit is measured to be 0. Note that the  $r_{12}^x$  and  $r_{23}^x$  operations commute since they involve different  $\gamma$  operators, making the operation order irrelevant (as well as allowing simultaneous operations). In general, each maximally entangled state can be prepared by measuring all three qubits to project to a single basis state, extending the outer topological regions towards the middle topological region for a small time, returning them to their original position, then projectively measuring the middle qubit. Once the state  $|\Phi_{\rm E}^-\rangle$  is prepared, one can test the version of Bell's inequality given in Section 1.3,

$$P_{=}(a,b) + P_{=}(b,c) + P_{=}(a,c) \ge 1, \tag{4.8}$$

where  $P_{=}(L, R)$  is the probability that the left and right qubits are equal after being rotated by angles L and R, respectively. The left side of the inequality, which will be called the "Bell quantity", can be interpreted as the probability that at least one of the rotation combinations will make the left and right qubits equal.

According to quantum mechanics the probability that the qubits are equal after rotations L and R is  $\cos^2\left(\frac{L-R}{2}\right)$ . Only the relative angles between rotations are physically relevant, so one can set  $A \equiv a - c$  and  $B \equiv b - c$  to write the Bell quantity as

$$\cos^{2}\left(\frac{A-B}{2}\right) + \cos^{2}\left(\frac{A}{2}\right) + \cos^{2}\left(\frac{B}{2}\right),\tag{4.9}$$

which is plotted in Fig. 4.3. Quantum mechanics predicts the Bell quantity can be as low as 3/4 (for the relative angles  $A = 2\pi/3$  and  $B = 4\pi/3$ , or vice-versa) and is inconsistent with local hidden variable theories, which require the Bell quantity to be greater than or equal to 1. In principle, Bell's inequality could be experimentally tested in our proposal by repeatedly preparing maximally entangled states, performing the three rotation combinations in Eq. (4.8), and measuring the qubits to find the probability of each state.

In practice however, almost every experiment that tests Bell's theorem uses the CHSH inequality discussed in Appendix 1.3,

$$|\langle L_1, R_1 \rangle + \langle L_2, R_1 \rangle + \langle L_1, R_2 \rangle - \langle L_2, R_2 \rangle| \le 2, \tag{4.10}$$

where  $\langle L, R \rangle = P_{=}(L, R) - P_{\neq}(L, R)$  is the expectation value of the combined parity of the left and right qubits after being rotated by angles L and R, respectively. The left



Figure 4.3: Contour plot for the quantum mechanical prediction of the Bell quantity for the state  $|\Phi_{\rm E}^-\rangle$ . Local hidden variable theories require that the Bell quantity be greater than or equal to 1, but it is predicted to be less than 1 for relative rotation angles inside the white triangles.

side of the inequality, which will be called the "CHSH quantity", must be less than or equal to 2 in local hidden variable theories.

According to quantum mechanics, the expectation value discussed above for general rotation angles L and R is simply  $\cos(L - R)$ . Again, only the relative angles of rotation are physically significant, so we introduce angles  $A \equiv L_1 - R_2$ ,  $B \equiv R_1 - L_1$ , and  $C \equiv L_2 - R_1$ , [see Fig. 4.4], making the CHSH quantity

$$\left|\cos(A) + \cos(B) + \cos(C) - \cos(A + B + C)\right|, \tag{4.11}$$

which has a maximum of  $2\sqrt{2}$  when  $A = B = C = \pi/4$ , contradicting the local hidden variable prediction. The inequality can be tested experimentally by repeatedly preparing the state  $|\Phi_{\rm E}^-\rangle$ , extending the topological regions to perform one of the four rotation combinations involved in Eq. (4.10), then returning the topological regions to their original position to measure the qubits.



Figure 4.4: Top: Angles of rotation in CHSH inequality. The left qubit is rotated by either angle  $L_1$  or  $L_2$ , while the right qubit is rotated by either angle  $R_1$  or  $R_2$ . Bottom: A) One of the four rotation combinations is performed by extending the outer topological regions, B) then returned for measurement.

#### 4.4 Semiconductor Hamiltonian and Simulation

Consider a more realistic model of the semiconductor system by re-introducing the one-dimensional lattice Hamiltonian

$$H_{\rm S} = H_{\rm TB} + H_{\rm SO} + H_{\rm Z} + H_{\rm SC},$$
 (4.12)

where

$$H_{\rm TB} = \sum_{j\sigma} \left[ (2t_0 - \mu_j) \hat{c}^{\dagger}_{j\sigma} \hat{c}_{j\sigma} - t_0 \hat{c}^{\dagger}_{j\pm 1,\sigma} \hat{c}_{j\sigma} \right], \qquad (4.13)$$

$$H_{\rm SO} = \sum_{j\sigma} \left[ \alpha \, s(\sigma) \left( \hat{c}^{\dagger}_{j\bar{\sigma}} \hat{c}_{j+1,\sigma} - \hat{c}^{\dagger}_{j+1,\bar{\sigma}} \hat{c}_{j\sigma} \right) \right], \tag{4.14}$$

$$H_{\rm Z} = \sum_{j\sigma} \left[ s(\sigma) V^z \hat{c}^{\dagger}_{j\sigma} \hat{c}_{j\sigma} + V^{s(\sigma)} \hat{c}^{\dagger}_{j\bar{\sigma}} \hat{c}_{j\sigma} \right], \qquad (4.15)$$

$$H_{\rm SC} = \sum_{j} \left( \Delta \hat{c}_{j\uparrow}^{\dagger} \hat{c}_{j\downarrow}^{\dagger} + \Delta^* \hat{c}_{j\downarrow} \hat{c}_{j\uparrow} \right), \qquad (4.16)$$

where all the terms are the same as defined in the uniform case of Sec. 3.3, except that the chemical potential,  $\mu_i$ , is allowed to vary at each site.

When  $\mu > \mu_{\rm T} \equiv \sqrt{V_{\perp}^2 - \Delta^2}$ , where  $V_{\perp}^2 = (V^z)^2 + (V^x)^2$  is the Zeeman field perpendicular to the spin-orbit quantization axis, the wire is a normal superconductor. In the other case,  $\mu < \mu_{\rm T}$ , a topologically distinct state emerges with Majorana bound states localized at the ends of the wire. If the chemical potential varies spatially and crosses the topological limit at multiple locations, then multiple Majorana bound states will be present and the setup discussed in Sec. 4.2 is possible.

Specifically, this proposal separates the wire into three topological regions, leading to six Majorana bound states, one at the end of each region, which are sufficiently separated to prevent them from fusing together. The Majorana bound states from each topological region can be paired together to form conventional fermions [e.g.,  $\hat{d}_n =$  $(\hat{\gamma}_A + i\hat{\gamma}_B)/2$ ] that correspond to three zero-energy (in the limit of an infinite wire) Bogoliubov excitations, separated from the higher-energy bulk states by a topological gap  $\Delta_{\rm T}$  [see Fig. 4.5]. Alternatively, these excitations can be thought of as the zero-energy eigenstate solutions to the Bogoliubov-de Gennes equations for the system. Just as in the simpler model, the occupation number of the eigenstates localized to the left and right serve as the logical qubits, while the occupation number of the middle eigenstate does not contain unique quantum information since the total parity is conserved.

The spatial distribution of these excitations is contained in the coefficients uand v from the Bogoliubov transformation

$$\hat{d}_n^{\dagger} = \sum_{j\sigma} (u_{j\sigma}^n \hat{c}_{j\sigma}^{\dagger} + v_{j\sigma}^n \hat{c}_{j\sigma}), \qquad (4.17)$$

which can be used to form the parity operators of the wire segments,  $\hat{P}_n \equiv 1 - 2\hat{d}_n^{\dagger}\hat{d}_n$ .

The simulation begins by finding the coefficients for the lowest three eigenstates of the Hamiltonian in Eqs. (4.12)-(4.16) with parameters corresponding to  $m^* = 0.015m_{\rm e}$ , a = 15nm leading to  $t_0 = 11.3$ meV, g = 50,  $B = B^z = 0.625$ T leading to  $V_{\perp} = 0.9$ meV, Rashba parameter  $\alpha_{\rm R} = 225$ meV·Å corresponding to  $\alpha \equiv \alpha_{\rm R}/a = 1.5$ meV, and  $\Delta = 0.5$ meV. Thus the chemical potential marking the threshold between topological phases is  $\mu_{\rm T} = 1.06$ meV. The wire has 600 sites leading to length  $l = 9\mu$ m, with non-periodic boundary conditions. At the domain walls the chemical potential smoothly alternates between 0 and  $2\mu_{\rm T}$  over a length of approximately  $4\lambda = 0.04l$  with the profile function  $\pm \mu_{\rm T} \tanh(x/\lambda)$ , as shown in Fig. 4.5.

Each of the lowest three eigenstates has two peaks localized at the ends of the topological region, indicating the location of the Majorana bound states. The spatial distribution of the left and right bound states of each wire segment can be found numerically by considering the self-conjugate combinations of the three eigenstates, corresponding to  $\hat{d}_n + \hat{d}_n^{\dagger}$  and  $-i(\hat{d}_n - \hat{d}_n^{\dagger})$ , shown as solid and dashed curves in Fig. 4.5, respectively. Though the peaks decrease exponentially, their small, but non-zero over-



Figure 4.5: Bottom Right: A spatially varying chemical potential with three regions below the topological threshold of  $\sqrt{V_{\perp}^2 - \Delta^2}$ , with domain wall lengths of ~ 4 $\lambda$ . Top Right: This leads to six Majorana bound states, one at the end of each region, that form three conventional eigenstates. The simulated spatial distribution of the Bogoliubov coefficients  $\sum_{\sigma} (|u_{\sigma}|^2 + |v_{\sigma}|^2)$  along the length of the wire for the left and right bound states of each region s plotted in solid and dashed, respectively. Left: The energy spectrum of these eigenstates is plotted in log scale, as well as the lowest-energy bulk state separated by a topological gap of  $360\mu$ eV. The splitting of the "zero"-energy states is due to the exponentially small overlap in peaks, which is larger for the shorter topological region of the middle segment. The topological region lengths are 13.3% and 9.5% of the wire length for the outer and middle regions, respectively.

laps cause the eigenstates to split from zero-energy. Thus the topological regions must be long enough to prevent these overlaps from splitting the excitations and coupling them to the bulk states. The lengths between the domain walls are set as  $l_1 = x_1 = 0.133l$ ,  $l_2 = x_3 - x_2 = 0.095l$ , and  $l_3 = l - x_4 = 0.133l$ . These lengths were chosen to minimize the overlaps between the Majorana bound states of the same region, as well as the overlap between neighboring regions.

Once the chemical potential is tuned as described above, it can be varied dynamically to perform operations on the qubits. The only operation needed to test entanglement inequalities are  $r_{12}^x$  and  $r_{23}^x$ , which can be performed simultaneously by extending the outer topological regions towards the middle region. Specifically the domain wall positions  $x_i$  alternate back and forth according to the function

$$\pm \Lambda \left[ \tanh\left(\frac{t}{\tau}\right) - \tanh\left(\frac{t-D}{\tau}\right) \right], \tag{4.18}$$

which smoothly shifts the domain walls  $2\Lambda$  over an approximate transition time of  $4\tau$  for a duration D between the center of the two transitions as shown in Fig. 4.6.

One only needs to consider the dynamics of the zero-energy states, which won't mix with the bulk states above the topological gap as long as the domain wall trajectories are adiabatic. This constraint can be treated with the Landau-Zener condition [66, 122]: the rate the chemical potential changes must satisfy  $\hbar |d\mu/dt| \ll 2\pi \Delta_{\rm T}^2$ . To test this in the simulation, it finds the probability that the basis states of Eqs. (4.2) remain in the zero-energy subspace of the same total parity after evolution by using the following procedure.

The initial state  $|\phi\rangle$  is assumed to be in the set of even parity basis states given in Eqs. (4.2),  $S_{\rm E} = \{|000\rangle, |011\rangle, |110\rangle, |101\rangle\}$ , where  $|000\rangle$  is defined as the state such that  $\hat{d}_n |000\rangle = 0$  for all n, including those corresponding to bulk states. The zero-energy



Figure 4.6: Top Left: Trajectories for the two left domain walls showing the amplitude, transition time, and duration for the  $r_{12}^x(\pi/4)$  operation. Top Right: Average infidelity of even states after performing  $r_{12}^x(\pi/4)$ , plotted against transition time for various amplitudes (labeled as percentages of the wire length l) showing exponential behavior in general agreement with the Landau-Zener formula until limited by the Runge-Kutta step-size. The  $\Lambda = 0.055l$  data is fit with a line that scales as  $\exp(-\beta\tau)$  with  $\beta = 240$ Ghz, reasonably close to the predicted value of 214 GHz. Bottom: Probabilities that the initial state  $|000\rangle$  remains in  $|000\rangle$  or transitions to  $|110\rangle$  when acted on by  $r_{12}^x$  with various duration times. The simulated operation agrees well with the expected rotation, with a minimum probability of ~ 0.3% for the  $|000\rangle$ , due to the very small unintended overlaps of the bound states.

eigenstates are evolved using fourth-order Runge-Kutta in the eigenstate basis, possibly leaking into the bulk states if the transition time is too short, then acted on by the zeroenergy eigenstate projector  $\hat{P}_0$  to find the sub-matrix  $\hat{U}_0 = \hat{P}_0 \hat{U} \hat{P}_0$  of the full evolution matrix  $\hat{U}$ . Using  $\hat{U}_0$ , the time-evolved occupation operators in the Heisenberg picture are,

$$\hat{N}_n(t) = \hat{U}_0^{\dagger}(t)\hat{d}_n^{\dagger}(0)\hat{d}_n(0)\hat{U}_0(t), \qquad (4.19)$$

for n = 1, 2, 3, which are bilinear combinations of the original  $\hat{d}_n(0)$ , including anomalous terms (e.g.,  $\hat{d}_1 \hat{d}_2$ ) since the Hamiltonian contains superconductor pairing. The three  $\hat{N}_n(t)$  are then used to form the projector for each multi-particle state in the threequbit basis, allowing one to calculate the probability that the corresponding state is occupied. For example, the state  $|110\rangle$  has the projector  $\hat{N}_1 \hat{N}_2 (\hat{1} - \hat{N}_3)$ , which yields 0 when acting on any other basis state. Since the projector's eigenvalue for  $|110\rangle$  is 1, the expectation value is equal to the probability, and the probability that the initial state will be measured in the state  $|110\rangle$  after time t is

$$P_{110}(\phi) = \langle \phi | \hat{N}_1(t) \hat{N}_2(t) [\hat{1} - \hat{N}_3(t)] | \phi \rangle, \qquad (4.20)$$

which can be easily calculated for any  $|\phi\rangle \in S_{\rm E}$ . Similarly, the probabilities that  $|\phi\rangle$  evolves into the other states in  $S_{\rm E}$  are found using their respective multi-particle projectors, which are summed to give the fidelity from  $|\phi\rangle$  to  $S_{\rm E}$ ,

$$F_{\rm E}(\phi) = P_{000}(\phi) + P_{011}(\phi) + P_{110}(\phi) + P_{101}(\phi). \tag{4.21}$$

This is calculated for all  $|\phi\rangle \in S_{\rm E}$  after simulating the operation  $r_{12}^x(\pi/4)$ , and plot the average infidelity for even states,  $1 - \langle F_{\rm E} \rangle$ , versus the transition time for various amplitudes in Fig. 4.6. The results are compared with the Landau-Zener formula[66, 122] by using the maximum of the chemical potential rate

$$\frac{d\mu}{dt} = \frac{\partial\mu}{\partial x_1} \frac{\partial x_1}{\partial t},\tag{4.22}$$

which occurs halfway through the transitions when  $\operatorname{sech}^2(0) = 1$ , giving

$$\left(\frac{d\mu}{dt}\right)_{\rm max} = \frac{\mu_{\rm T}\Lambda}{\lambda\tau}.$$
(4.23)

Thus  $1 - \langle F_{\rm E} \rangle$  should scale as  $\exp(-\beta \tau)$  with

$$\beta = \frac{2\pi\lambda\Delta_{\rm T}^2}{\hbar\mu_{\rm T}\Lambda},\tag{4.24}$$

in general agreement with the data. For example, the fitted line for  $\Lambda = 0.055l$  in the logarithmic plot in Fig. 4.6 has a slope that corresponds to  $\beta_{\rm fit} = 240$  GHz, while the value predicted from Eq. (4.24), using  $\Delta_{\rm T} = 0.36$  meV found in the simulation, is  $\beta \simeq 214$  GHz. All the amplitudes fit the expected exponential behavior reasonably well until limited by the Runge-Kutta step-size, with the exception of small plateaus that appear at different transition times for different amplitudes. This indicates that the coefficient for the average infidelity contains some amplitude-dependent factors, but these factors are insignificant compared to the exponential scaling and unimportant for this proposal. The adiabatic constraint is easily satisfied by proceeding with our simulation using  $\Lambda \sim 0.06l$  and  $\tau = 100$  ps.

In order to ensure the  $r_{12}^x$  operation is performed as expected, the probabilities for the basis states are found using the initial state  $|000\rangle$  after the domain wall trajectory in Fig. 4.6 is simulated. As anticipated, the probabilities  $P_{000}(000)$  and  $P_{110}(000)$ oscillate, with negligible probabilities (on the order of our step-size limit of  $10^{-6}$ ) found in the states with incorrect total parity. However, the operation doesn't complete a full bit-flip for the duration expected to correspond to a  $\pi$  rotation, with ~ 0.3% of the probability found in the  $|011\rangle$  and  $|101\rangle$  instead of  $|110\rangle$ . This is consistent with a 0.3% shift of the rotation axis away from  $\hat{x}$  due to small, undesired overlaps between Majorana bound states not involved in the operation. For example, a small overlap between the bound states of the middle and right regions would lead to an additional rotation  $r_{23}^x$  that effectively shifts the axis of rotation very slightly.

Using an amplitude  $\Lambda = 0.06l$  for  $r_{12}^x$  and  $r_{23}^x$ , the simulated rotations have a period of ~ 0.2ns. Since the operations are achieved by bringing together exponentially decaying peaks, the overlap-dependent coupling between topological regions (e.g.,  $\Gamma$  in the Majorana Hamiltonian) is exponentially sensitive to the trajectory amplitude. Thus, longer rotation periods can be achieved by slightly decreasing the amplitude. On the other hand, greater amplitudes give shorter periods, but can also risk fusing the adjacent Majorana bound states if increased too much, which begins to occur in the simulation near  $\Lambda \sim 0.07l$ . Thus, the typical operation time (including adiabatic transitions) for the parameters is on the order of 0.4 ns.

Finally, the CHSH inequality is tested in the simulation by simultaneously performing the  $r_{12}^x$  and  $r_{23}^x$  rotations on the initial state  $|\Phi_{\rm E}^-\rangle$  and finding the probabilities for each basis state. The CHSH quantity is a function of three relative angles, making it more difficult to visualize and compare to our simulation. Instead, consider one of the planes involving the maximum violation, namely when  $R_2 = 0$  and  $L_1 = \pi/4$ [i.e.,  $A = \pi/4$  in Eq. (4.11)]. The theoretical prediction and simulation are plotted in Fig. 4.7, showing very good agreement and a significant range of angles that violates the inequality. Thus the simulation indicates that the more realistic semiconductor Hamiltonian is consistent with the simpler Majorana model and the proposal is feasible for demonstrating entanglement in a Majorana wire. The code for this simulation is provided in Appendix B.

Before discussing experimental considerations, consider a simplification to the CHSH experiment that only requires projective measurement and the repeated use of two operations, namely  $r_{12}^x(L)$  and  $r_{23}^x(R)$  with specific values for L and R. Ideally  $L = R = \pi/4$ , but they remain unspecified here with the thought that the experiment could be attempted with angles that differ slightly from the ideal case.

The experiment begins by tuning the chemical potential to create three topological regions and measuring all of their occupation parities to project to one of the eight basis states. For concreteness, consider only the states that lead to the  $|\Phi_{\rm E}^-\rangle$  state,



Figure 4.7: Left: Theoretical contour plot of the quantum mechanical prediction of the CHSH quantity for  $|\Phi_{\rm E}^-\rangle$  for the  $A = \pi/4$  plane. Local hidden variable theories are inconsistent with a CHSH quantity above 2, which occurs inside the white lines. Right: Simulated contour plot showing agreement with the global maximum at  $B = C = \pi/4$  present, violating the CHSH inequality by approximately 40%.

so the procedure only continues if the measurement of the middle parity matches the total parity [see Table 4.1]. Alternatively, the full experiment is carried out regardless of the measurement outcomes, but the cases when the parities do not match are disregarded. Then the operations  $r_{12}^x(L)$  and  $r_{23}^x(R)$  are simultaneously performed twice before measuring the middle parity, proceeding only when this parity matches the initial result. For the ideal angles  $L = R = \pi/4$ , this procedure prepares maximally entangled states.

This preparation is followed by one of the four rotation combinations given in Eq. (4.10), with  $L_1 = L$ ,  $L_2 = 3L$ ,  $R_1 = 0$ , and  $R_2 = 2R$ . For example, the combination with  $L_2$  and  $R_1$  is performed by carrying out  $r_{12}^x(L)$  three times but leaving the right domain walls stationary. After one of the rotation combinations is performed, all three parities are measured and the results are recorded. This is repeated several times for each combination to find the corresponding probabilities and calculate the CHSH quantity in Eq. (4.10). The quantum mechanical prediction for the CHSH quantity using the above procedure is easily calculated (though not particularly illuminating in written form) and is plotted in Fig. 4.8. As expected, the CHSH quantity has a maximum at  $L = R = \pi/4$ , with a wide range of angles spanning from approximately  $\pi/8$  to  $3\pi/8$  confirming Bell's theorem. Thus this procedure can be used for a broad range of angles, demonstrating entanglement in Majorana wires, even with limited accuracy in the tuning operations.



Figure 4.8: Contour plot of the quantum mechanical prediction of the CHSH quantity with  $L_1 = L$ ,  $L_2 = 3L$ ,  $R_1 = 0$  and  $R_2 = 2R$ . Local hidden variable theories are inconsistent with a CHSH quantity above 2, which occurs inside the white lines. The plot repeats with a period of  $\pi$  for both L and R.

#### 4.5 Experimental Considerations

This section discusses some aspects of this proposal that may be significant for an experimental realization. One of the first hurdles that must be overcome is the development of reliable projective measurement of the occupation. Aside from directly probing the wire with point contacts, there have been several proposals for observing the presence of Majorana bounds states such as using the Aharanov-Casher effect [50], transmons [49], and Shapiro step doubling in the AC-Josephson effect [99]. Without committing to a particular readout scheme, oe should note that this proposal requires measurement of a single topological region during the procedure in order to project to a maximally entangled state.

Another non-trivial experimental task is the fine-tuning of the chemical potential to minimize undesired Majorana peak overlaps. The simplest way to mitigate these overlaps is to use a longer wire, which exponentially reduces the overlaps. The simulation indicates that a wire length on the order of  $5 - 10\mu$ m is sufficient. Alternatively, a setup that links together several shorter wires may also be possible if longer wires are experimentally unavailable.

The adiabatic constraint that was found using the Landau-Zener formula is rather lenient, only requiring transition times on the order of 0.1ns. This is due to the generous topological gap of  $\sim 0.35$  meV that separated the zero-energy and bulk states, due to the relatively large proximity effect and g-factor. In addition, the topological phase requires a relatively large spin-orbit coupling. The parameters are reasonable when compared to experiments [82, 24, 20, 36, 99], but the need for a robust topological gap should be considered when selecting materials, and further advances of the proximity effect and spin-orbit in relevant materials would be helpful.

The operation time for performing the ideal rotation angles can be found experimentally by reproducing the probability plot in Fig. 4.6. Indeed, calibration is needed to find the effective coupling that accounts for additional effects such as disorder, local tunings, and the small coupling between regions via the s-wave superconductor [124]. For example, the ideal  $r_{12}^x(\pi/4)$  for the simplified CHSH experiment can be calibrated in the following way. First, all three parities are measured to project to a basis state, then gates are gradually tuned to shift the chemical potential in the left non-topological region for an operation time  $\sim 0.5$  ns, followed by a final measurement of all three parities. Note that the operations can be achieved by moving a single domain wall if this is easier experimentally; the simulation moved both the outer and middle domain walls to suppress overlaps in the topological regions, but this may be unnecessary in longer wires. This procedure is repeated several times with the same operation time, tracking the percentage of times the state remains in the initially measured state. This is repeated with several slightly different operation times, until the percentage is near  $\cos^2(\pi/8) \simeq 0.85$ . Once the rotations that correspond to  $L = R = \pi/4$  are roughly calibrated, the simplified CHSH experiment can be carried out.

As noted by Rainis et. al. [93], one must also consider the phenomena of quasiparticle poisoning in any system that uses superconductors to achieve the topological phase. While a superconductor at T = 0 (which our simulation assumed) will only form Cooper pairs, at finite temperatures less than ~ 160mK a small, fixed population of quasi-particles seems to be present [23]. Quasi-particle poisoning occurs when a single quasi-particle tunnels between the superconductor and semiconductor, changing the total parity of the system and destroying the quantum information. Thus, both the measurement and operation times must be much shorter than the average time of quasiparticle tunneling, constraining the operation time in the opposite limit as the adiabatic condition. Fortunately, estimates of the average time for quasi-particle tunneling in Majorana wire systems are 100ns or greater in typical experimental systems, depending on the specific parameters [93]. These calculations depend linearly on the resistance of the wire; thus longer wires will have shorter tunneling times. Nevertheless, with the adiabatic constraint only requiring operation times  $\sim 0.1$ ns, there is a wide window for the few operations needed in this proposal. Thus the constraint on the operation time of measurements is largely unaffected by the proposed operations.

Another non-trivial aspect inherent to the entanglement inequalities is the need to find probabilities rather than single measurements outcomes, requiring a high level of precision in the gate tuning. While this may make it difficult to reproduce the exact predictions of quantum mechanics, the large violation of the CHSH inequality by > 40%, and the wide range of angles that violate hidden variable theories may still be sufficient for demonstrating entanglement.

One way to circumvent the precision requirement is to perform the test proposed by Greenberger, Horne, Zeilinger [45] and Mermin [78] (GHZM). The GHZM experiment requires three logical qubits (thus four topological regions), but tests hidden variable theories with a single measurement, rather than involving probabilities and inequalities. Indeed, there are many interesting experiments that demonstrate entanglement, such as quantum teleportation, that are possible with one additional qubit.

This and other relevant systems are still a new and emerging area of research for theorists and experimentalists alike. With that said, there are many recent developments which are not reflected in the above Hamiltonian. For example, it seems that the experiments on the Majorana wire systems are not strictly one-dimensional, and must be analyzed as multi-channel wires to explain some of the experimental findings [110]. Others note that phenomena like Andreev reflection, disorder, finite temperature and the Kondo effect may need to be further understood in these systems [110, 71, 94]. Despite these complications, one should note that almost any convincing manifestation of Majorana bound states must demonstrate entanglement, which will likely be easier than full braiding. While this proposal discussed the specific setup of Majorana wires, the general idea of using non-topological, proximity-induced operations to test the entanglement inequalities as a stepping stone to braiding operations, as well as other aspects like separating logical qubits with a parity qubit, may be applied to many systems that potentially support Majorana bound states, such as topological insulators [39, 85, 40].

#### 4.6 General Preparation Procedure

Sec. 4.3 discussed the Bell and CHSH inequalities with the state  $|\Phi_{\rm E}^-\rangle$ . This section considers more general procedures for preparing any maximally entangled state and testing the inequalities.

Entanglement can be demonstrated for any initial condition with very simple alterations to the proposal, rather than discarding data for the incorrect initial state or measurement outcome.

The procedure shown in Fig. 4.2 prepares one of the eight maximally entangled states,

$$|\Phi_{\rm E}^{\pm}\rangle = \frac{|000\rangle \pm |101\rangle}{\sqrt{2}}, \qquad |\Psi_{\rm E}^{\pm}\rangle = \frac{|011\rangle \pm |110\rangle}{\sqrt{2}}, \qquad (4.25)$$

$$|\Phi_{\mathcal{O}}^{\pm}\rangle = \frac{|010\rangle \pm |111\rangle}{\sqrt{2}}, \qquad |\Psi_{\mathcal{O}}^{\pm}\rangle = \frac{|001\rangle \pm |100\rangle}{\sqrt{2}}, \qquad (4.26)$$

by measuring all three parities to project to one basis state from Eqs. (4.2), performing the operations  $r_{12}^x(\pi/2)$  and  $r_{23}^x(\pi/2)$ , then measuring the middle parity. The state that is prepared depends on the overall parity and middle parity measurements, as shown in Table 4.1. Note that the results for the even and odd total parity are equivalent upon

the exchange  $0 \leftrightarrow 1$  for the middle parity qubit.

Table 4.1: Maximally entangled state prepared for various total parity and middle parity measurements. The even and odd total parities give the same results if upon exchanging  $0 \leftrightarrow 1$  for the middle parity.

Total Parity	Even				Odd			
Initial Middle Parity	0		1		1		0	
Final Middle Parity	0	1	0	1	1	0	1	0
Resulting State	$\Phi_{\rm E}^-$	$\Psi_{\rm E}^+$	$\Phi_{\rm E}^+$	$\Psi_{\rm E}^-$	$\Phi_0^-$	$\Psi_{\rm O}^+$	$\Phi_{\rm O}^+$	$\Psi_{\rm O}^-$

Any of the maximally entangled states can be used to demonstrate the violation of the Bell and CHSH inequalities, but with different rotation angles. For example, quantum mechanics predicts that  $\langle L, R \rangle$ , the expectation value of the combined parity of the left and right qubits after being rotated by angles L and R, respectively, for  $\Phi_{\rm E}^+$ is  $\cos(L+R)$  rather than  $\cos(L-R)$  for  $\Phi_{\rm E}^-$ . Clearly the CHSH quantity in Eq. (4.10) is the same except with  $R \to -R$ , which can be returned to the case in Sec. 4.3 by substituting  $\{L_1, L_2, R_1, R_2\} \to \{L_1, L_2, -R_1, -R_2\}$ . The relevant probabilities and angle transformations for the Bell and CHSH inequalities are listed in Table 4.2 for each maximally entangled state. These changes can be accounted for by designing the experiment to perform different rotations depending on the middle parity measurement outcomes found during the preparation of the maximally entangled states.

Table 4.2: Probabilities and expectation values predicted by quantum mechanics for the various maximally entangled states. The set of angles that corresponds to the case in the body of the paper for the CHSH violation is given as well. The results are the same for even and odd parity, so we suppress the corresponding subscript.

State	$P_{=}(L,R)$	$\langle L, R \rangle$	CHSH Angles		
$\Phi^-$	$\cos^2\left(\frac{L-R}{2}\right)$	$\cos(L-R)$	$L_1, L_2, R_1, R_2$		
$\Phi^+$	$\cos^2\left(\frac{L+R}{2}\right)$	$\cos(L+R)$	$L_1, L_2, -R_1, -R_2$		
$\Psi^-$	$\sin^2\left(\frac{L-R}{2}\right)$	$-\cos(L-R)$	$L_1, L_2, R_1, R_2$		
$\Psi^+$	$\sin^2\left(\frac{L+R}{2}\right)$	$-\cos(L+R)$	$L_1, L_2, -R_1, -R_2$		

#### 4.7 Majorana Wire Proposal Summary

The motion of domain walls in a Majorana wire system was analyzed with three separate topological regions using a simple Hamiltonian analogous to Kitaev's toy model [61]. Using the occupation number parity of the Majorana bound states in each region, a three qubit basis was defined with two logical qubits and one qubit forfeited to total parity conservation. In this basis, x-axis rotations are performed by extending the outer topological regions to isolate a single coupling between different topological regions. While these rotations are not topologically protected from local perturbations, they can demonstrate entanglement by testing the Bell and CHSH inequalities.

With the simpler model as a guide, the domain wall motion was simulated using a more realistic semiconductor Hamiltonian. Results indicate that the topological regions can be well separated in wires of length  $\sim 5 - 10 \,\mu$ m with reasonable parameters compared to recent experiments. Adiabatic changes in the chemical potential were simulated with operation times on the order of 0.5ns, consistent with the Landau-Zener condition applied to excitations from zero-energy to the bulk. Extending the topological regions results in the rotations predicted by the simpler model. Finally, the CHSH experiment was simulated and the expected inconsistency with hidden variable theories predicted by Bell's theorem was found, indicating that the simpler Majorana Hamiltonian approximates the Majorana wire system well.

A simplified version of the CHSH experiment was introduced that only requires projective measurement and repeated use of two  $\pi/4$  rotations, finding a wide range of operation angles from  $\pi/8$  to  $3\pi/8$  violate the entanglement inequalities. Thus a keyboard gating setup needs to be relatively precise, but moderate inaccuracy is tolerable. Calibrating methods were provided for the rotations and potential hurdles for an experimental realization were discussed. The analysis and simulation indicate that there is a large window of operation times, spanning three orders of magnitude, that satisfy the adiabatic and quasi-particle poisoning constraints. Hence this proposal is viable for demonstrating entanglement in Majorana wires if methods for projective measurement and precise gate tuning are available.

While there have been proposals to test entanglement [57, 17, 100, 6, 67], observing entanglement by directly testing Bell's theorem would be a novel opportunity in solid state systems in general, and an important advance for Majorana wire systems specifically. In addition to being a crucial ingredient for quantum information, entanglement may rule out alternative explanations inconsistent with non-local effects, providing a useful tool beyond local measurements. More so, experimental work on gate tuning is already required for braiding operations, and this proposal could serve as a useful benchmark towards that goal. Thus, the observation of entanglement would support current models of Majorana wires and provide a significant piece of evidence supporting the presence of Majorana bound states.

## Chapter 5

# Conclusion

This dissertation considered two proposals for fault-tolerant qubits: pairs of electron spins in semiconductor double quantum dots and Majorana zero modes bound to the ends of semiconductor nanowires coupled to conventional superconductors. Analysis and simulation of these systems showed that both proposals are viable and could enable further advances in the construction of quantum information systems.

In particular, the additional spin-orbit effects that result from the motion of double quantum dots do not outweigh the benefits of suppressing the hyperfine dephasing. Thus it is possible to incorporate adiabatic motion into quantum dot proposals, providing another avenue for decreasing the infidelity of quantum operations closer to the quantum error correction threshold. Similarly, the adiabatic motion of Majorana bound states in nanowires can be used to perform quantum operations within the constraints of quasiparticle poisoning. Aside from the potential for braiding operations, these operations are important since they allow one to both setup and demonstrate entanglement, a crucial ingredient to any quantum information system. Rather than depending solely on experimental advances, these proposals aim to prevent decoherence by building on qubit schemes that are fundamentally immune to certain local errors. Furthermore, these proposals are constructed from conventional materials and are within reach of current experiments. While these fault-tolerant designs require additional complications such as adiabatic motion and novel particles, the included analysis and simulations indicate that these added complexities are not prohibitive. In summary, the use of adiabatic motion to implement and operate on qubits is a viable method for getting closer to the ambitious goal of realizing a quantum computer.
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#### Appendix A

# **Double Quantum Dots Simulation**

### Code

```
1 /*
 2 * RK4 program written by David Drummond
    */
3
4
5 #include <iostream>
6 #include <cstdlib>
7 #include <string>
8 #include <fstream>
9 #include <complex>
10 #include <Eigen/Dense>
11 #include "grev_rnd.h"
12
13 #define PI 3.1415926535897932
14
15 using namespace std;
16 using namespace Eigen;
17
18 //global variables
{\rm 19} \ {\tt Matrix2cd} \ {\tt U1} \ , \ {\tt U2} \ , \ {\tt B0} \ , \ {\tt I} \ , \ {\tt sx} \ , \ {\tt sy} \ , \ {\tt sz} \ ,
            \texttt{H1}\;,\;\;\texttt{H2}\;,\;\;\texttt{k1}\;,\;\;\texttt{k2}\;,\;\;\texttt{k3}\;,\;\;\texttt{k4}\;;
20
21 double chi=0.5; //approximate percentage of time spent moving
22 double T=400.0; //total time to complete 1 circle, including rest time
23 double tau=T*chi/2; //time spent to make one semi-circle
24 int steps=400*3200; //good convergence at 3200
25 int tt = 0;
26 double TT = 40000.0;
27 double stepSize=T/steps;
28 \text{ double beta} = 0.075;
29 double b=2*PI/*+0.001*/;
30 double t1 = 0.0;
31 double t2=TT/2;
```

```
32 double t3=TT;
33
34 //prototypes
35 void plotOriginal();
36 double angle(double t, int dir, int dot);
37 double velx(double t, int dir, int dot);
38 double vely(double t, int dir, int dot);
39 double nonStatic(double t, int dir, int dot, int field);
40 Matrix2cd Ham(double t, int dir, int dot);
41 void plotAngle(int dir, int reps, int dot);
42 void plotVel(int dir, int reps, int dot);
43 void plotSpli();
44 void plotField(int dir, int reps, int dot);
45 void plotPhi(int dir, int reps);
46 void display(int dir, int reps);
47 double infid(Matrix2cd & u1, Matrix2cd & u2);
48 Matrix4cf kronProd(Matrix2cd & A, Matrix2cd & B);
49 void initMatrices();
50 void rk4(int dir, int reps);
51 void repeatSim(int cycles, int dir, int reps);
52
       //plot original field from do_field
53
       void plotOriginal()
54
       {
55
                  cout<<endl<<endl;</pre>
56
                  cout<< "# Plotting original field "<<endl;</pre>
57
                  for (int i=0; i<256*2;i++)
58
                           cout << i*0.5 <<"\t"<<do_field(i*0.5,0)
59
                                     <<" \ t" << do_field(i * 0.5, 1)
60
                                     <<"\t"<<do_field(i*0.5,2)
61
                                     <<"\t"<<do_field(i*0.5,3)
62
                                     <<"\t"<<do_field(i*0.5,4)
63
                                     <<"\t"<<do_field(i*0.5,5)
64
                                     <<"\t"<<do_field(i*0.5,6)
65
                                     <<"\t"<<do_field(i*0.5,7)
66
                                     <<"\t"<<do_field(i*0.5,8)
67
                           <<endl;
68
                  cout<<endl<<endl;</pre>
69
70
71
        //returns number between 0 and 256 representing angle between 0 and 2*pi
72
         //for dot going in direction dir - allows t>T using fmod function
73
74 double angle(double t, int dir, int dot)
75
                  return 64*(2 + \text{dir}* \tanh(5*(\text{fmod}(t+\text{dot}*T/2,T)-T/4)/\text{tau})
76
                                               + dir*tanh(5*(fmod(t+dot*T/2,T)-3*T/4)/tau));
77
78 }
79
so double velx(double t, int dir, int dot)
81 {
                  \texttt{return} \; \texttt{dir} * \texttt{sin} \left( 2 * \texttt{PI} * \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \right) / \texttt{T} \right) * \left( \texttt{pow} \left( \texttt{cosh} \left( 5 * \left( \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) \right) \right) \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} / 2 \,, \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \leftrightarrow \texttt{T} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{T} \rightarrow \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{dot} * \texttt{T} , \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{fmod} , \texttt{fmod} \right) + \texttt{fmod} \left( \texttt{t} + \texttt{fmod} \right) + \texttt{fmo
82
                             )-T/4)/tau), -2)
                           + pow(cosh(5*(fmod(t+dot*T/2,T)-3*T/4)/tau),-2));
83
84
85
86 double vely(double t, int dir, int dot)
87
       ł
                  return 3.5*dir*cos(2*PI*fmod(t+dot*T/2,T)/T)*(pow(cosh(5*(fmod(t+dot*T\leftrightarrow T)))))
88
                             (2,T)-T/4)/tau), -2)
                           + pow(cosh(5*(fmod(t+dot*T/2,T)-3*T/4)/tau),-2));
89
```

```
90 }
91
92 inline
93 double nonStatic(double t, int dir, int dot, int field)
94
   {
       /*return (t/TT)*do_field(angle(t, dir, dot), field)
95
      +((TT-t)/TT)*do_field(angle(t, dir, dot), field+3);
96
97
       */
          return ((t-t2)*(t-t3)/((t1-t2)*(t1-t3)))*do_field(angle(t, dir, \leftrightarrow
98
              dot), field)
       + ((t-t1)*(t-t3)/((t2-t1)*(t2-t3)))*do_field(angle(t, dir, dot), \leftrightarrow
99
            \texttt{field}+3)
       + ((t-t2)*(t-t1)/((t3-t2)*(t3-t1)))*do_field(angle(t, dir, dot), \leftrightarrow
100
            \texttt{field}+6;
101
102
103
104
105
   inline
106 Matrix2cd Ham(double t, int dir, int dot)
107
   /* return B0
108
      + sx*do_field(angle(t, dir, dot), 0)
109
      + sy*do_field(angle(t, dir, dot),1);
110
      + sz*do_field(angle(t, dir, dot), 2);
111
      // + sx*beta*(-1.0*velx(t, dir, dot)+vely(t, dir, dot))
112
      // + sy*beta*(-1.0*velx(t, dir, dot)+vely(t, dir, dot));
113
114
          return BO
115
         + sx*beta*(-1.0*velx(t, dir, dot)+vely(t, dir, dot))
116
117
         + sy*beta*(-1.0*velx(t, dir, dot)+vely(t, dir, dot))
118
         + sx*nonStatic(t, dir, dot, 0)
         + sy*nonStatic(t, dir, dot,1)
119
         + sz*nonStatic(t, dir, dot, 2);
120
121
122
123
   void plotAngle(int dir, int reps, int dot)
124
125
       cout<<endl<<endl;</pre>
126
       cout <<"# Plotting angle for dot "<<dot+1<<" with direction "<<dir<<↔
127
           endl:
       tt=0:
128
       int tempDir=1; //used to reverse direction if dir=-1
129
       for (int j=0; j<reps; j++)
130
131
       {
          for (int i=0; i < steps; i++)
132
133
          {
             tt++:
134
             cout \ll tt*stepSize \ll 't' \ll angle(tt*stepSize, tempDir, dot) \ll endl \leftrightarrow
135
136
          tempDir*=dir; //changes direction if dir=-1
137
138
       ł
139
   ł
140
141 void plotVel(int dir, int reps, int dot)
142
   ł
      cout<<endl;</pre>
143
```

```
cout <<"# Plotting velocity for "<<dot+1<<" with direction "<<dot+1↔
144
       tt=0;
145
       int tempDir=1; //used to reverse direction if dir=-1
146
147
       for (int j=0; j < reps; j++)
       {
148
          for (int i=0; i < steps; i++)
149
150
          {
              tt++;
151
              cout<< tt*stepSize<<"\t"<<velx(tt*stepSize, tempDir, dot)</pre>
152
                           <<"\t"<<vely(tt*stepSize, tempDir, dot)<<endl;</pre>
153
154
          tempDir*=dir; //changes direction if dir=-1
155
156
157
158
159
   void plotSpli()
160
161
       cout<<endl<<endl;</pre>
162
       for (int t=0;t<TT;t++)
          cout << t << ' t' <<( (t-t2)*(t-t3)/((t1-t2)*(t1-t3)) )
163
             << ' \ t <<( (t-t1)*(t-t3)/((t2-t1)*(t2-t3)))
164
          <\!\!< ' \ t' <\!\!<\!\!( \ (t-t2) * (t-t1) / ((t3-t2) * (t3-t1)) ) ) <\!\!<\!\!endl;
165
   }
166
167
   void plotField(int dir, int reps, int dot)
168
169
   {
170
       cout<<endl<<endl;</pre>
       cout <<"# Plotting angle for dot "<<dot+1<<" with direction "<<dir<<↔
171
           endl;
172
       tt=0:
       int tempDir=1; //used to reverse direction if dir=-1
173
174
       for (int j=0; j<reps; j++)
175
       ł
          for (int i=0; i < steps; i++)
176
177
          {
              tt++;
178
              //cout << tt * stepSize <<"\t"<< do_field (angle (tt * stepSize , tempDir , ↔
179
                   dot),2)<<endl;
              cout \ll tt*stepSize \ll t'* < nonStatic(tt*stepSize,tempDir,dot,2) < \leftrightarrow
180
                  endl;
          }
181
          tempDir*=dir; //changes direction if dir=-1
182
       }
183
   3
184
185
   void plotPhi(int dir, int reps)
186
187
   ł
       cout<<endl<<endl;</pre>
188
       cout << "# Plotting phi with direction "<<dir<<endl;
189
190
       tt=0;
       double phi1=0.0;
191
       double phi2=0.0;
192
       int tempDir=1; //used to reverse direction if dir=-1
193
       for (int j=0; j<reps; j++)
194
195
       ł
          for (int i=0; i < steps; i++)
196
197
          ł
198
              tt++;
              phi1+=do_field(angle(tt*stepSize,tempDir, 0),2);
199
```

```
phi2+=do_field(angle(tt*stepSize,tempDir, 1),2);
200
             cout << tt * stepSize << "\ t "<< phi1<< "\ t "<< phi2<< endl;
201
          }
202
203
          tempDir*=dir; //changes direction if dir=-1
204
       }
205
   }
206
   void display(int dir, int reps)
207
208
   ł
       plotOriginal();
209
       plotAngle(dir, reps, 0);
210
       plotAngle(dir, reps, 1);
211
       plotField(dir, reps, 0);
212
       plotField(dir, reps, 1);
213
       plotSpli();
214
215
       plotVel(dir, reps, 1);
216
      plotPhi(dir, reps);
217
218
   //initialize matrices before rk4 iterations
219
   void initMatrices()
220
   {
221
      U1 << 1.0, 0.0, 0.0, 1.0; //Initialize U to 1 for dot 1
222
      U2 \ll 1.0, 0.0, 0.0, 1.0; //Initialize U to 1 for dot 2
223
      I(0,0)=complex < double > (0.0, -stepSize); // Eigen doesn't like to mix \leftrightarrow
224
           complex numbers with matrices
      I(1,1) = complex < double > (0.0, -stepSize); //so an i=sqrt(-1) matrix is \leftrightarrow
225
                                          //also include -step from the \leftrightarrow
226
                                              differential equation for efficiency
      sx \ll 0.0, 1.0, 1.0, 0.0;
227
      sz \ll 1.0, 0.0, 0.0, -1.0;
228
      BO \ll b, 0.0, 0.0, -b; //Applied B field
229
       sy(0,1) = complex < double > (0.0, -1.0);
230
       sy(1,0) = complex < double > (0.0,1.0);
231
232
233
    //returns the kronecker product of 2x2 matrices A and B
234
   Matrix4cf kronProd(Matrix2cd & A, Matrix2cd & B)
235
236
237
      Matrix4cf K;
       for(int i=0; i<4; i++)</pre>
238
239
          for (int j=0; j<4; j++)
240
             K(i,j) = A(floor(i/2), floor(j/2)) * B(i\%2, j\%2);
241
242
       return K;
243
244 }
245
   //outputs the average infidelity of u1 and u2
246
247 double infid(Matrix2cd & u1, Matrix2cd & u2)
248
   {
      Matrix4cf U=kronProd(u1,u2); //forms kronecker product of the \leftrightarrow
249
      U(0,0)=0; U(3,0)=0; U(0,3)=0; U(3,3)=0; //project onto logical \leftrightarrow
250
           subspace
       Matrix4cf UUdag=U*U.adjoint();
251
       return (1-real(abs(U.trace()))*abs(U.trace()) + UUdag.trace())/6);
252
253 }
254
```

```
255 //performs RK4
   void rk4(int dir, int reps)
256
257
   {
258
       cout<<endl<<endl;</pre>
259
       tt=0:
       int tempDir=1; //used to reverse direction if dir=-1
260
       H1 = Ham(0, tempDir, 0); //initial H value
261
       H2 = Ham(0, tempDir, 1);
262
263
       for (int j=0; j<reps; j++)
264
       {
           for (int i=0; i < steps; i++)
265
266
           {
               k1 = I*H1*U1; //evaluate k1 using the previous H value
267
               \texttt{H1} = \texttt{Ham}((\texttt{tt+tempDir}*0.5)*\texttt{stepSize}, \texttt{tempDir}, 0); //\texttt{update} \text{ for } \leftarrow
268
                   half step
               k2 = I * H1 * (U1 + 0.5 * k1);
269
270
               k3 = I * H1 * (U1 + 0.5 * k2);
               H1 = Ham((tt+tempDir)*stepSize, tempDir, 0); //update H for next \leftrightarrow
271
               k4 = I * H1 * (U1+k3);
272
               U1 + = (1/6.0) * (k1 + 2 * k2 + 2 * k3 + k4);
273
274
               k1 = I*H2*U2; //same for dot 2
275
               H2 = Ham((tt+tempDir*0.5)*stepSize, tempDir, 1);
276
               k2 = I * H2 * (U2 + 0.5 * k1);
277
               k3 = I * H2 * (U2 + 0.5 * k2);
278
               H2 = Ham((tt+tempDir)*stepSize, tempDir, 1);
279
               k4 = I * H2 * (U2+k3);
280
               U2 + = (1/6.0) * (k1 + 2 * k2 + 2 * k3 + k4);
281
282
283
               tt++:
               if ((i < 16000) || (i > (steps - 16000))) // if 16000 steps within \leftrightarrow
284
                   commensurate time
                   { if (tt%10==1)
285
                      printf("%10.9g t %10.9g n",tt*stepSize, infid(U1,U2));
286
287
               else
288
                   if (tt%3200==1)
289
                      printf("\%10.9g \setminus t \%10.9g \setminus n",tt*stepSize, infid(U1,U2));
290
291
           tempDir*=dir; //changes direction if dir=-1
292
       }
293
294
   J
295
    //repeats sim cycles times, each with different steps, or TT
296
   void repeatSim(int cycles, int dir, int reps)
297
298
   {
       for(int i=0;i<cycles;i++)</pre>
299
300
       {
           cout << "#TT="<< TT << endl;
301
           rk4(dir, reps);
302
           steps *=2;
303
           TT - = 2000;
304
           t2 = TT / 2;
305
           t3=TT;
306
307
           tt=0:
           stepSize=T/steps;
308
309
           initMatrices();
310
       }
311 }
```

```
312
313 int main (int argc, const char *argv[])
314 {
      initMatrices();
315
316
      do_init(argc,argv);
      int reps=40;
317
      int dir=-1;
318
      cout<<"# Running rk4 with step size: "<<stepSize<<endl</pre>
319
         <<"# "<<reps<<" Reps with direction="<<dir<<" with TT="<<TT<<endl;</pre>
320
      rk4(dir, reps);
321
322 // repeatSim(5, dir, reps);
323
      kill_spline();
324
      return 0;
325 }
```

#### Appendix B

### Majorana Wire Simulation Code

```
/* Author - David Drummond
1
   * Majorana program to simulate Hamiltonian of semiconductor wire
2
   * Note that unit tests are commented out, rather than deleted
3
4
   * verbosity key
\mathbf{5}
   * v=0 nothing
6
7
   * v=1 parameter comment
   * v=2 matrices comment
8
   * v=4 domain walls
9
   * v=8 chemical potential - animation compatible if TT!=0
10
   * v=16 eigenvalues for zero-energy and lowest bulk state
11
12
   * v=32 Zero-energy spatial distribution - animation compatible if TT!=0
   * v=64 non-adiabatic leakage
13
14 */
15 #include "c-major.h"
                          Vy Vz delta mu alpha lambda AL AR durL ↔
L1 L2 L3 orig1 orig2 orig3 tSteps TT dispRate↔
16
             t
                     Vx
      durR TauL TauR L1
       nx ny nz step per mN v */
17
18 params prm = \{11.3, 0.0, 0.0, 0.9, 0.5, 2.12, 2, 0.06, 0.11, 0.115, 496.6, \leftrightarrow \}
       442, 200, 200, 0.13, 0.09, 0.1355, 0.0, 0.5, 1.0, 100, 800, 4, \leftrightarrow
         400, 0, 0, 3, 3, 56;
19 int NMAX=1000; /* matrix is initialized with size 4*NMAX, but reduced to \leftrightarrow
      actual size
      * after prm are initialized 4nn by 4nn hamiltonian of dynamic size X \leftrightarrow
20
         by X of
      * complex double coefficients initially full of zeros. 4nn for \leftrightarrow
21
         creation (c*) and
      * annihilation (c), as well as spin,
22
      * operators basis is clu, cld, c2u, c2d,...c(nn)u, c(nn)d, c*lu, c*ld,\leftrightarrow
23
           c*2u, c*2d, ..., c*(nn)u, c*(nn)d
     */
^{24}
_{25} MatrixXcd ham, evec, maj, Usub, Iden, Uinit, Td, projMat1L, projMat2L, \leftrightarrow
      projMat3L, projMat1R, projMat2R, projMat3R;
26 MatrixXcd addMat1, addMat2, addMat3, /*addMat1R, addMat2R, addMat3R,*/ \leftrightarrow
      locL , locR , symL , symR , tempMaj;
27 MatrixXcd maj1L, maj2L, maj3L, maj1R, maj2R, maj3R, fixMat1L, fixMat2L, \leftrightarrow
      fixMat3L, fixMat1R, fixMat2R, fixMat3R;
28 VectorXd chem, Vx, Vy, Vz, eval;
29 VectorXcd del, initState, prob;
30 SelfAdjointEigenSolver<MatrixXcd> sol; //workspace to solve eigen \leftrightarrow
     equations
```

```
31 std::vector<br/>T> tripletList; //vector to hold triplets of row, col,<br/> \leftarrow
        elements for sparse Matrix
32 SparseMatrix<complex<double>> sparseHam;
33 double latSize, stepSize, startL, startR;
34 int nn, mN, size;
35 IOFormat comMat(7, 0, "", " \setminus n", "\# [", "]");
       //precision, alignRows, coeffSep, rowSep, rowPrefix, rowSuffix, \leftarrow
36
            matPrefix, matSuffix
37
   void init_params(int argc, char *argv[]){ /* read cmd line params */
38
         int i, ok=1;
39
         for (i=1; i < argc; ++i) {
40
              if(
41
                sscanf(argv[i],"t=%lg",&prm.t)
                                                            42
               sscanf (argv [i], "Vx=%lg",&prm.Vx) ||

sscanf (argv [i], "Vy=%lg",&prm.Vx) ||

sscanf (argv [i], "Vy=%lg",&prm.Vy) ||

sscanf (argv [i], "delta=%lg",&prm.delta)

sacanf (argv [i], "endetta=%lg",&prm.delta)
43
44
45
46
                                                                        sscanf(argv[i], "mu=%lg",&prm.ucluar) ||
sscanf(argv[i], "mu=%lg",&prm.u)||
sscanf(argv[i], "alpha=%lg",&prm.alpha)||
sscanf(argv[i], "lambda=%lg",&prm.lambda)||
47
48
49
                sscanf(argv[i],"AL=%lg",&prm.AL)|
50
                sscanf(argv[i], "AR=%lg",&prm.AR)||
51
                sscanf(argv[i],"durL=%lg",&prm.durL)||
52
                sscanf(argv[i],"durR=%lg",&prm.durR)
53
                sscanf(argv[i],"TauL=%lg",&prm.TauL)||
54
                sscanf(argv[i],"TauR=%lg",&prm.TauR)||
55
                sscanf(argv[i],"L1=%lg",&prm.L1)|
56
                sscanf(argv[i],"L2=%lg",&prm.L2)||
57
                sscanf(argv[i], "L3=\%lg", &prm.L3)
58
                sscanf(argv[i], "orig1=%lg",&prm.orig1)||
59
                sscanf(argv[i], "orig2=%lg",&prm.orig2)||
60
                sscanf(argv[i],"orig3=%lg",&prm.orig3)||
61
                sscanf(argv[i],"tSteps=%d",&prm.tSteps)||
62
                sscanf(argv[i],"TT=%d",&prm.TT)||
63
               sscanf(argv[i], "Ir /nk ,&prm.Ir)||
sscanf(argv[i], "dispRate=%d",&prm.dispRate)||
sscanf(argv[i], "nx=%d",&prm.nx)||
sscanf(argv[i], "ny=%d",&prm.ny)||
sscanf(argv[i], "nz=%d",&prm.nz)||
64
65
66
67
                sscanf(argv[i],"step=%d",&prm.step)||
68
               sscanf(argv[i], "per=%d",&prm.per)||
sscanf(argv[i], "mN=%d",&prm.mN)||
69
70
                sscanf(argv[i], "v=%d",&prm.v)
71
72
73
74
                if (prm.v&1)
                    printf("# main: scanned arg[%d]=\"%s\"\n",i,argv[i]);
75
76
               ok++;
77
           }
78
       if ((prm.nx+prm.ny+prm.nz)>NMAX)
79
           cout <<" ERROR: Number of lattice sites is larger than NMAX" << endl;
80
         if(ok !=argc)
81
           cout << " invalid input " << endl
82
               <<"<br/>USAGE: [t=#] [Vx=#] [Vy=#] [Vz=#] [delta=#] [mu=#] [alpha=#] \leftrightarrow
83
                     [lambda=\#] [AL=\#] [AR=\#] [durL=\#] [durR=\#] [TauL=\#] [TauR=\#] \leftrightarrow
                       [L1=\#] \ [L2=\#] \ [L3=\#] \ [orig1=\#] \ [orig2=\#] \ [orig3=\#] \ [tSteps \leftrightarrow
                    =#] [TT=#] [dispRate=#] [n?=#] [per=#] [mN=#] [step=#] [v=#]↔
               <<endl<<"where ?=x,y,z"<<endl;
```

```
84
```

```
111
```

```
if (prm.v&1)
 85
 86
           printf("# params: t=%g Vx=%g Vy=%g Vz=%g delta=%g mu=%g alpha=%g ↔
 87
               lambda=%g AL=%g AR=%g durL=%g durR=%g TauL=%g L1=%g L2↔
               =%g L3=%g orig1=%g orig2=%g orig3=%g tSteps=%d TT=%d dispRate=%↔
               d nx=%d ny=%d nz=%d step=%d per=%d mN=%d v=%d\n",
              prm.t, prm.Vx, prm.Vy, prm.Vz, prm.delta, prm.mu, prm.alpha, prm↔
 88
                   .lambda, prm.AL, prm.AR, prm.durL, prm.durR, prm.TauL, prm.\leftrightarrow
                   TauR, prm.L1, prm.L2, prm.L3, prm.orig1, prm.orig2, prm. \leftrightarrow
                   \tt orig3 \ , \ \tt prm.tSteps \ , \ \tt prm.TT \ , \ \tt prm.dispRate \ , \ \tt prm.nx \ , \ \tt prm.ny \ , \ \tt prm \ \leftrightarrow
                   .nz, prm.step, prm.per, prm.mN, prm.v);
           cout \ll \# verbosity key: v=0 nothing, v=1 parameter comment, v=2 \leftrightarrow
 89
               =16 eigenvalues, v=32 Zero-energy spatial distribution, v=64 \leftrightarrow
               non-adiabatic leakage"<<endl;</pre>
 90
       }
 91
       nn=prm.nx+prm.ny+prm.nz;
 92
       mN=prm.mN;
 93
       size=powInt(2,mN);
       latSize = 1.0/nn;
 ^{94}
       stepSize = 1.0 / prm.tSteps;
 95
       chem = VectorXd::Zero(nn);
 96
       del = VectorXcd::Zero(nn);
 97
       Vx = VectorXd :: Zero(nn);
 98
       Vv = VectorXd::Zero(nn);
99
       Vz = VectorXd :: Zero(nn);
100
       prob = VectorXcd::Zero(size);
101
       initState = VectorXcd::Zero(size);
102
       initState(0) = 0.5 * sqrt(2);
103
       \texttt{initState}(0) = 1;
104
       \texttt{initState}(1) = 0;
105
       \texttt{initState}(2) = 0;
106
       \texttt{initState}(3) = 0;
107
       initState(3) = -0.5 * sqrt(2);
108
       \texttt{initState}(4) = 0;
109
       \texttt{initState}(5) = 0;
110
       \texttt{initState}(6) = 0;
111
       initState(7) = 0;
112
       ham = MatrixXcd:: Zero(4*nn, 4*nn);
113
       maj = MatrixXcd::Zero(4*nn,mN);
114
       maj1L = MatrixXcd::Zero(4*nn,1); maj2L = MatrixXcd::Zero(4*nn,1); \leftrightarrow
115
           maj3L = MatrixXcd:: Zero(4*nn, 1);
       maj1R = MatrixXcd:: Zero(4*nn, 1); maj2R = MatrixXcd:: Zero(4*nn, 1); \leftrightarrow
116
           maj3R = MatrixXcd::Zero(4*nn,1);
       tempMaj = MatrixXcd::Zero(4*nn,1);
117
       locL = MatrixXcd:: Zero(4*nn, mN); locR = MatrixXcd:: Zero(4*nn, mN);
118
       symL = MatrixXcd::Zero(4*nn,mN); symR = MatrixXcd::Zero(4*nn,mN);
119
120
    //#include "savedMaj400.dat"
121
       projMat1L = MatrixXcd::Zero(4*nn, 4*nn); projMat2L = MatrixXcd::Zero(4* \leftrightarrow
122
           nn, 4*nn; projMat3L = MatrixXcd::Zero(4*nn, 4*nn);
123
       projMat1R = MatrixXcd::Zero(4*nn,4*nn); projMat2R = MatrixXcd::Zero(4*\leftrightarrow
           nn, 4*nn; projMat3R = MatrixXcd::Zero(4*nn, 4*nn);
       \texttt{fixMat1L} = \texttt{MatrixXcd}::\texttt{Zero}(4*\texttt{nn}, 4*\texttt{nn}); \texttt{fixMat2L} = \texttt{MatrixXcd}::\texttt{Zero}(4*\texttt{nn} \leftrightarrow \texttt{nn})
124
            ,4*nn); fixMat3L = MatrixXcd::Zero(4*nn, 4*nn);
       \texttt{fixMat1R} = \texttt{MatrixXcd}::\texttt{Zero}(4*nn, 4*nn); \texttt{fixMat2R} = \texttt{MatrixXcd}::\texttt{Zero}(4*nn \leftrightarrow \texttt{fixMat2R})
125
            ,4*nn); fixMat3R = MatrixXcd::Zero(4*nn, 4*nn);
126
       int cut = nn * 0.2;
       for (int i=0; i<nn; i++){
127
           if (i <= cut) {
128
```

```
projMat1L(2*i, 2*i) = 1;
129
             projMat1L(2*i + 1, 2*i + 1) = 1;
130
             projMat1L(2*i + 2*nn, 2*i + 2*nn) = 1;
131
             projMatlL(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
132
133
          }
          if(i > cut \&\& i \le 2*cut + 0.2*nn)
134
             projMat2L(2*i, 2*i) = 1;
135
             projMat2L(2*i + 1, 2*i + 1) = 1;
136
             projMat2L(2*i + 2*nn, 2*i + 2*nn) = 1;
137
             projMat2L(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
138
139
          if(i > 2*cut + 0.2*nn){
140
             projMat3L(2*i, 2*i) = 1;
141
             projMat3L(2*i + 1, 2*i + 1) = 1;
142
             projMat3L(2*i + 2*nn, 2*i + 2*nn) = 1;
143
             projMat3L(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
144
145
146
       }
147
       \operatorname{cut} = \operatorname{nn} * 0.4;
       for (int i=0; i<nn; i++){
148
          if(i \leq cut)
149
             projMat1R(2*i, 2*i) = 1;
150
             projMat1R(2*i + 1, 2*i + 1) = 1;
151
             projMat1R(2*i + 2*nn, 2*i + 2*nn) = 1;
152
             projMatlR(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
153
          }
154
          if(i > cut \&\& i \le 2*cut)
155
             projMat2R(2*i, 2*i) = 1;
156
             projMat2R(2*i + 1, 2*i + 1) = 1;
157
             projMat2R(2*i + 2*nn, 2*i + 2*nn) = 1;
158
159
             projMat2R(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
160
          if(i > 2*cut){
161
             projMat3R(2*i, 2*i) = 1;
162
             projMat3R(2*i + 1, 2*i + 1) = 1;
163
             projMat3R(2*i + 2*nn, 2*i + 2*nn) = 1;
164
             projMat3R(2*i + 2*nn + 1, 2*i + 2*nn + 1) = 1;
165
166
167
       addMat1 = MatrixXcd::Zero(mN, mN); addMat2 = MatrixXcd::Zero(mN, mN); \leftrightarrow
168
            addMat3 = MatrixXcd :: Zero(mN, mN);
       for (int i=0; i < mN; i++)
169
          addMat1(i, 0) = 1;
170
          addMat2(i,1) = 1;
171
          addMat3(i,2) = 1;
172
173
       }
174
      Usub = MatrixXcd :: Zero(4*nn, 2*mN);
175
      Uinit = MatrixXcd :: Zero(4*nn, 2*mN);
176
      Td = MatrixXcd :: Identity (2*mN, 2*mN);
177
       Iden = MatrixXcd :: Identity (2*mN, 2*mN);
178
179
       sparseHam.resize(4*nn,4*nn);
       tripletList.reserve(28*nn-16*prm.per);
180
       /* reserves space for the number of non-zero elements in sparse Ham nn↔
181
           on-site h,
        * nn on-site V, nn on-site Delta...*2 for spin plus nn-1 t, nn-1 \hookleftarrow
182
            alpha...*2 for
        * left and right movers and *2 for spin gives 6*nn+8*(nn-1)=14*nn\leftrightarrow
183
            -8...*2 for Nambu
```

```
* "symmetry" = 28*nn-16 if periodic all movers have nn terms instead \leftrightarrow
184
             of nn-1
        */
185
       prm.TauL=5000*max(prm.AL,prm.AR);
186
187
       prm.TauR=prm.TauL;
       startL=prm.TauL;
188
       startR=prm.TauR;
189
       prm.TT=2*max(prm.TauL, prm.TauR)+max(prm.durL,prm.durR);
190
       prm.dispRate=prm.TT/20; //display 20 times
191
192
   ł
193
   void disp(double t)
194
195
       if (prm.v&2) //Display hamiltonian matrix
196
197
       ł
           cout<<"# Hamiltonian"<<endl<<ham.format(comMat)<<endl;</pre>
198
199
           cout<<"# Multi-Particle Basis"<<endl;</pre>
200
           multBasis();
201
       if (prm.v&4) //plot domain walls with the time
202
203
       ł
           cout \ll t \ll ' \setminus t';
204
                                //if 1 topological region, L2 and orig2 used for \leftarrow
           if(prm.step==1)
205
              cout << (prm.orig2-prm.L2/2) << ' t ' << (prm.orig2+prm.L2/2+moveWallL( \leftrightarrow
206
                   t))<<endl;
                                 //if 2 regions, L1 and L3 are used typically \leftrightarrow
           if (prm.step==2)
207
               toward the ends
              cout << (prm.orig1-prm.L1) << ' \setminus t ' << (prm.orig1+prm.L1+moveWallL(t)) \leftrightarrow
208
                  <<'\t'
              <<(prm.orig3-prm.L3-moveWallR(t))<<'\t'<<(prm.orig3+prm.L3)<<<
209
                   endl;
210
           if (prm.step==3)
                               //if 3 regions, all are used
              \texttt{cout} <\!\!<\!\!(\texttt{prm.orig1-prm.L1}) <\!\!<\!\!' \setminus \texttt{t'} <\!\!<\!\!(\texttt{prm.orig1+prm.L1+moveWallL(t)}) \leftrightarrow
211
                  <<'\t'
              <<(prm.orig2-prm.L2/2-moveWallL(t))<<'\t'<<(prm.orig2+prm.L2/2+\leftrightarrow
212
                   moveWallR(t)) << ' \setminus t'
              <<(prm.orig3-prm.L3-moveWallR(t))<<'\t'<<(prm.orig3+prm.L3)<<<
213
                   endl:
214
       if (prm.v&8) //output chemical potential spatial distribution
215
216
           for (int i=0; i < nn; i++)
217
218
              cout<<i*latSize<<'\t'<<chem(i)<<endl;</pre>
219
220
           cout << endl ; //new index
221
222
       if (prm.v&16) //output mN eigenvalues + 1 bulk eigenvalue
223
224
       ł
           cout<<"# Eigenvalues"<<endl;</pre>
225
           for (int i=0; i < mN+1; i++)
226
227
           ł
              cout \ll eval(2*nn+i) \ll ' \setminus t';
228
           }
229
           cout<<endl<<endl; //new index
230
231
       if (prm.v&32)
                         //Majorana spatial distribution
232
233
       {
           cout<<"# Location"<<'\t'<<"u-Spin Up"<<'\t'<<"u-Spin Down"<<'\t'
234
```

```
<<"v-Spin Up"<<'\t'<<"v-Spin Down"<<endl;
235
           for (int i=0; i<nn; i++) //i iterates thru the spatial coordinate
236
           {
237
               cout<<i*latSize;</pre>
238
               for(int j=0; j<mN; j++) //j iterates thru the lowest energy \leftrightarrow
239
                   level eigenvectors
240
               ł
                  printf("\t %10.9g \t %10.9g \t %10.9g \t %10.9g",
241
                          real(symL(2*i,j)), real(symL(2*i+1,j)), real(symL(2*nn \leftrightarrow
242
                              +2*i,j), real(symL(2*nn+2*i+1,j)));
243
               }
               cout<<endl;</pre>
244
           }
245
           cout<<endl<<endl; //new index
246
247
           cout<"# Location"<<'\t'<<"u-Spin Up"<<'\t'<<"u-Spin Down"<<'\t'
248
              <<"v-Spin Up"<<'\t'<<"v-Spin Down"<<endl;
249
           for(int i=0; i<nn; i++) //i iterates thru the spatial coordinate
250
251
           ł
252
               cout<<ii*latSize;</pre>
               for(int j=0; j<mN; j++) //j iterates thru the lowest energy \leftrightarrow
253
                   level eigenvectors
254
               ł
                  printf("\t %10.9g \t %10.9g \t %10.9g \t %10.9g",
255
                          real(symR(2*i,j)), real(symR(2*i+1,j)), real(symR(2*nn \leftrightarrow
256
                              +2*i,j), real(symR(2*nn+2*i+1,j)));
               }
257
258
               cout<<endl;</pre>
           }
259
           cout<<endl<<endl; //new index
260
261
           cout<<"# Location"<<'\t'<<"u-Spin Up"<<'\t'<<"u-Spin Down"<<'\t'
262
              <<"v-Spin Up"<<'\t'<<"v-Spin Down"<<endl;
263
           for (int i=0; i<nn; i++) //i iterates thru the spatial coordinate
264
265
           Ł
               cout<<i*latSize;</pre>
266
               for(int j=0; j<mN; j++) //j iterates thru the lowest energy \leftrightarrow
267
                    level eigenvectors
268
               Ł
                   printf("\t %10.9g \t %10.9g \t %10.9g \t %10.9g",
269
                          \operatorname{norm}(\operatorname{locL}(2*i,j)), \operatorname{norm}(\operatorname{locL}(2*i+1,j)), \operatorname{norm}(\operatorname{locL}(2*nn \leftrightarrow
270
                              +2*i, j)), norm(locL(2*nn+2*i+1, j)));
               }
271
272
               cout << endl:
           }
273
           cout<<endl<<endl; //new index
274
275
           cout <<"# Location"<<'\t'<<"u-Spin Up"<<'\t'<<"u-Spin Down"<<'\t'
276
              <<"v-Spin Up"<<'\t'<<"v-Spin Down"<<endl;
277
           for (int i=0; i<nn; i++) //i iterates thru the spatial coordinate
278
279
               cout<<i*latSize;</pre>
280
               for (int j=0; j<mN; j++) //j iterates thru the lowest energy \leftrightarrow
281
                   level eigenvectors
282
               ł
                  printf("\t %10.9g \t %10.9g \t %10.9g \t %10.9g",
283
                          \operatorname{norm}(\operatorname{locR}(2*i,j)), \operatorname{norm}(\operatorname{locR}(2*i+1,j)), \operatorname{norm}(\operatorname{locR}(2*nn \leftrightarrow
284
                              +2*i, j), norm(locR(2*nn+2*i+1, j)));
285
               }
               cout<<endl;</pre>
286
```

```
}
           cout<<endl<<endl; //new index
288
289
           cout<"# Location"<<'\t'<"u-Spin Up"<<'\t'<"u-Spin Down"<<'\t'
290
              <<"v-Spin Up"<<'\t'<<"v-Spin Down"<<endl;
291
           for (int i=0; i < nn; i++) //i iterates thru the spatial coordinate
292
293
           {
               cout<<i*latSize;</pre>
294
               for(int j=0; j<mN; j++) //j iterates thru the lowest energy \leftrightarrow
295
                   level eigenvectors
296
               {
                   printf("\t %10.9g \t %10.9g \t %10.9g \t %10.9g",
297
                          \texttt{norm}(\texttt{maj}(2*\texttt{i},\texttt{j})), \texttt{ norm}(\texttt{maj}(2*\texttt{i}+1,\texttt{j})), \texttt{ norm}(\texttt{maj}(2*\texttt{nn}+2*\texttt{i} \leftrightarrow \texttt{i})))
298
                               (,j)), norm(maj(2*nn+2*i+1,j)));
               ł
299
               cout<<endl;</pre>
300
301
302
           cout << endl ; //new index
303
304
                          //check whether subgap space is unitary to test leakage
       if (prm.v&64)
305
306
       {
           307
308
309
310
   ł
311
    //finds sign of maximum absolute coefficient of vec
312
   int maxSign(VectorXcd vec){
313
       int sign = 1;
314
315
       double \max = 0;
       for (int i=0; i<nn; i++){
316
               if((norm(vec(2*i)) + norm(vec(2*i+1)))) > max) \{
317
                      \max = (\operatorname{norm}(\operatorname{vec}(2*i)) + \operatorname{norm}(\operatorname{vec}(2*i+1)));
318
                       if ((real(vec(2*i)) + real(vec(2*i+1))) > 0)
319
                          sign = 1;
320
                       else
321
                          sign = -1;
322
               }
323
324
325
       return sign;
326
327
   int maxLoc(VectorXcd vec){
328
       int loc = 0;
329
       double max = 0;
330
       for (int i=0; i<nn; i++){
331
               if ( (norm(vec(2*i)) + norm(vec(2*i+1))) > max) \{
332
                      \max = (\operatorname{norm}(\operatorname{vec}(2*i)) + \operatorname{norm}(\operatorname{vec}(2*i+1)));
333
                      loc = i;
334
               }
335
336
       ł
337
       return loc;
338
   ł
339
   double moveWallL(double t)
340
341
       return (prm.AL*0.5) *(tanh((5/prm.TauL)*(t-startL))-tanh((5/prm.TauL)*(\leftrightarrow
342
            t-startL-prm.durL)));
343 }
```

287

```
345 double moveWallR(double t)
346
       return (prm.AR*0.5)*(tanh((5/prm.TauR)*(t-startR))-tanh((5/prm.TauR)*(\leftarrow
347
           t-startR-prm.durR)));
348
349
    //returns a step that goes from 0 to 1, over length lambda, and is \leftrightarrow
350
        centered at origin
   double step(double x, double lambda, double origin)
351
352
       return 0.5*(tanh((5/lambda)*(x - origin)) + 1);
353
354
355 double delFunction(double x, double t)
356
       return prm.delta+x*0.0+t*0.0; //replace with position and/or time \leftrightarrow
357
           dependent delta if desired
358
   double magXFunction(double x, double t)
359
360
       return prm.Vx+x*0.0+t*0.0; //replace with position and/or time \leftrightarrow
361
           dependent if desired
   }
362
   double magYFunction(double x, double t)
363
364
       return prm.Vy+0.0*t+0.0*x; //replace with position and/or time \leftrightarrow
365
           dependent if desired
366
   }
   double magZFunction(double x, double t)
367
368
       return prm.Vz+0.0*t+0.0*x; //replace with position and/or time \leftrightarrow
369
           dependent if desired
370
371
   void update_chem(double t) //t is the percentage of time elapsed
372
373
       for (int i=0; i < nn; i++)
374
375
       ł
           if(prm.step==0)
376
              chem(i)=prm.mu;
377
                               //if 1 topological region, L2 and orig2 used for \leftarrow
           if(prm.step==1)
378
              the middle
              chem(i) = prm.mu*(1+(
379
                     -1*step(i*latSize, prm.lambda, prm.orig2-prm.L2/2)
380
                     +1*step(i*latSize, prm.lambda, prm.orig2+prm.L2/2+ \leftrightarrow
381
                         moveWallL(t))
                            ));
382
           if (prm.step==2)
                                //if 2 regions, L1 and L3 are used typically \leftrightarrow
383
              toward the ends
              chem(i) = prm.mu*(1+(
384
                     -step(i*latSize, prm.lambda, prm.orig1-prm.L1)
385
                     + \texttt{step}(\texttt{i*latSize}, \texttt{prm.lambda}, \texttt{prm.orig1+prm.L1+moveWallL}(\texttt{t}) \leftrightarrow \texttt{constraint}
386
                     -step(i*latSize, prm.lambda, prm.orig3-prm.L3-moveWallR(t)↔
387
                     +step(i*latSize, prm.lambda, prm.orig3+prm.L3)
388
                            ));
389
           if (prm.step==3)
                                //if 3 regions, all are used
390
              chem(i) = prm.mu*(1+(
391
                     -step(i*latSize, prm.lambda, prm.orig1-prm.L1)
392
```

344

```
+ \texttt{step}(\texttt{i*latSize}, \texttt{prm.lambda}, \texttt{prm.orig1+prm.L1+moveWallL}(\texttt{t}) \leftrightarrow
393
                      -step(i*latSize, prm.lambda, prm.orig2-prm.L2/2-moveWallL(↔
394
                           t))
                      +step(i*latSize, prm.lambda, prm.orig2+prm.L2/2+moveWallR(↔
395
                           t))
                      -\texttt{step}(\texttt{i} * \texttt{latSize}, \texttt{prm}.\texttt{lambda}, \texttt{prm}.\texttt{orig3-prm}.\texttt{L3-moveWallR}(\texttt{t}) \leftrightarrow
396
                           )
                      +step(i*latSize, prm.lambda, prm.orig3+prm.L3)
397
398
                              ));
399
       }
400
401
   void update_del(double t)
402
403
       for (int i=0; i < nn; i++)
404
405
           del(i)=delFunction(double(i),t);
406
407
408
   void update_V(double t)
409
410
   {
       for (int i=0; i < nn; i++)
411
412
       ł
           Vx(i)=magXFunction(double(i),t);
413
           Vy(i)=magYFunction(double(i),t);
414
           Vz(i)=magZFunction(double(i),t);
415
416
417 }
   void init_ham(double t) // init the Hamiltonian matrix
418
419
   ł
420
       int j;
       double temp; //used for efficiency
421
       update_chem(t);
422
       update_del(t);
423
       update_V(t);
424
       /* Helpful key for Hamiltonian basis
425
426
        * +1 changes spin
427
        * +2 moves to the right in real space
428
        * +2nn switches particle to hole
429
430
        */
431
       for (int i=0; i<nn; i++) //on-site term, chemical potential, and Z \leftrightarrow
432
           Zeeman term
       {
433
           temp=2.0*prm.t-chem(i);
434
           ham(2*i, 2*i) = temp + prm. Vz;
435
           ham(2*i+1,2*i+1)=temp-prm.Vz;
436
           ham(2*i+2*nn, 2*i+2*nn) = -1.0*temp-prm.Vz;
437
           ham(2*i+2*nn+1,2*i+2*nn+1) = -1.0*temp+prm. Vz;
438
439
       for (int i=0; i<nn-1; i++) //hopping terms
440
441
       {
           j = 2 * i + 2;
442
           ham(2*i, j) = -1.0*prm.t;
443
           ham(j, 2*i) = -1.0*prm.t;
444
445
           ham(2*i+1,j+1) = -1.0*prm.t;
           ham(j+1,2*i+1) = -1.0*prm.t;
446
           ham(2*i+2*nn, j+2*nn)=prm.t;
447
```

```
ham(j+2*nn,2*i+2*nn)=prm.t;
448
           ham (2*i+2*nn+1, j+2*nn+1)=prm.t;
449
           ham(j+2*nn+1,2*i+2*nn+1)=prm.t;
450
451
       for (int i=0; i < nn; i++) //X and Y Zeeman terms
452
453
       {
           j=2*i+1;
454
           ham(2*i,j) = complex < double > (Vx(i), -1.0*Vy(i));
455
           ham(j, 2*i) = complex < double > (Vx(i), 1.0*Vy(i));
456
           ham(2*i+2*nn, j+2*nn) = complex < double > (-1.0*Vx(i), -1.0*Vy(i));
457
           ham(j+2*nn, 2*i+2*nn) = complex < double > (-1.0*Vx(i), 1.0*Vy(i));
458
           //Eigen isn't overloaded to handle real scalar multiplication by \leftrightarrow
459
                complex vectors
460
       for (int i=0; i<nn-1; i++) //Spin-orbit terms
461
462
       {
463
           j = 2 * i + 2;
464
           temp=0.5*prm.alpha;
465
           ham(2*i+1,j)=temp;
466
           ham(j, 2*i+1) = temp;
           ham(2*i, j+1) = -1.0*temp;
467
           ham(j+1,2*i) = -1.0*temp;
468
           ham(2*i+1+2*nn, j+2*nn) = -1.0*temp;
469
           ham(j+2*nn, 2*i+1+2*nn) = -1.0*temp;
470
           ham(2*i+2*nn, j+1+2*nn) = temp;
471
472
           ham(j+1+2*nn, 2*i+2*nn) = temp;
473
       for (int i=0; i < nn; i++)
                                          //Superconductor pairing terms
474
475
       {
           j=2*i+2*nn;
476
477
           ham(2*i,j+1)=del(i);
478
           \mathtt{ham}(\mathtt{j+1}, 2 \mathtt{i}) = \mathtt{del}(\mathtt{i});
           ham(2*i+1,j) = complex < double > (-1.0,0)*del(i);
479
           ham(j, 2*i+1) = complex < double > (-1.0, 0) * del(i);
480
481
482
483
    void init_ham_sparse(double t)
                                             // init the Hamiltonian matrix as a \leftrightarrow
484
485
       int j;
486
       double temp;
                          //used for efficiency
487
       update_chem(t);
488
       update_del(t);
489
       update_V(t);
490
       for(int i=0; i<nr; i++) //on-site term, chemical potential, and Z \leftrightarrow
491
            Zeeman term
492
       {
           temp=2.0*prm.t-chem(i);
493
           tripletList.push_back(T(2*i,2*i,temp+prm.Vz));
494
           tripletList.push_back(T(2*i+1,2*i+1,temp-prm.Vz));
495
           \texttt{tripletList.push\_back}(\texttt{T}(2*\texttt{i}+2*\texttt{nn}, 2*\texttt{i}+2*\texttt{nn}, -1.0*\texttt{temp-prm}.\texttt{Vz}));
496
           \texttt{tripletList.push\_back}(\texttt{T}(2*\texttt{i}+2*\texttt{nn}+1,2*\texttt{i}+2*\texttt{nn}+1,-1.0*\texttt{temp+prm.Vz}));
497
498
       for (int i=0; i<nn-1; i++) //hopping terms
499
500
       ł
           j = 2 * i + 2;
501
           tripletList.push_back(T(2*i, j, -1.0*prm.t));
502
           tripletList.push_back(T(j, 2*i, -1.0*prm.t));
503
           tripletList.push_back(T(2*i+1,j+1,-1.0*prm.t));
504
```

```
tripletList.push_back(T(j+1,2*i+1,-1.0*prm.t));
505
                         \texttt{tripletList.push\_back}(\texttt{T}(2*\texttt{i}+2*\texttt{nn},\texttt{j}+2*\texttt{nn},\texttt{prm}.\texttt{t}));
506
                         tripletList.push_back(T(j+2*nn,2*i+2*nn,prm.t));
507
                         tripletList.push_back(T(2*i+2*nn+1,j+2*nn+1,prm.t));
508
                         tripletList.push_back(T(j+2*nn+1,2*i+2*nn+1,prm.t));
509
510
                 for (int i=0; i<nn; i++) //X and Y Zeeman terms
511
512
                 {
                         j=2*i+1;
513
                         tripletList.push_back(T(2*i,j,complex<<u>double</u>>(Vx(i),-1.0*Vy(i))));
514
515
                         tripletList.push_back(T(j, 2*i, complex < double > (Vx(i), 1.0*Vy(i))));
                         \texttt{tripletList.push\_back} (\texttt{T}(2*\texttt{i}+2*\texttt{nn},\texttt{j}+2*\texttt{nn},\texttt{complex}<\texttt{double}) (-1.0*\texttt{Vx}(\texttt{i}) \leftrightarrow \texttt{vac}) (-1.0*\texttt{Vx}(\texttt{i})) (-1.0*\texttt{Vx}(\texttt{i}
516
                                    , -1.0 * Vy(i)));
                         \texttt{tripletList.push\_back} (\texttt{T}(\texttt{j}+2*\texttt{nn}, 2*\texttt{i}+2*\texttt{nn}, \texttt{complex} < \texttt{double} > (-1.0*\texttt{Vx}(\texttt{i}) \leftrightarrow \texttt{vac}) 
517
                                    (1.0 * Vy(i)));
                         //Eigen isn't overloaded to handle real scalar multiplication by \leftrightarrow
518
                                  complex vectors
519
                 for(int i=0; i<nn-1; i++) //Spin-orbit terms</pre>
520
521
                 ł
                         j=2*i+2;
522
                         temp=0.5*prm.alpha;
523
                         tripletList.push_back(T(2*i+1,j,temp));
524
                         tripletList.push_back(T(j,2*i+1,temp));
525
                         tripletList.push_back(T(2*i, j+1, -1.0*temp));
526
                         tripletList.push_back(T(j+1,2*i,-1.0*temp));
527
                         tripletList.push_back(T(2*i+1+2*nn, j+2*nn, -1.0*temp));
528
                         tripletList.push_back(T(j+2*nn, 2*i+1+2*nn, -1.0*temp));
529
                         tripletList.push_back(T(2*i+2*nn, j+1+2*nn, temp));
530
                         tripletList.push_back(T(j+1+2*nn, 2*i+2*nn, temp));
531
532
                 for (int i=0; i < nn; i++)
533
                                                                                           //Superconductor pairing terms
534
                 {
                         j=2*i+2*nn;
535
                         tripletList.push_back(T(2*i,j+1,del(i)));
536
                         tripletList.push_back(T(j+1,2*i,del(i)));
537
                         tripletList.push_back(T(2*i+1,j,complex<double>(-1.0,0)*del(i)));
538
                         tripletList.push_back(T(j,2*i+1,complex<double>(-1.0,0)*del(i)));
539
540
541
542
         void per_ham(double k) //add periodic terms if prm.per=1
543
544
                 double temp;
                                                         //Periodic SO terms
545
                 temp=0.5*prm.alpha;
546
                ham(0, 2*nn-1) = polar(temp, k);
547
                ham(2*nn-1,0) = polar(temp, -1.0*k);
548
                ham(1, 2*nn-2) = polar(-1.0*temp, k);
549
                ham(2*nn-2,1) = polar(-1.0*temp, -1.0*k);
550
                ham(2*nn, 4*nn-1) = polar(-1.0*temp, -1.0*k);
551
                ham(4*nn-1,2*nn) = polar(-1.0*temp,k);
552
                ham(2*nn+1, 4*nn-2) = polar(temp, -1.0*k);
553
                ham(4*nn-2,2*nn+1)=polar(temp,k);
554
555
                                                                          //Periodic hopping terms
                temp = -1.0*prm.t;
556
                \mathtt{ham}(0, 2*\mathtt{nn}-2) = \mathtt{polar}(\mathtt{temp}, \mathtt{k});
557
                ham(2*nn-2,0) = polar(temp, -1.0*k);
558
                ham(1, 2*nn-1) = polar(temp, k);
559
                ham(2*nn-1,1) = polar(temp, -1.0*k);
560
                 ham(2*nn, 4*nn-2) = polar(-1.0*temp, -1.0*k);
561
```

```
ham(4*nn-2,2*nn) = polar(-1.0*temp,k);
562
       ham(2*nn+1,4*nn-1) = polar(-1.0*temp,-1.0*k);
563
       ham(4*nn-1,2*nn+1) = polar(-1.0*temp,k);
564
565
        temp=0.5*prm.alpha; //Periodic SO terms
566
        tripletList.push_back(T(0, 2*nn-1, polar(temp, k)));
567
        tripletList.push_back(T(2*nn-1,0,polar(temp,-1.0*k)));
568
        tripletList.push_back(T(1,2*nn-2,polar(-1.0*temp,k)));
569
        \texttt{tripletList.push_back}(\texttt{T}(2*\texttt{nn}-2,1,\texttt{polar}(-1.0*\texttt{temp},-1.0*\texttt{k})));
570
571
        tripletList.push_back(T(2*nn, 4*nn-1, polar(-1.0*temp, -1.0*k)));
        tripletList.push_back(T(4*nn-1,2*nn,polar(-1.0*temp,k)));
572
        \texttt{tripletList.push\_back}(\texttt{T}(2*\texttt{nn}+1, 4*\texttt{nn}-2, \texttt{polar}(\texttt{temp}, -1.0*\texttt{k})));
573
        tripletList.push_back(T(4*nn-2,2*nn+1,polar(temp,k)));
574
575
       temp = -1.0*prm.t;
                                  //Periodic hopping terms
576
        tripletList.push_back(T(0,2*nn-2,polar(temp,k)));
577
578
        tripletList.push_back(T(2*nn-2,0,polar(temp,-1.0*k)));
579
        tripletList.push_back(T(1,2*nn-1,polar(temp,k)));
580
        tripletList.push_back(T(2*nn-1,1,polar(temp,-1.0*k)));
581
        \texttt{tripletList.push_back}(\texttt{T}(2*\texttt{nn}, 4*\texttt{nn}-2, \texttt{polar}(-1.0*\texttt{temp}, -1.0*\texttt{k})));
        tripletList.push_back(T(4*nn-2,2*nn,polar(-1.0*temp,k)));
582
        tripletList.push_back(T(2*nn+1,4*nn-1,polar(-1.0*temp,-1.0*k)));
583
        \texttt{tripletList.push_back}(\texttt{T}(4*\texttt{nn}-1,2*\texttt{nn}+1,\texttt{polar}(-1.0*\texttt{temp},\texttt{k})));
584
   }
585
586
    void update_ham(double t) // update the Hamiltonian matrix in real space
587
588
    {
        double temp;
                          //used for efficiency
589
        update_chem(t);
590
        for (int i=0; i < nn; i++)
                                          //on-site term, chemical potential, and Z \leftrightarrow
591
            Zeeman terms
592
           temp=2.0*prm.t-chem(i);
593
           sparseHam.coeffRef(2*i,2*i)=temp+prm.Vz;
594
           sparseHam.coeffRef(2*i+1,2*i+1)=temp-prm.Vz; //sparse version
595
           sparseHam.coeffRef(2*i+2*nn,2*i+2*nn) = -1.0*temp-prm.Vz;
596
           sparseHam.coeffRef(2*i+2*nn+1,2*i+2*nn+1) = -1.0*temp+prm.Vz;
597
598
599
600
    void do_solve(void)
601
602
       sol.compute(ham);
603
        evec = sol.eigenvectors();
604
        eval = sol.eigenvalues();
605
        maj = evec.block(0, 2*nn, 4*nn, mN);
                                                    //starting row, starting col, row \leftrightarrow
606
        // maj vector corresponds to lowest 3 positive eigenvectors
607
        // which are stored in a 4*nn x mN matrix, with the lowest energy \leftrightarrow
608
            eigenvector on the left
609
610
        //make the particle-hole symmetric states
611
        \texttt{maj1L} = \texttt{evec.col}(2*\texttt{nn} - 1) + \texttt{evec.col}(2*\texttt{nn});
612
       \texttt{maj2L} = \texttt{evec.col}(2*\texttt{nn} - 2) + \texttt{evec.col}(2*\texttt{nn} + 1);
613
614
       maj3L = evec.col(2*nn - 3) + evec.col(2*nn + 2);
615
       \texttt{maj1R} = \texttt{evec.col}(2*\texttt{nn} - 1) - \texttt{evec.col}(2*\texttt{nn});
616
        \texttt{maj2R} = \texttt{evec.col}(2*\texttt{nn} - 2) - \texttt{evec.col}(2*\texttt{nn} + 1);
        \texttt{maj3R} = \texttt{evec.col}(2*\texttt{nn} - 3) - \texttt{evec.col}(2*\texttt{nn} + 2);
617
618
```

```
// check to which is more to the right, switch if needed
619
       if( maxLoc(maj1L) > maxLoc(maj1R) ){
620
          tempMaj = maj1L;
621
          maj1L = maj1R;
622
          maj1R = tempMaj;
623
624
       if ( maxLoc(maj2L) > maxLoc(maj2R) ) {
625
          tempMaj = maj2L;
626
627
          maj2L = maj2R;
          maj2R = tempMaj;
628
629
       if( maxLoc(maj3L) > maxLoc(maj3R) ){
630
          tempMaj = maj3L;
631
          maj3L = maj3R;
632
          maj3R = tempMaj;
633
634
635
636
       fixMat1L = projMat1L * maxSign(projMat1L*maj1L)
637
             + projMat2L * maxSign(projMat2L*maj1L)
             + projMat3L * maxSign(projMat3L*maj1L);
638
       fixMat2L = projMat1L * maxSign(projMat1L*maj2L)
639
             + projMat2L * maxSign(projMat2L*maj2L)
640
             + projMat3L * maxSign(projMat3L*maj2L);
641
      fixMat3L = projMat1L * maxSign(projMat1L*maj3L)
642
             + projMat2L * maxSign(projMat2L*maj3L)
643
             + projMat3L * maxSign(projMat3L*maj3L);
644
      fixMat1R = projMat1R * maxSign(projMat1R*maj1R)
645
             + projMat2R * maxSign(projMat2R*maj1R)
646
             + projMat3R * maxSign(projMat3R*maj1R);
647
       fixMat2R = projMat1R * maxSign(projMat1R*maj2R)
648
649
             + projMat2R * maxSign(projMat2R*maj2R)
             + projMat3R * maxSign(projMat3R*maj2R);
650
       fixMat3R = projMat1R * maxSign(projMat1R*maj3R)
651
             + projMat2R * maxSign(projMat2R*maj3R)
652
             + projMat3R * maxSign(projMat3R*maj3R);
653
654
       symL << fixMat1L*maj1L, fixMat2L*maj2L, fixMat3L*maj3L;</pre>
655
       symR << fixMat1R*maj1R, fixMat2R*maj2R, fixMat3R*maj3R;</pre>
656
657
     / \text{ symL} \ll \text{maj1L}, \text{ maj2L}, \text{ maj3L};
658
   // \text{ symR} \ll \text{maj1R}, \text{maj2R}, \text{maj3R};
659
660
       //below are the localized versions of the maj
661
      locL += projMat1L*symL*addMat1;
662
      locL += projMat2L*symL*addMat2;
663
      locL += projMat3L*symL*addMat3;
664
665
      locR += projMat1R*symR*addMat1;
666
      locR += projMat2R*symR*addMat2;
667
      locR += projMat3R*symR*addMat3;
668
669
      maj = 0.5*(locL + locR);
670
671
       //re-normalize
672
       MatrixXcd norms = MatrixXcd::Identity(3,3);
673
674
       double sum:
       for(int i=0; i < mN; i++) //i iterates thru the spatial coordinate
675
676
          ł
             sum = 0;
677
```

```
for(int j=0; j<nn; j++) //j iterates thru the lowest energy \leftrightarrow
678
                  level eigenvectors
              ł
679
                  sum += norm(maj(2*j,i)) + norm(maj(2*j+1,i)) + norm(maj(2*nn \leftrightarrow
680
                       +2*i,i) + norm(maj(2*nn+2*i+1,i));
681
              }
              norms(i,i) = 1.0 / sqrt(sum);
682
683
          }
684
       maj = maj * norms;
685
   ł
686
   int orderStates(int parity) //re-orders the maj vector to place parity \leftrightarrow
687
       state first, then right to left
688
       int rowCount=0, j=1;
689
690
       double weight, threshold = 0.02;
                                            //value of eigenstate which clearly \leftrightarrow
           marks majorana location
       MatrixXcd reOrder = MatrixXcd::Zero(mN,mN); //matrix to reorder state↔
691
            vectors
          //i is current location of state and j is final locations
692
       reOrder(parity, 0)=1; //makes the parity state the new 0
693
                                     //goes thru the wire from right to left
       for (int i=nn-1; i \ge 0; i--)
694
695
       {
          for(int state=0; state<mN; state++) //checks if each state hits \leftrightarrow
696
              threshold
697
          {
              weight = (norm(maj(2*i, state)) + norm(maj(2*i+1, state)))
                                                                             //weight \leftrightarrow
698
                  of state at current site
                     +norm(maj(2*nn+2*i, state))+norm(maj(2*nn+2*i+1, state)));
699
              if ( (threshold<weight)&&(weight<threshold+0.003)) //upper bound \leftrightarrow
700
                   to reduce sites with-in threshold
701
              {
                 cout << "Mode found at site "<<i *1.0/nn<<" for state "<< state <<↔
702
                 rowCount=0; //resets rowCount
703
                 for (int k=0; k < mN; k++)
704
                 {
705
                     if (reOrder(state,k)=complex<double>(1,0)) //counts \leftrightarrow
706
                         coefficients in row of state
                        rowCount++;
707
708
                 if (rowCount==0) //makes sure state isn't already used by \leftrightarrow
709
                     checking the matrix row and col
                 {
710
                     reOrder(state,j)=1; //makes the right-most state next \leftrightarrow
711
                         state
                     j++; //move to new row
712
                     if (j==mN)
                                 //once matrix is done, re-order maj and exit \leftrightarrow
713
                         function
714
                     {
                        maj=maj*reOrder;
715
                        return 0;
716
717
                     }
                 }
718
              }
719
          3
720
721
722
       return 1;
723 }
724
```

```
725 MatrixXcd makeU(void) //returns zero-energy states with Nambu-symmetry
726
727
       MatrixXcd temp = MatrixXcd:: Zero(4*nn, 2*mN);
       for (int i=0; i<2*nn; i++)
728
729
       {
          for (int j=0; j < mN; j++)
730
731
          {
             temp(i, mN+j)=maj(i,j); //top right comes from top
732
             temp(2*nn+i, mN+j)=maj(2*nn+i, j); //bottom right comes from \leftrightarrow
733
              temp(i, mN-1-j)=conj(maj(2*nn+i,j)); //top left comes from \leftarrow
734
                 bottom
             temp(2*nn+i, mN-1-j)=conj(maj(i,j)); //bottom left comes from \leftrightarrow
735
                  top
          }
736
       }
737
738
       return temp;
739
740
   //performs RK4 with mN vectors
741
   void rk4(void)
742
   {
743
       MatrixXcd k1 = MatrixXcd:: Zero(4*nn, mN);
744
       MatrixXcd k2 = MatrixXcd :: Zero(4*nn, mN);
745
       MatrixXcd k3 = MatrixXcd::Zero(4*nn,mN);
746
       MatrixXcd k4 = MatrixXcd::Zero(4*nn,mN);
747
       complex < double > I (0.0, -1.0*stepSize); //-i and step size combined for \leftrightarrow
748
            efficiency
       int dur=prm.tSteps*prm.TT;
749
750
       for (int tt=0; tt<dur; tt+=2)
751
       {
          k1 = I*sparseHam*maj;
752
          update_ham((tt+1.0)*stepSize);
753
          k2 = I * sparseHam * (maj+k1);
754
          k3 = I * sparseHam * (maj+k2);
755
          update_ham((tt+2.0)*stepSize);
756
          k4 = I*sparseHam*(maj+2.0*k3);
757
          maj=maj+(k1+2.0*k2+2.0*k3+k4)/3.0;
758
                                                            //only display every ↔
          if ((tt+2)\%(prm.tSteps*prm.dispRate) == 0)
759
              time unit
760
             Usub=makeU();
761
             Td=(Uinit.adjoint())*Usub;
762
             disp((tt+2.0)*stepSize);
763
          }
764
       }
765
   }
766
767
   int powInt(int x, int y) //overload pow for integers
768
769
   {
       if(x==0)
770
          return 0;
771
       if(y==0)
772
          return 1;
773
       int temp=x;
774
775
       for(int i=1;i<y;i++)</pre>
776
          temp*=x;
777
       return temp;
778 }
779
```

```
780 int biForm(int i) //converts decimal to binary
781
   {
782
       int x=0;
       for (int j=mN; j>0; j--)
783
784
       {
           if(i \ge powInt(2, j-1))
785
786
           {
              x + = powInt(10, j-1);
787
              i = powInt(2, j-1);
788
           }
789
790
       }
       return x;
791
792
   ł
793
   void dispBi(int i)
794
795
   ł
       printf("\%0*d", mN, i); //\% flag to pad with 0's to form * digits with\leftrightarrow
796
             *=mN
797
   1
798
   int decForm(int i, int length)
                                           //converts binary with length bits to \leftrightarrow
799
800
   ł
       int x=0:
801
       for (int j=length; j>0; j--)
802
803
       ł
           if (i>=powInt(10, j-1))
804
805
           {
              x = powInt(2, j-1);
806
807
              i = powInt(10, j-1);
808
           ł
809
       }
810
       return x;
811
   812
   int partNum(int i, int length)
                                           //returns number of particles in the \leftrightarrow
813
        first length bits
814
815
       int p=0;
       for(int j=mN; j>(mN-length); j--) //runs loop to check length digits on↔
816
             the left
817
       ł
           if(i \ge powInt(2, j-1))
818
           {
819
              p++;
820
              i = powInt(2, j-1);
821
           }
822
                  \operatorname{cout} \ll \operatorname{powInt}(2, j-1) \ll t' \ll x \ll t' t' \ll endl;
823
824
       }
825
       return p;
826
   }
827
   void multBasis(void) //displays multiparticle basis
828
829
   ł
       MatrixXcd temp = MatrixXcd::Zero(size,size);
830
       int rowE=0, colE=0, rowO=size/2, colO=size/2;
831
       int c,r;
832
       for(int i=0;i<size;i++)</pre>
833
834
       ł
```

```
if (partNum(i,mN)\%2==0) //checks if the bra of matrix element is \leftrightarrow
835
                even
            {
836
               r = rowE;
837
               rowE++;
838
            }
839
            else
840
841
            {
               r = row0;
842
               row0++;
843
844
            }
            for(int j=0;j<size;j++)</pre>
845
            {
846
                if(partNum(j,mN)%2==0) //checks if the ket of matrix element is↔
847
                     even
                {
848
849
                    c=colE;
850
                    colE++;
                }
851
                else
852
                {
853
                    c=col0;
854
                    col0++;
855
                }
856
                           cout << "c is "<<c<<" and r is "<<r<endl:
857
               temp(r,c)=complex<double>(biForm(i),biForm(j));
858
                           printf("\%0*d,\%0*d", mN, biForm(i), mN, biForm(j)); \leftrightarrow
859
                     //% flag to pad with 0's to form * digits with *=mN
            }
860
861
            colE = 0;
                          //reset the column count for next row
            colO=size/2;
862
863
            //no need to reset since matrix is done
        for(int i=0;i<size;i++)</pre>
864
865
        {
            cout << "# [ ";
866
            for(int j=0;j<size;j++)</pre>
867
868
            {
                \texttt{printf} \left( \verb"\%0*d, \verb\%0*d" , \verb"mN, \verb"int(real(temp(i,j))), \verb"mN, "int(imag( \leftrightarrow \texttt{mag})) \right) \right) = \texttt{mag}(\texttt{mag})
869
                    temp(i,j))); //\% flag to pad with 0's to form * digits with \leftrightarrow
                      *=mN
870
            }
            cout \ll " ] \setminus n";
871
        }
872
873 }
874
   int movesNeeded(int b, int k) //returns number of operators needed to \leftrightarrow
875
        match bra with ket
876
   {
        int xB, xK, x=0;
877
        for (int j=mN; j>0; j--)
878
879
        {
           xB=0;
880
            xK = 0;
881
            if(b \ge powInt(2, j-1))
882
883
            ł
               xB=1;
884
               b = powInt(2, j-1);
885
886
            if(k \ge powInt(2, j-1))
887
888
            Ł
```

```
xK = 1;
889
               k = powInt(2, j-1);
890
            }
891
            if(xB!=xK)
892
               x + = powInt(10, j-1);
893
894
        return partNum(decForm(x,mN), mN);
895
896
   1
897
   complex < double > multDiag(int k, int i) // calculates diagonal terms for \leftrightarrow
898
         multiparticle basis
899
        complex < double > x = 0.0;
900
       dispBi(biForm(i));
901
        \operatorname{cout} <<' \ ';
902
        dispBi(biForm(i));
903
904
        for (int j=mN; j>0; j--) // j iterates thru the bit form left to right
905
906
        ł
            if(i \ge powInt(2, j-1))
907
908
            ł
               \mathbf{x} += \texttt{conj}(\texttt{Td}(\texttt{k},\texttt{mN}-1+\texttt{j})) * \texttt{Td}(\texttt{k},\texttt{mN}-1+\texttt{j});
909
               i = powInt(2, j-1);
910
               cout <<' t'<< conj(Td(k,mN-1+j))*Td(k,mN-1+j);
911
           }
912
           else
913
            {
914
               x = conj(Td(k,mN-j))*Td(k,mN-j);
915
               \operatorname{cout} \ll \operatorname{t'} = \operatorname{conj} (\operatorname{Td}(k, mN-j)) * \operatorname{Td}(k, mN-j);
916
917
            }
918
        }
919
       cout<<endl<<endl;</pre>
920
       return x;
921
922
    //calculates coefficients for normal ordered term x, and complement y of \leftarrow
923
         multiparticle basis
    void checkCre(complex<double> &x, complex<double> &y, int k, int bra, int\leftrightarrow
924
         ket)
925
                               //int to hold the bits of the bra and ket
926
        int xKet, xBra;
        int secCreate=0;
927
        int secAnn=0;
928
        int tempBra=bra;
929
                               //use temp ints to keep bra and ket for second use \leftrightarrow
        int tempKet=ket;
930
             if needed
        dispBi(biForm(bra));
931
        \operatorname{cout} <<' \ ';
932
        dispBi(biForm(ket));
933
        cout<<endl;</pre>
934
935
        for (int j=mN; j>0; j--)
                                     //starting from the left
936
937
        ł
           xBra=0; //default binary digit is 0, unoccupied
938
           xKet = 0;
939
            if (tempBra>=powInt(2,j-1)) //check if the bit of bra state is \leftrightarrow
940
            ł
941
               xBra=1; //if so, change the digit to 1
942
                tempBra-=powInt(2, j-1); //and change tempBra for next bit test
943
```

```
944
            if(tempKet >= powInt(2, j-1))
945
            ł
946
               xKet = 1;
947
               tempKet = powInt(2, j-1);
948
949
            if((xBra-xKet) == 1)
                                     //check if bra is occupied and ket isn't
950
951
            ł
                if ((partNum(bra,mN-j)\%2)==1) //change sign if odd number of \leftrightarrow
952
                    fermion exhanges needed
953
                {
                             //i.e number of 1's to the left of the position \leftrightarrow
                   x *= -1;
954
                        being checked in the bra
                   y *= -1;
955
                ł
956
               if (secCreate==0) //if first creation operator used
957
               {
958
959
                   x*=conj(Td(k,mN+j-1)); //positive matrix entry for T_i*
960
                   y = Td(k, mN-j); //corresponding negative entry for T_--i
961
                   secCreate++;
                   bra=powInt(2, j-1);
962
                   \operatorname{cout} \ll \operatorname{"Bra} \operatorname{tKet} \operatorname{tx} \operatorname{ty} \operatorname{tT}_{i} \ast \operatorname{tT}_{-i} \operatorname{"} \ll \operatorname{endl};
963
                   dispBi(biForm(bra));
964
                   cout <<' \ t':
965
                   dispBi(biForm(ket));
966
                   967
                        mN-j)<<endl;
968
               }
969
               else
                      //if second creation operator
970
971
                ł
972
                   x = Td(k, mN-j); //negative entry T_-j
                   y = conj(Td(k,mN+j-1)); //T_j *
973
                   bra = powInt(2, j-1);
974
                   cout \ll "Bra tKet tx ty tT_-j tT_j*" \ll endl;
975
                   dispBi(biForm(bra));
976
                   \operatorname{cout} <<' \ t';
977
                   dispBi(biForm(ket));
978
                   cout <<' t'<<x<<' t'<<Td(k,mN-j)<<' t'<<conj(Td(k,mN+))
979
980
               }
981
982
                                       //check if ket is occupied and bra isn't
            if((xKet-xBra)==1)
983
984
                if ((partNum(ket,mN-j)\%2)==1) //change sign if odd number of \leftrightarrow
985
                    fermion exhanges needed
               {
986
                             //i.e number of 1's to the left of the position \leftrightarrow
                   x * = -1;
987
                       being checked in the ket
                   y *= -1;
988
989
               }
               if (secAnn==0) //if first annihilation operator used
990
               {
991
                   x*=Td(k,mN+j-1); //positive matrix entry for T_j
992
                   y*=conj(Td(k,mN-j)); //corresponding negative entry for T_-j*
993
                   \mathtt{ket} = \mathtt{powInt} (2, \mathtt{j} - 1);
994
                   secAnn++;
995
                   \operatorname{cout} \ll \operatorname{"Bra} \operatorname{tKet} \operatorname{tx} \operatorname{ty} \operatorname{tT}_j \operatorname{tT}_- j \ast \operatorname{"} \ll \operatorname{endl};
996
                   dispBi(biForm(bra));
997
```

```
\operatorname{cout} \ll \operatorname{'\!\!\!\!/} t';
998
                     dispBi(biForm(ket));
999
                     cout <<'\t'<<x<<'\t'<<Td(k,mN+j-1)<<'\t'<<conj(Td(k,↔
1000
                         mN-j))<<endl;
1001
                 }
1002
                 else
                        //if second annihilation operator
1003
                 {
1004
                     x*=conj(Td(k,mN-j)); //negative entry T_-i*
1005
                     y = Td(k, mN+j-1);
                                            //T_i
1006
                     \mathtt{ket} = \mathtt{powInt} (2, j-1);
1007
                     \operatorname{cout} \ll \operatorname{"Bra} \operatorname{tKet} \operatorname{tx} \operatorname{ty} \operatorname{tT}_{-i} * \operatorname{tT}_{i} \ll \operatorname{endl};
1008
                     dispBi(biForm(bra));
1009
                     \operatorname{cout} <<' \ ';
1010
                     dispBi(biForm(ket));
1011
1012
                     cout <<' t'<<x<<' t'<<conj(Td(k,mN-j)))<<' t'<<Td(k,mN-j))
                         +j-1 << endl;
1013
1014
                 }
             }
1015
1016
         }
1017
     ł
1018
    MatrixXcd multProj(int k) //returns N matrix in multiparticle basis
1019
1020
         MatrixXcd temp = MatrixXcd::Zero(size,size);
1021
         int rowE=0, colE=0, rowO=size/2, colO=size/2;
1022
         int c,r;
1023
         complex<double> x;
1024
1025
         complex<double> y;
1026
         for(int i=0;i<size;i++)</pre>
1027
         ł
             if (partNum(i, mN)\%2==0) //checks if the bra of matrix element is \leftrightarrow
1028
                 even
1029
             ł
                 r = rowE;
1030
                 rowE++;
1031
1032
             }
             else
1033
1034
             {
1035
                 r = row0;
                 row0++;
1036
1037
             for(int j=0;j<size;j++)</pre>
1038
1039
                             //reset x and y for each bra and ket
                 x = 1.0;
1040
                 y = 1.0;
1041
                 if(partNum(j, mN)%2==0) //checks if the ket of matrix element is↔
1042
                       even
1043
                 {
1044
                     c=colE;
1045
                     colE++;
                 }
1046
                 else
1047
                 {
1048
                     c=col0;
1049
                     col0++;
1050
                 }
1051
                 cout \ll c is "\ll c \ll" and r is "\ll r \ll c endl;
1052
                 if(i==j)
1053
```
```
1054
              ł
                  temp(r,c)=multDiag(k,i);
1055
              }
1056
              else if (movesNeeded(i,j)==2)
1057
              {
1058
                 checkCre(x, y, k, i, j);
1059
1060
                 temp(r,c)=x-y;
1061
              }
1062
              else
1063
                 \texttt{temp}(\texttt{r},\texttt{c}) = 0;
1064
           }
1065
                        //reset the column count for next row
           colE = 0;
1066
           colO=size/2;
1067
          //no need to reset since matrix is done
1068
       }
       return temp;
1069
1070
1071
    void findProb(void) //finds probabilities using all combos of projectors↔
1072
         and initState
1073
                          //matrix to hold projector product
       MatrixXcd temp;
1074
       int rowE=0, rowO=size/2; //places even probabilties in top half
1075
       int r; //holds the row after adjusting for even and odd sectors
1076
       int tempI; //temp iterator to store value of i
1077
       for(int i=0;i<size;i++) //i iterates thru the
1078
       {
1079
           temp = MatrixXcd::Identity(size,size); //initialize matrix
1080
           if (partNum(i,mN)%2==0) //checks if the row is even
1081
1082
           {
1083
              r=rowE;
1084
              rowE++;
           }
1085
           else
1086
1087
           {
              r = row0:
1088
              row0++;
1089
           }
1090
           cout <<"`i="<<i <<'\t'<<"`r="<<rdd];
1091
           tempI=i;
1092
           //determine projectors for each state
1093
           for (int j=mN; j>0; j--)
                                    //starts from the left
1094
1095
           ł
              if (tempI>=powInt(2,j-1)) //determine if jth bit is 1
1096
              {
1097
                 cout <<"tempI="<<tempI<<" j=1"<<'\t'<<mN-1+j<<endl;
1098
                  temp*=multProj(mN-1+j); //if 1, use positive projector
1099
                                            //decreases i to find binary form, \leftrightarrow
                  tempI = powInt(2, j-1);
1100
                     but doesn't affect i
              }
1101
              else //if jth bit is 0
1102
              {
1103
                  temp*=multProj(mN-j); //if 0, use negative projector
1104
                 cout <<"tempI="<<tempI<<" j=0"<<'\t'<<mN-j<<endl;
1105
              }
1106
           }
1107
1108
           prob(r)=( (initState.adjoint())*temp*initState );
1109
1110
       ł
1111 }
```

```
1112
    /* prints maj for future use so diagonalization isn't repeated for the \leftrightarrow
1113
        same parameters
     * then file is inserted as a header
1114
1115
     */
1116 void copyMaj(void)
1117
    {
        for (int i=0; i<4*nn; i++)
1118
1119
        {
            for (int j=0; j \le MN; j++)
1120
1121
               printf(maj(d,d)=13.12g; n, i, j, real(maj(i,j)));
1122
1123
            }
1124
        }
        cout<<endl;</pre>
1125
1126
    ł
1127
    int main (int argc, char *argv[])
1128
1129
    {
         init_params(argc, argv);
1130
        \texttt{init\_ham}(0.0);
1131
        init_ham_sparse(0.0);
1132
        if(prm.per==1)
1133
           per_ham(0.0);
1134
        sparseHam.setFromTriplets(tripletList.begin(), tripletList.end());
1135
        if (prm.v>3)
1136
1137
        {
            do_solve();
1138
1139
            orderStates(1);
                                 //argument is the column number of the parity \leftrightarrow
                eigenvector
1140
            copyMaj();
1141
        }
        disp(0.0);
1142
        if(prm.v>4)
1143
            Uinit=makeU();
1144
        if (prm.v>3)
1145
            rk4();
1146
        if (prm.v&128)
1147
1148
        {
1149
            findProb();
            for (int i=0; i < size / 2; i++)
1150
1151
            {
               cout \ll real(prob(i)) \ll ' \setminus t';
1152
            }
1153
1154
        }
        cout<<endl;</pre>
1155
        return 0;
1156
1157 }
```